

The Geometry of Nonlinear Mixed Integer Programming Formulations

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Joint work with Miles Lubin and Ilias Zadik

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Mixed Integer **Convex** Optimization (MICP)

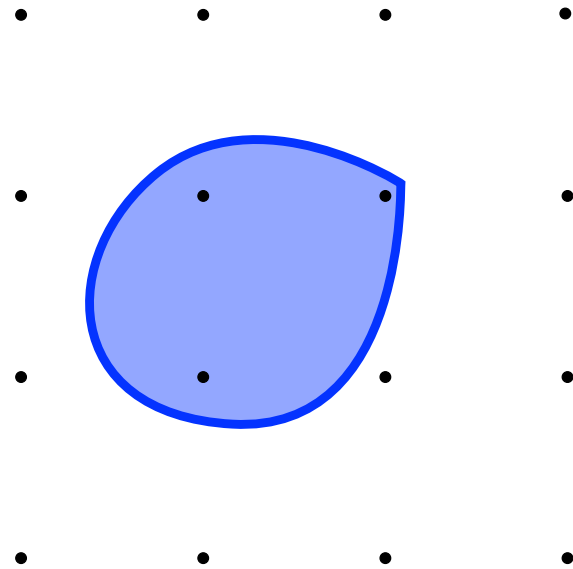
$$\min f(x)$$

s.t.

