

# Mixed Integer Programming Models for Non-Separable Piecewise Linear Cost Functions

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# Outline

- Introduction
- Modeling Piecewise Linear Functions
- Computational Results
- Conclusions

# Piecewise Linear Optimization

$$\min f_0(x)$$

*s.t.*

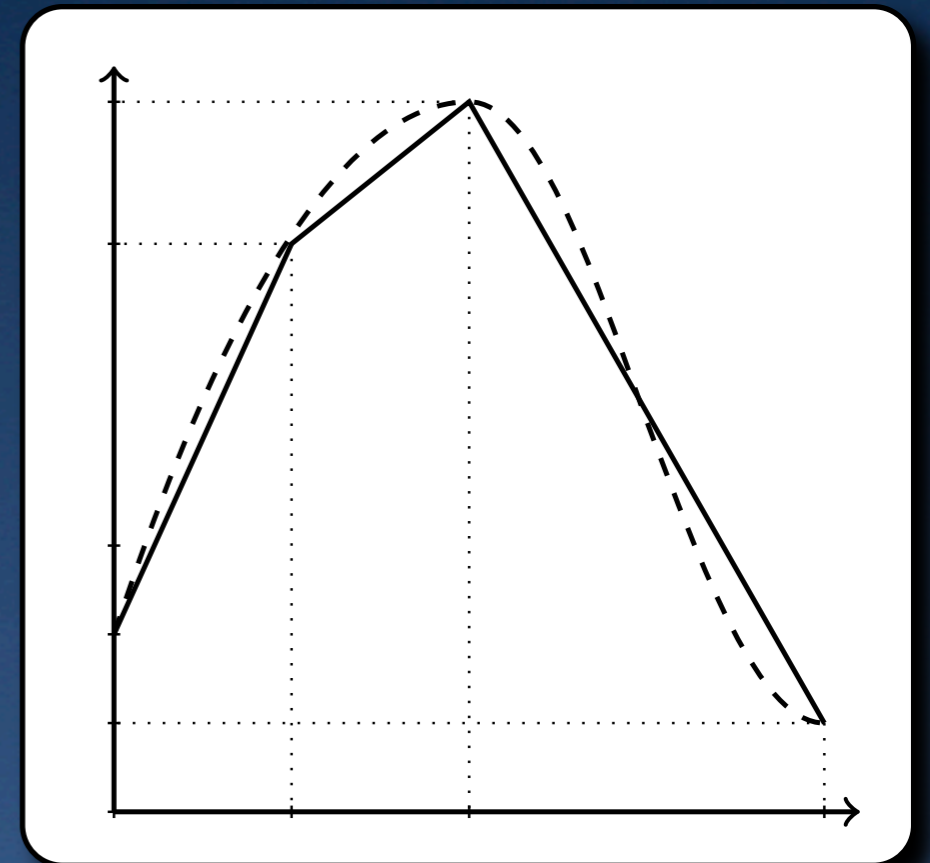
$$f_i(x) \leq 0 \quad \forall i \in I$$

$$x \in X \subset \mathbb{R}^n.$$

- $f_i(x) : D \rightarrow \mathbb{R}$  is a piecewise linear function  $\forall i \in \{0\} \cup I$ .
- $X$  is any compact set.

# Piecewise Linear Functions (PLF)

- Approximate non-linearities, discounts for volume, etc.
- Many Applications.
- Convex = Linear Programming.
- Non-Convex = NP Hard.
- Specialized algorithms (Tomlin 1981, ..., de Farias et al. 2008 ) or **Mixed Integer Programming Models** (12+ papers)