$$x \in C$$

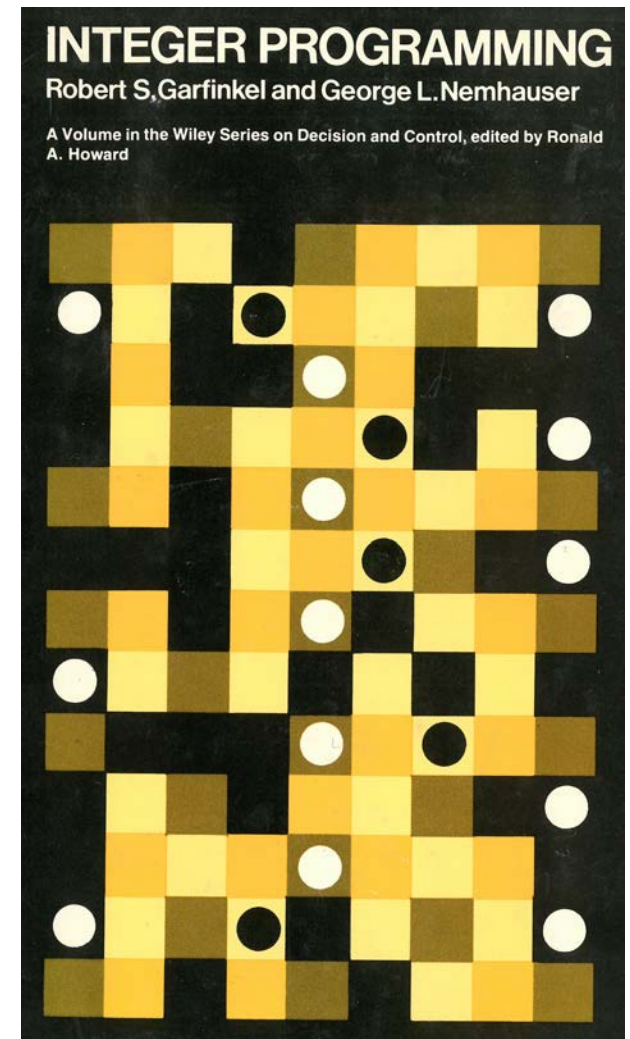
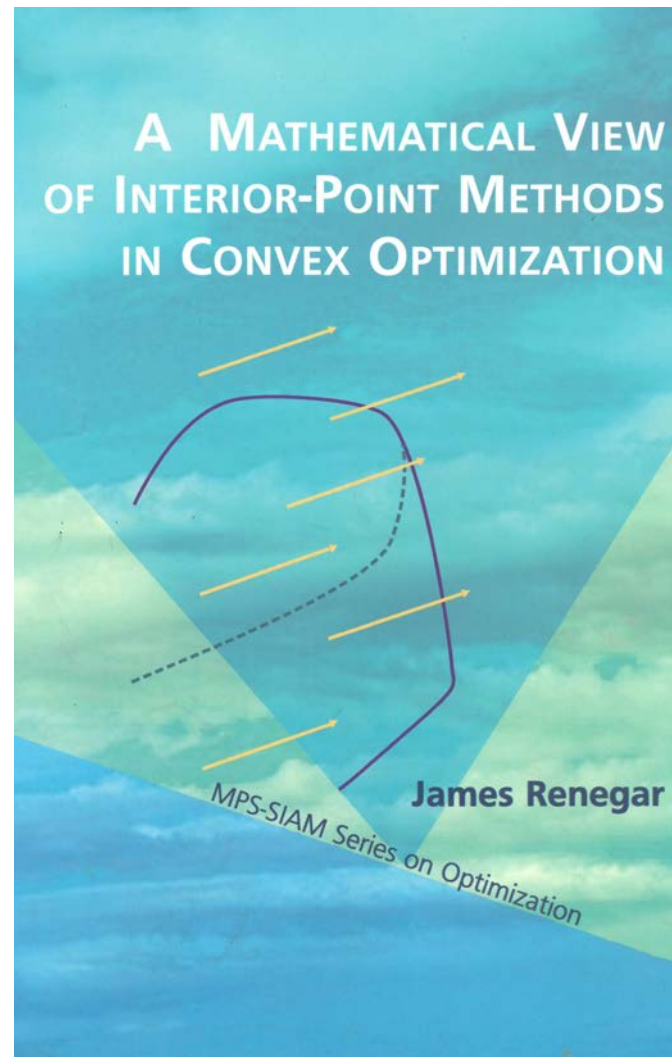
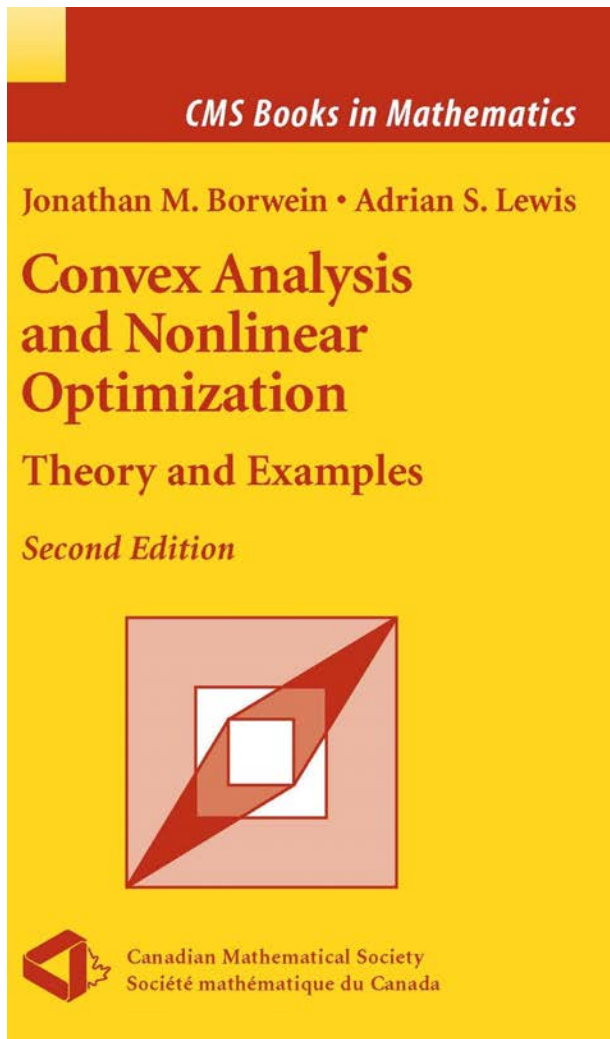
$$x_i \in \mathbb{Z} \quad i \in I$$

Convex f and C .