# Non-Separable = Multivariate

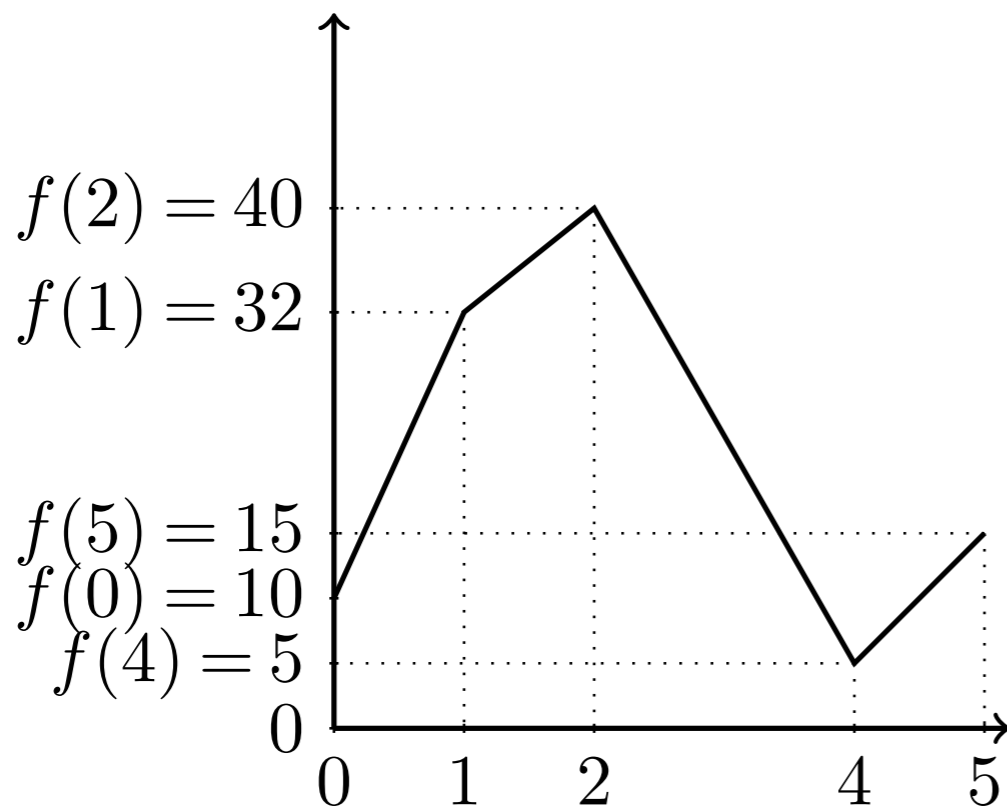
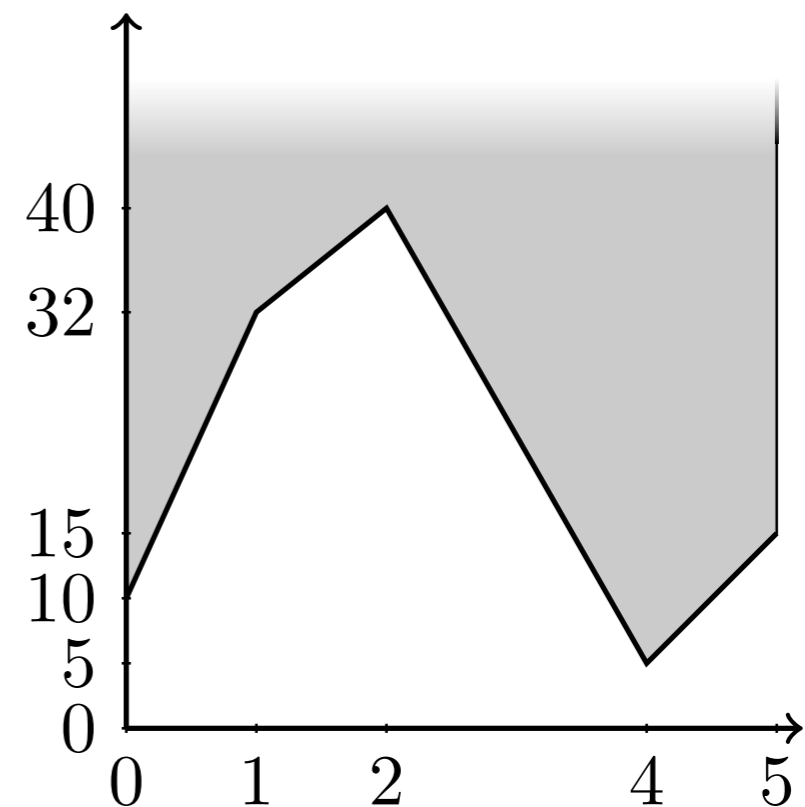
- Separable function:

$$f(x) = \sum_{j=1}^n f_j(x_j) \text{ for } f_j(x_j) : \mathbb{R} \rightarrow \mathbb{R}$$

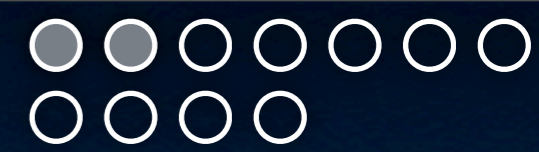
- Functions can sometimes be separated:
  - Undesirable for numerical reasons and strength.
  - Not possible for interpolated functions.

# Modeling Function = Epigraph

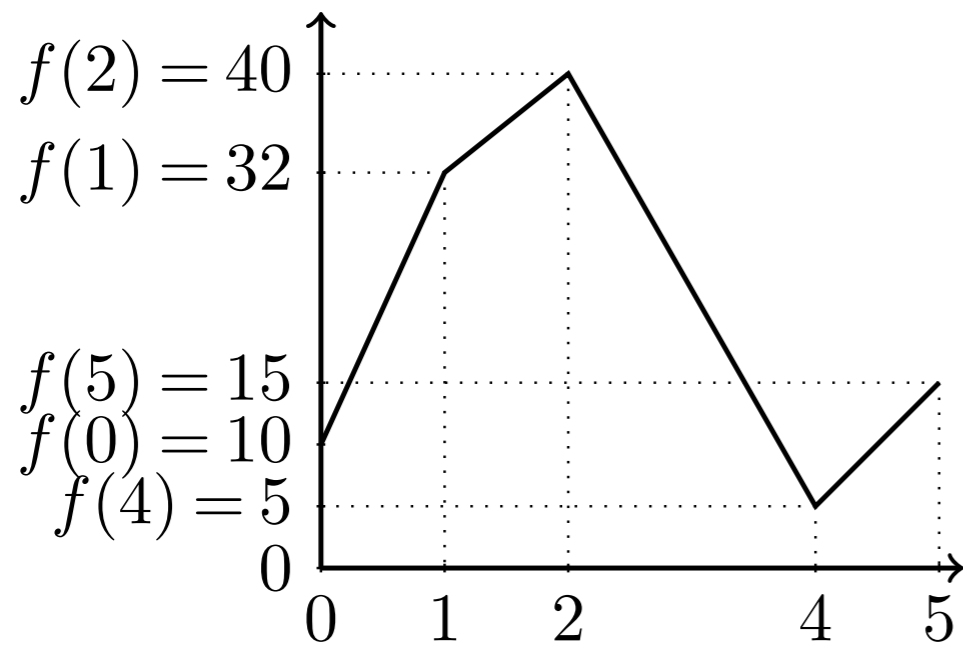
- $\text{epi}(f) := \{(x, z) \in D \times \mathbb{R} : f(x) \leq z\}$ .

(a)  $f$ .(b)  $\text{epi}(f)$ .

- Example:  $f(x) \leq 0 \Leftrightarrow (x, z) \in \text{epi}(f), z \leq 0$



# Piecewise Linear Functions: Definition



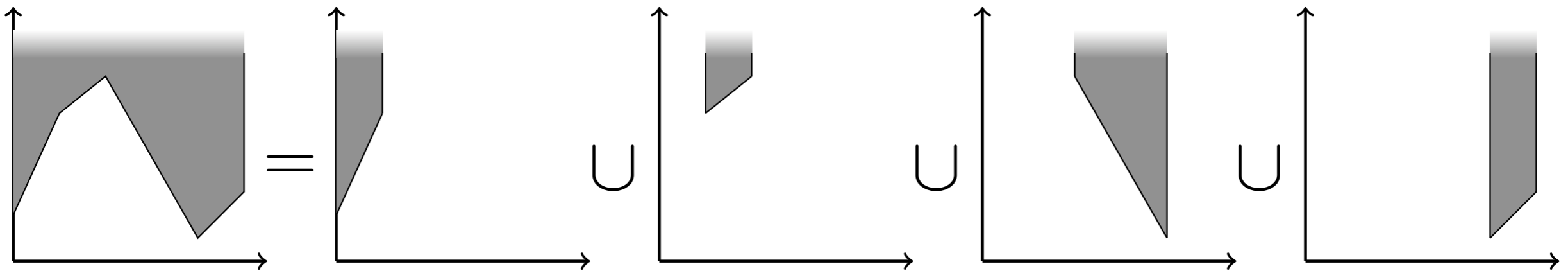
$$f(x) := \begin{cases} 22x + 10 & x \in [0, 1] \\ 8x + 24 & x \in [1, 2] \\ -17.5x + 75 & x \in [2, 4] \\ 10x - 35 & x \in [4, 5] \end{cases}$$

DEFINITION 1. Piecewise Linear  $f : D \subset \mathbb{R}^n \rightarrow \mathbb{R}$ :

$$f(x) := \begin{cases} m_P x + c_P & x \in P \quad \forall P \in \mathcal{P}. \end{cases}$$

for finite family of polytopes  $\mathcal{P}$  such that  $D = \bigcup_{P \in \mathcal{P}} P$