- Examples:
 - MI-Second Order Cone Programming (MISOCP)
 - MI-Semidefinite Programming (MISDP)

Standing On The Shoulders of Gigants



MI-SOCP Solvers and Applications

- Effective and improving solvers:



- Gurobi v6.0 (2014) – v6.5 (2015) MISOCP: 4.43 x (V., Dunning, Huchette, Lubin, 2017)

- Applications:

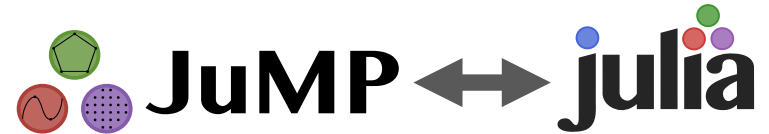
- Portfolio optimization, pricing, regression, experimental design, etc.



<http://www.gurobi.com/company/example-customers>

MI-SDP (& other cones) Solvers and Applications

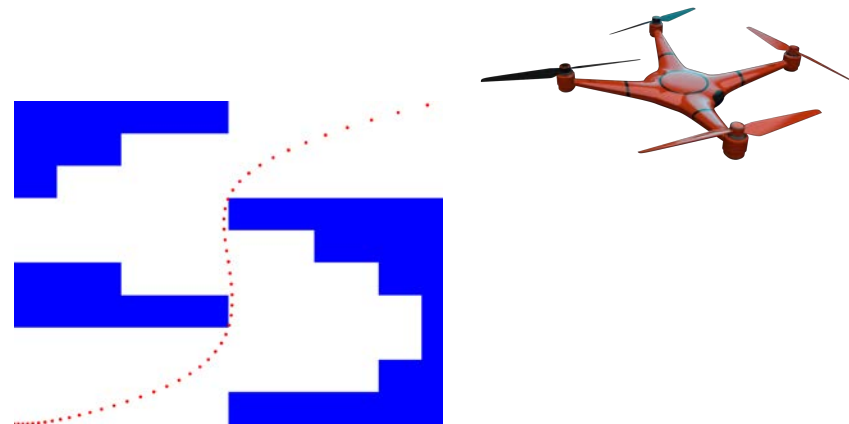
- Emerging Solvers:



Coey, Bent, Lubin, V. and Yamangil '16

- Applications:

- Collision avoidance with mixed integer sum-of-squares for optimal control of polynomial trajectories



Outline

- General Mixed Integer Convex Representability
 - What can be modeled with MICP ?
 - Joint work with Miles Lubin and Ilias Zadik
- 0-1 Mixed Integer Convex Representability
 - Unions of Convex Sets
 - Small and Strong Formulations

MICP Formulations and Representability

- A set $S \subseteq \mathbb{R}^n$ is MICP representable (MICPR) if it has an MICP formulation:
 - A closed convex set $M \subseteq \mathbb{R}^{n+p+d}$
 - auxiliary continuous variables $y \in \mathbb{R}^p$
 - auxiliary integer variables $z \in \mathbb{Z}^d$

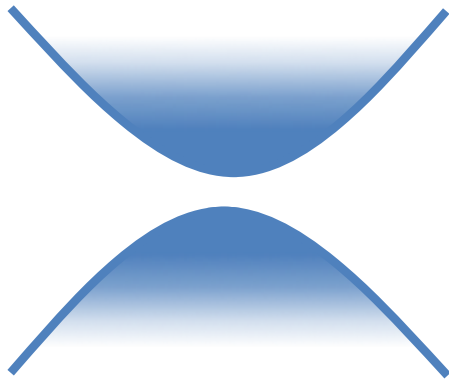
$$x \in S \quad \Leftrightarrow \quad \begin{array}{l} \exists (y, z) \in \mathbb{R}^p \times \mathbb{Z}^d \text{ s.t.} \\ (x, y, z) \in M \end{array}$$

or equivalently

$$S = \text{proj}_x \left(M \cap \left(\mathbb{R}^{n+p} \times \mathbb{Z}^d \right) \right)$$

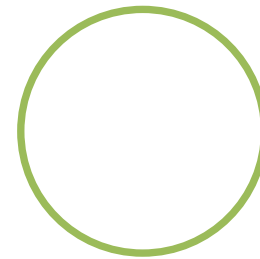
What Sets are MICP Representable (MICPR) ?

Two sheet hyperbola?



$$\{x \in \mathbb{R}^2 : 1 + x_1^2 \leq x_2^2\}$$

Spherical shell?



$$\{x \in \mathbb{R}^2 : 1 \leq \|x\| \leq 2\}$$

- Discrete subsets of the real line or natural numbers:
 - Dense discrete set? $\{\sqrt{2}x - \lfloor \sqrt{2}x \rfloor : x \in \mathbb{N}\} \subseteq [0, 1]$
 - Set of prime numbers?

A Simple Lemma for non-MICP Representability

- Obstruction for MICP representability of S :

$$\text{infinite } R \subseteq S \quad \text{s.t.} \quad \frac{u+v}{2} \notin S \quad \forall u, v \in R, u \neq v$$

Proof: Assume for contradiction there exists M such that:

$$S = \text{proj}_x \left(M \cap \left(\mathbb{R}^{n+p} \times \mathbb{Z}^d \right) \right)$$

$$\begin{aligned} (u, y_u, z_u) \in M \\ (v, y_v, z_v) \in M \end{aligned} \quad \Rightarrow \quad \frac{z_u + z_v}{2} \notin \mathbb{Z}^d$$

$$z_u \equiv z_v \pmod{2} \text{ component-wise} \quad \Rightarrow \quad \frac{z_u + z_v}{2} \in \mathbb{Z}^d$$

$$\text{component-wise parity classes} = 2^d < |R| = \infty \quad \Rightarrow \neq$$

Examples of non-MICPR Sets From Lemma

✗ Spherical shell $\{x \in \mathbb{R}^2 : 1 \leq \|x\| \leq 2\}$

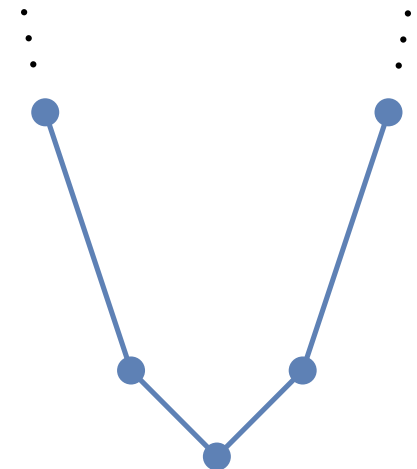


✗ Set of prime numbers

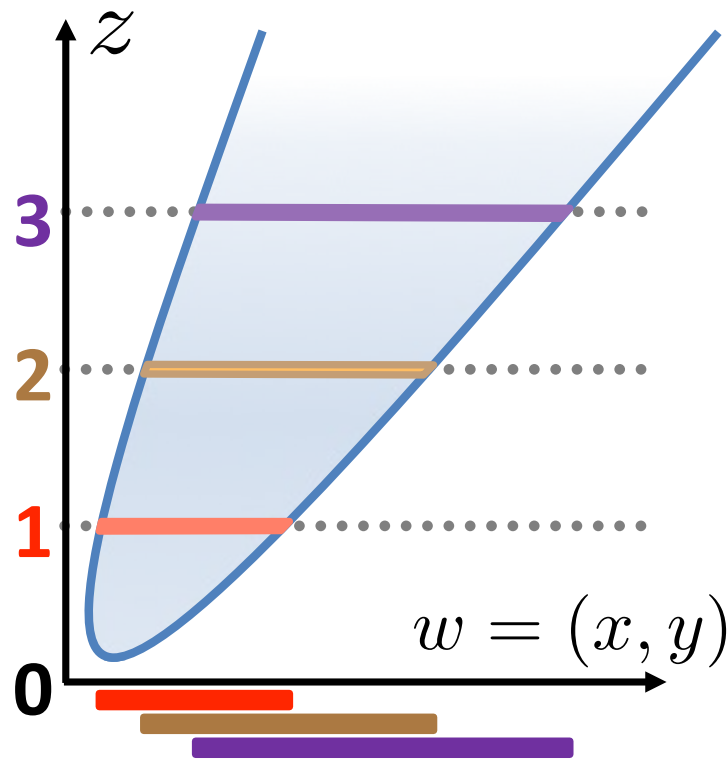
- However **prime numbers** has a **non-convex** polynomial **integer** programming formulation