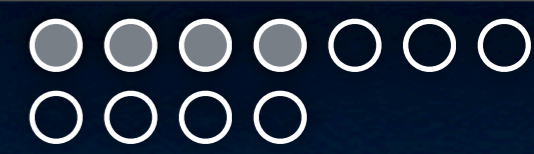
# Epigraph of PLF is Union of Polyhedra



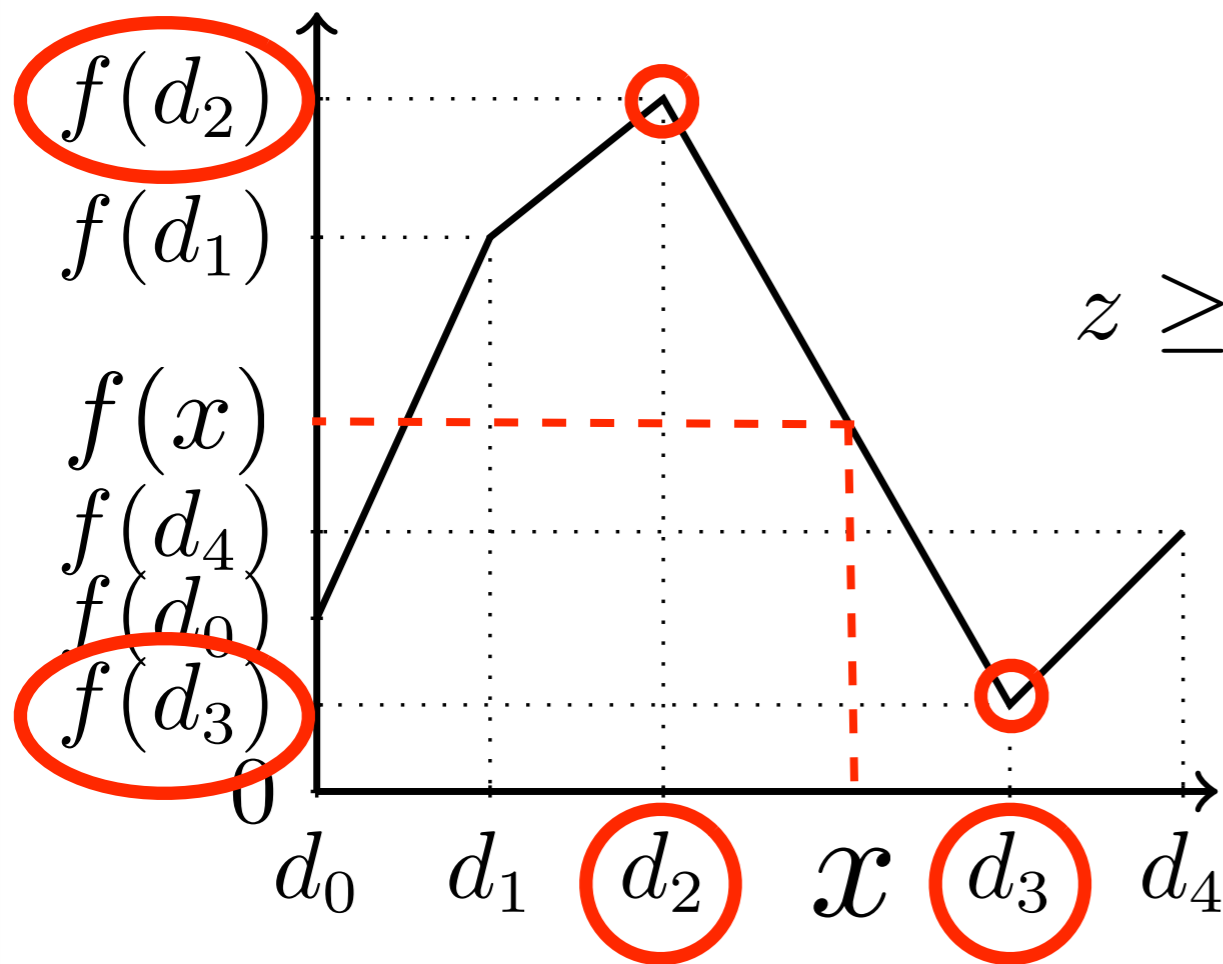
$$\begin{aligned} \text{epi}(f) &= C_n^+ + \bigcup_{P \in \mathcal{P}} \text{conv} \left( \{(v, f(v))\}_{v \in V(P)} \right) \\ &= C_n^+ + \bigcup_{P \in \mathcal{P}} \text{conv} \left( \{(v, m_P v + c_P)\}_{v \in V(P)} \right) \end{aligned}$$

$$C_n^+ := \{(0, z) \in \mathbb{R}^n \times \mathbb{R} : z \geq 0\}, \quad V(P) := \text{vertices of } P.$$