✗ Set of Matrices of rank at most k

✗ Piecewise linear interpolation of x^2 at all integers



MICPR = Convex Sets Indexed by Integers in Convex



closed convex M

$$S = \text{proj}_x (M \cap (\mathbb{R}^{n+p} \times \mathbb{Z}^d))$$



$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

I convex and $B_z : \mathbb{R}^d \Rightarrow \mathbb{R}^{n+p}$ closed and convex:

- $\left(w_m \in B_{z_m}, \quad (w_m, z_m) \xrightarrow{m} (w, z) \right) \Rightarrow w \in B_z$
- $\lambda B_z + (1 - \lambda) B_{z'} \subseteq B_{\lambda z + (1 - \lambda) z'}$

MI-Linear Programming: Rational Polyhedral M

$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z) = \bigcup_{z \in I \cap \mathbb{Z}^d} P_z$$

- P_z = rational polyhedra with the same recession cone
- Representation simplifies to (Jeroslow and Lowe '84):
 - $S = \bigcup_{i=1}^k P_i + \left\{ \sum_{i=1}^t \lambda_i r^i : \lambda \in \mathbb{Z}_+^t \right\}$
 - P_i = rational polytopes
 - Very regular infinite union
- Bounded or 0-1 MILPR / MICPR = Bounded I = Finite union
 - MILPR: of polyhedra with the same recession cone
 - MICPR: of non-polyhedral convex sets ...

Extra from MICP 1: Non-Polyhedral Unions

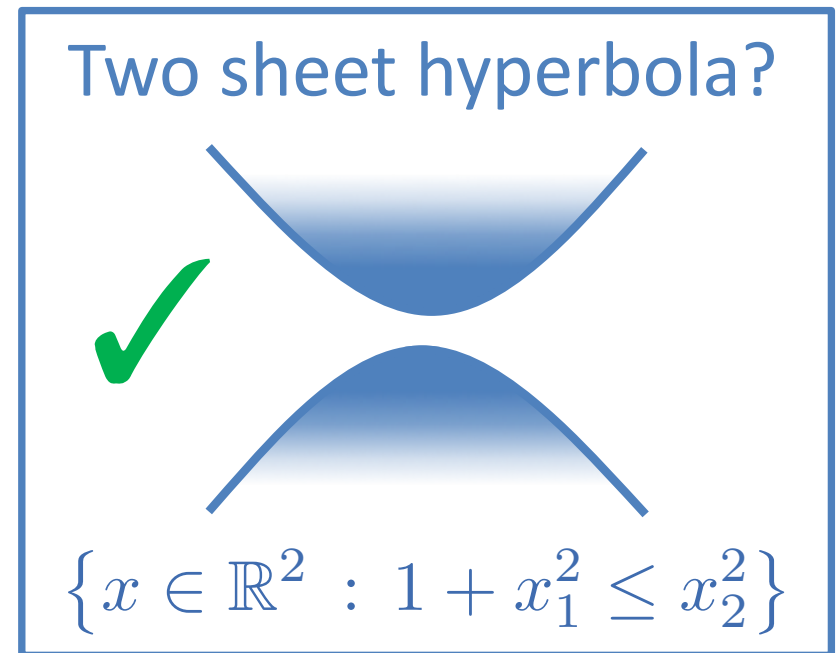
$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

1. Unions of **Non-Polyhedral sets**

Plus Projection:

2. Unions of **non-closed sets**

3. Unions of **convex sets with different recession cones**



$$B_z = \left\{ (x, y) \in \mathbb{R}^{n+1} : x \in C_z, \quad \|x\|_2^2 \leq y \right\}$$

Extra from MICP 2: Non-Polyhedral Index Set

$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$



"God made the integers, all else is the work of man"
- Leopold Kronecker

- Integers + **non-rational** unbounded **ray** = Trouble !