# Convex Combination Models



$$x = \lambda d_2 + (1 - \lambda) d_3$$

↓

$$z \geq f(x) = \lambda f(d_2) + (1 - \lambda) f(d_3)$$

$$(x, z) \in \text{epi}(f)$$



# Disaggregated Conv. Comb. (DCC)

$$\sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} v = x,$$

$$\sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} (m_P v + c_P) \leq z$$

$$\lambda_{P,v} \geq 0 \quad \forall P \in \mathcal{P}, v \in V(P),$$

$$\sum_{v \in V(P)} \lambda_{P,v} = y_P \quad \forall P \in \mathcal{P}$$

$$\sum_{P \in \mathcal{P}} y_P = 1,$$

$$y_P \in \{0, 1\} \quad \forall P \in \mathcal{P}.$$

- Croxton et al. (2003a), Jeroslow (1987), Jeroslow and Lowe (1984), Lowe (1984), Meyer (1976) and Serali (2001)



# Logarithmic DCC (DLog)

$$\sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} v = x, \quad \sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} (m_P v + c_P) \leq z$$

$$\lambda_{P,v} \geq 0 \quad \forall P \in \mathcal{P}, v \in V(P), \quad \sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} = 1$$

$$\sum_{P \in \mathcal{P}^+(B,l)} \sum_{v \in V(P)} \lambda_{P,v} \leq y_l, \quad \sum_{P \in \mathcal{P}^0(B,l)} \sum_{v \in V(P)} \lambda_{P,v} \leq (1 - y_l), y_l \in \{0, 1\} \quad \forall l \in L(\mathcal{P})$$

where  $B : \mathcal{P} \rightarrow \{0, 1\}^{\lceil \log_2 |\mathcal{P}| \rceil}$  is any injective function,  $L(\mathcal{P}) := \{1, \dots, \lceil \log_2 |\mathcal{P}| \rceil\}$ ,

$\mathcal{P}^+(B, l) := \{P \in \mathcal{P} : B(P)_l = 1\}$  and  $\mathcal{P}^0(B, l) := \{P \in \mathcal{P} : B(P)_l = 0\}$ .

- New? Direct from ideas in Ibaraki (1976), Vielma and Nemhauser (2008)



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$$\lambda_{P,v} \geq 0 \quad \forall P \in \mathcal{P}, v \in V(P), \quad \sum_{P \in \mathcal{P}} \sum_{v \in V(P)} \lambda_{P,v} = 1$$

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# Convex Combination (CC)

$$\begin{aligned} \sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v v &= x, & \sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v (m_P v + c_P) &\leq z \\ \lambda_v &\geq 0 \quad \forall v \in \mathcal{V}(\mathcal{P}), & \sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v &= 1 \\ \lambda_v &\leq \sum_{P \in \mathcal{P}(v)} y_P \quad \forall v \in \mathcal{V}(\mathcal{P}), & & \\ \sum_{P \in \mathcal{P}} y_P &= 1, & y_P &\in \{0, 1\} \quad \forall P \in \mathcal{P}, \end{aligned}$$

where  $\mathcal{P}(v) := \{P \in \mathcal{P} : v \in P\}$ .

- Dantzig (1963, 1960), Garfinkel and Nemhauser (1972), Jeroslow and Lowe (1985), Keha et al. (2004), Lee and Wilson (2001), Lowe (1984), Nemhauser and Wolsey (1988), Padberg (2000) and Wilson (1998)



# Logarithmic Conv. Comb. (Log)

$$\sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v v = x,$$

$$\sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v (m_P v + c_P) \leq z$$

$$\lambda_v \geq 0 \quad \forall v \in \mathcal{V}(\mathcal{P}),$$

$$\sum_{v \in \mathcal{V}(\mathcal{P})} \lambda_v = 1$$

$$\sum_{v \in L_s} \lambda_v \leq y_s,$$

$$\sum_{v \in R_s} \lambda_v \leq (1 - y_s),$$

$$y_s \in \{0, 1\} \quad \forall s \in \mathcal{S}.$$

- Requires Independent Branching Scheme.
- Vielma and Nemhauser (2008).



# Logarithmic Conv. Comb. (Log)

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# Multiple Choice (MC)

$$\sum_{P \in \mathcal{P}} x^P = x, \quad \sum_{P \in \mathcal{P}} (m_P x^P + c_P y_P) \leq z$$

$$A_P x^P \leq y_P b_P \quad \forall P \in \mathcal{P}$$

$$\sum_{P \in \mathcal{P}} y_P = 1, \quad y_P \in \{0, 1\} \quad \forall P \in \mathcal{P},$$

where  $A_P x \leq b_P$  is the set of linear inequalities describing  $P$ .

- Balakrishnan and Graves (1989), Croxton et al. (2003a), Jeroslow and Lowe (1984) and Lowe (1984)





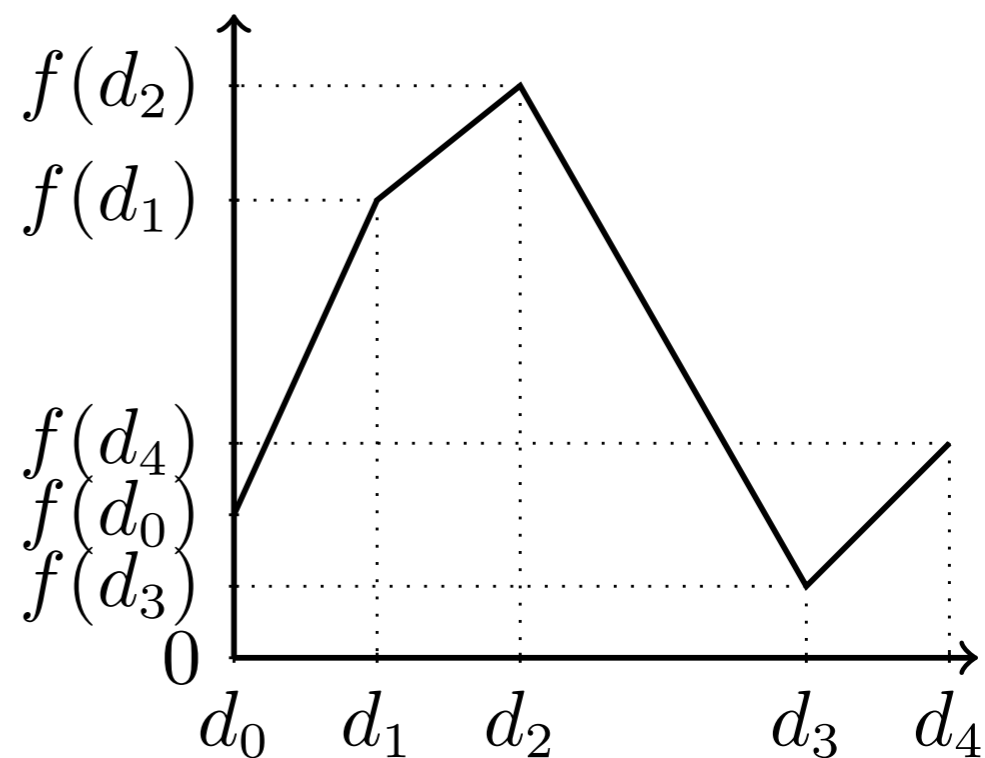
# Incremental or Delta (Inc)

$$d_0 + \sum_{k=1}^K \delta_k (d_k - d_{k-1}) = x$$

$$f(d_0) + \sum_{k=1}^K \delta_k (f(d_k) - f(d_{k-1})) \leq z$$

$$\delta_1 \leq 1, \quad \delta_K \geq 0, \quad \delta_{k+1} \leq y_k \leq \delta_k,$$

$$y_k \in \{0, 1\} \quad \forall k \in \{1, \dots, K-1\}.$$



- Similar for multivariate functions.
- Croxton et al. (2003a), Dantzig (1963, 1960), Keha et al. (2004), Markowitz and Manne (1957), Padberg (2000), Serali (2001), Vajda (1964) and Wilson (1998).

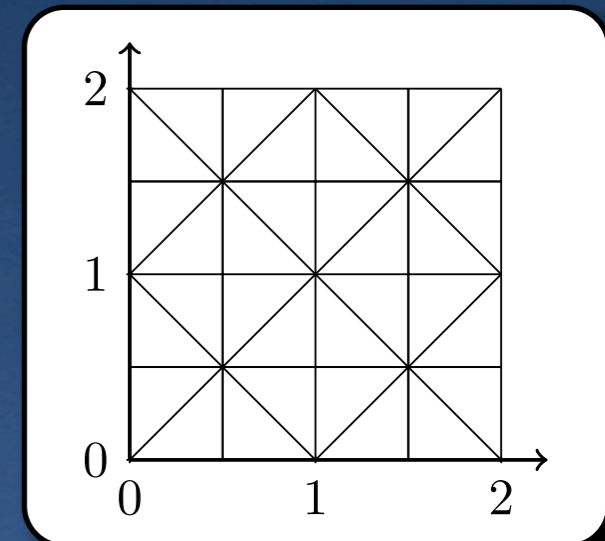


# Strength of the Models

- All models give the same LP relaxation bound:
  - LP relaxation is model of lower convex envelope (Sharp).
- In the absence of other constraints:
  - All models except for CC have integral vertices (Locally Ideal).

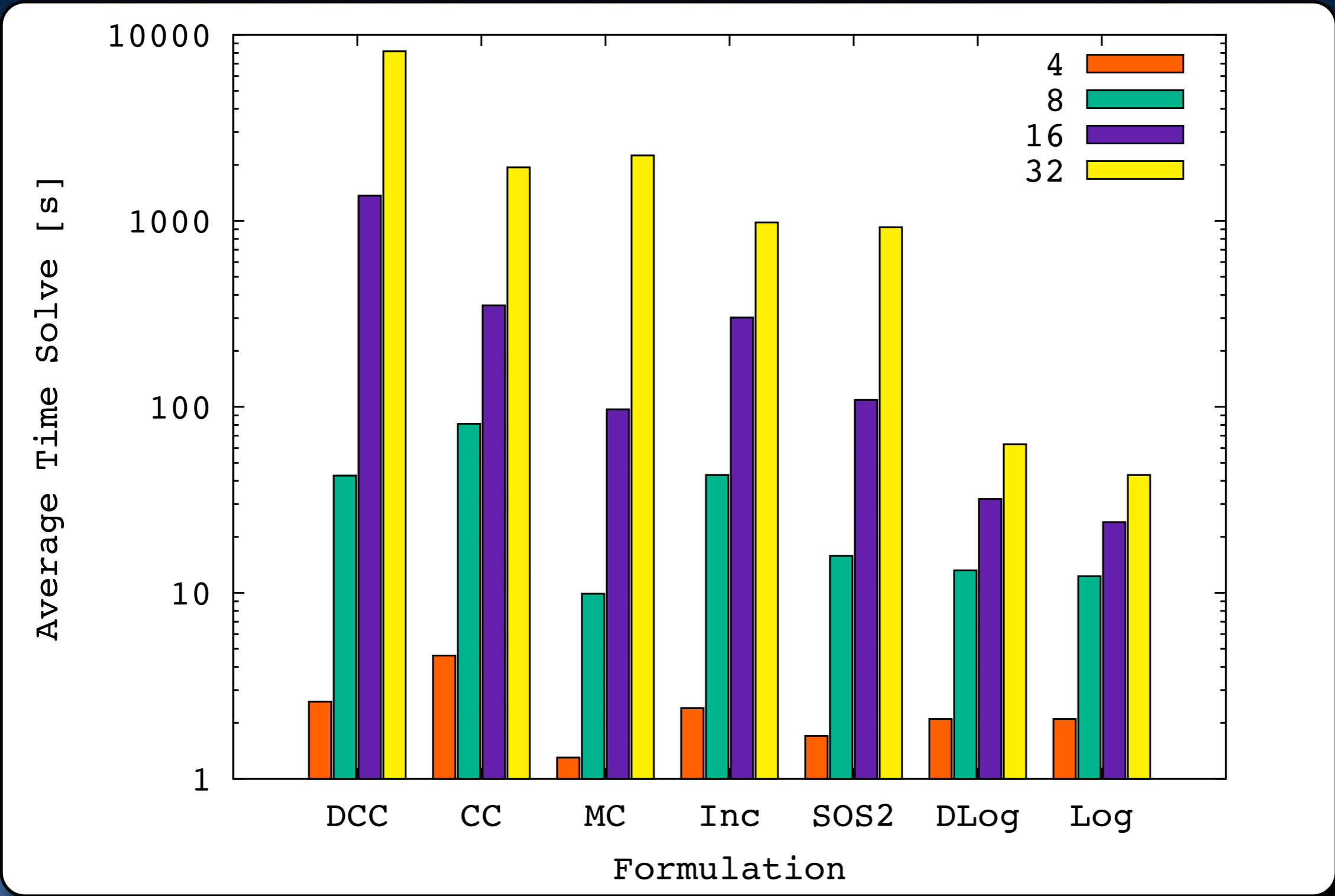
# Instances and Solvers

- Instances
  - Transportation problems (10x10 & 5x2).
  - Univariate: Concave Separable Objective.
  - Multivariate: Multi-commodity function.
  - Functions are affine in  $k$  segments or in a  $k \times k$  grid triangulation (100 instances per each  $k=4, 8, 16, 32$ ).
- Solver: CPLEX 11 on 2.4Ghz machine.



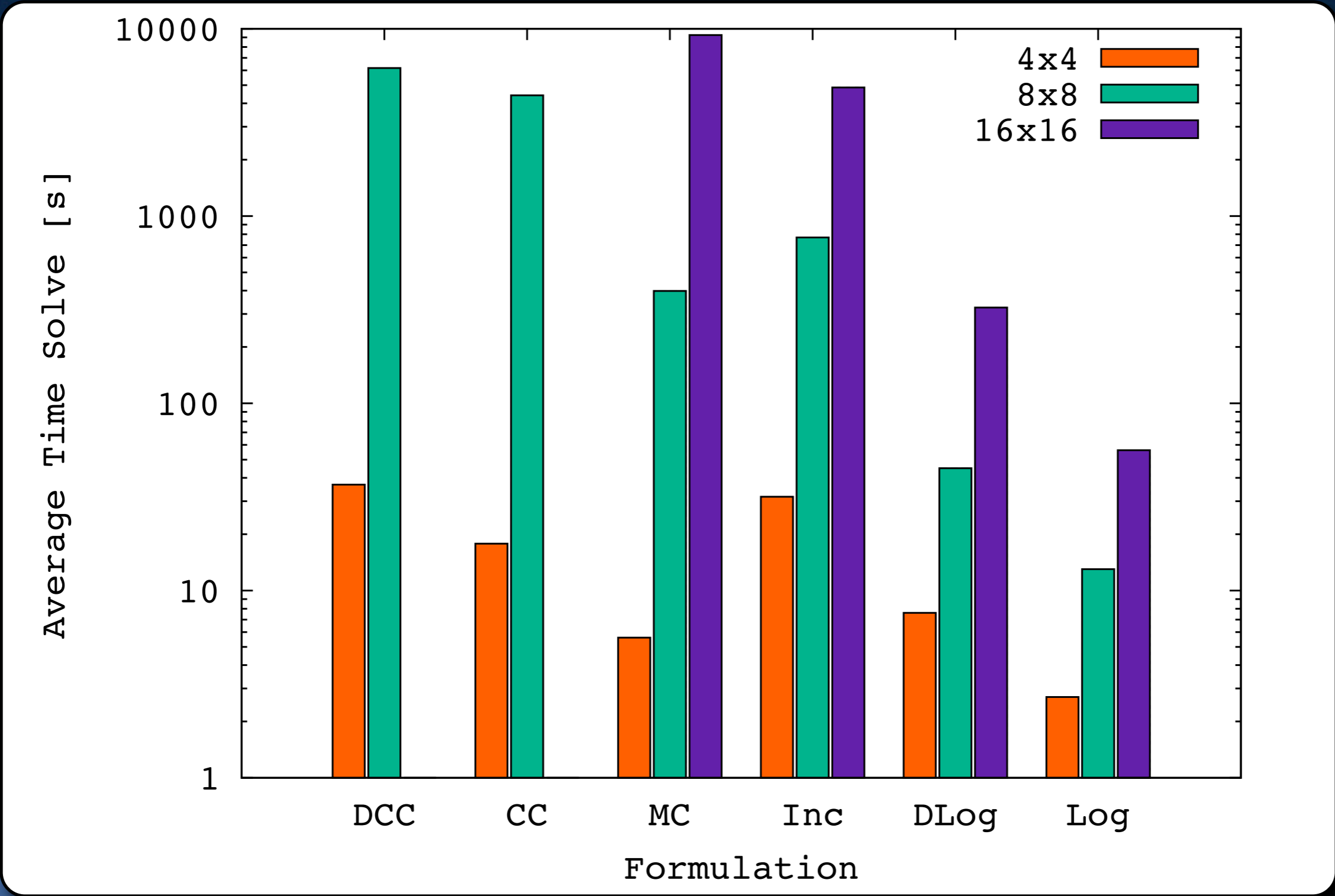


# Univariate Case (Separable)





# Multivariate Case (Non-Separable)



# Conclusions and Other Results

- Suggestions
  - For small  $k$  use MC or Inc instead of DCC or CC.
  - For large  $k$  use DLog or Log.
- DLog, DCC and MC can also be also used for Lower Semicontinuous Functions.