✓ Dense discrete set $\left\{ \sqrt{2}x - \lfloor \sqrt{2}x \rfloor : x \in \mathbb{N} \right\} \subseteq [0, 1]$

$$\|(z_1, z_1)\|_2 \leq z_2 + 1, \quad \|(z_2, z_2)\|_2 \leq 2z_1, \quad x_1 = y_2 - z_2,$$

$$\|(z_1, z_1)\|_2 \leq y_1, \quad \|(y_1, y_1)\|_2 \leq 2z_1, \quad z \in \mathbb{Z}^2$$

One (Somewhat Extreme) Way to Add Regularity

$$S = \text{proj}_x \left(M \cap \left(\mathbb{R}^{n+p} \times \mathbb{Z}^d \right) \right), \quad I = \text{proj}_z (M)$$

- $M = B + K$:
 - B compact convex set
 - K rational polyhedral cone
- Then
 - $S = \bigcup_{i=1}^k C_i + \left\{ \sum_{i=1}^t \lambda_i r^i : \lambda \in \mathbb{Z}_+^t \right\}$
 - $C_i =$ compact convex sets
- Less extreme, but still well behaved “**Rational** MICPR”:
 - Any **rational** affine mapping of index set I is bounded, or has a **rational** recession direction

Bounded or 0-1 MICP Formulations for Unions of Convex Sets

A Classical Strong Formulation for $\bigcup_{i=1}^k C_i$

$$C_i = \{x \in \mathbb{R}^n : A^i x \preceq_i b^i\}, \quad C_i^\infty = C_j^\infty$$

$$A^i x^i \preceq_i b^i z_i, \quad \forall i \in [k]$$

$$\sum_{i=1}^k x^i = x,$$

$$\sum_{i=1}^k z_i = 1, \quad z \in \{0, 1\}^k$$

$$x, x^i \in \mathbb{R}^n, \quad \forall i \in [k]$$

- Auxiliary continuous variables are copies of original variables
 - $y = (x^i)_{i=1}^k$
- “Ideal” Formulation Strength:
 - Extreme points of continuous relaxation satisfy **integrality constraints** on z
 - **Variable copies** crucial here, but **slow down computations** (usually worse than Big-M)

Generic Geometric Formulation = Gauge Functions

- For C such that $\mathbf{0} \in \text{int}(C)$ let:

$$\gamma_C(x) := \inf\{\lambda > 0 : x \in \lambda C\}$$

$$\text{epi}(\gamma_C) = \text{cone}(C \times \{1\})$$

- If $b^i \in C_i$ then **ideal** formulation:

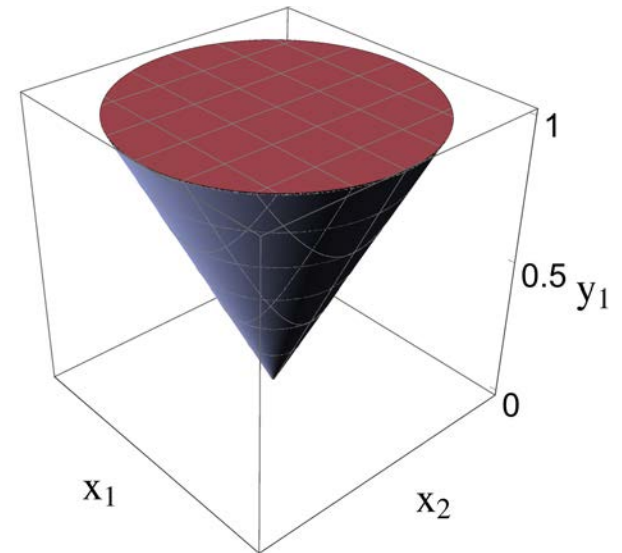
$$\gamma_{C^i - \{b^i\}}(x^i - z_i b^i) \leq z_i \quad \forall i \in [k]$$

$$\sum_{i=1}^k x^i = x$$

$$\sum_{i=1}^k z_i = 1$$

$$z \in \{0, 1\}^k$$

$$x, x^i \in \mathbb{R}^n \quad \forall i \in [k]$$



Simple Ideal Formulation without Variable Copies

- Unions of (nearly) Homothetic Closed Convex Sets (V. 17):

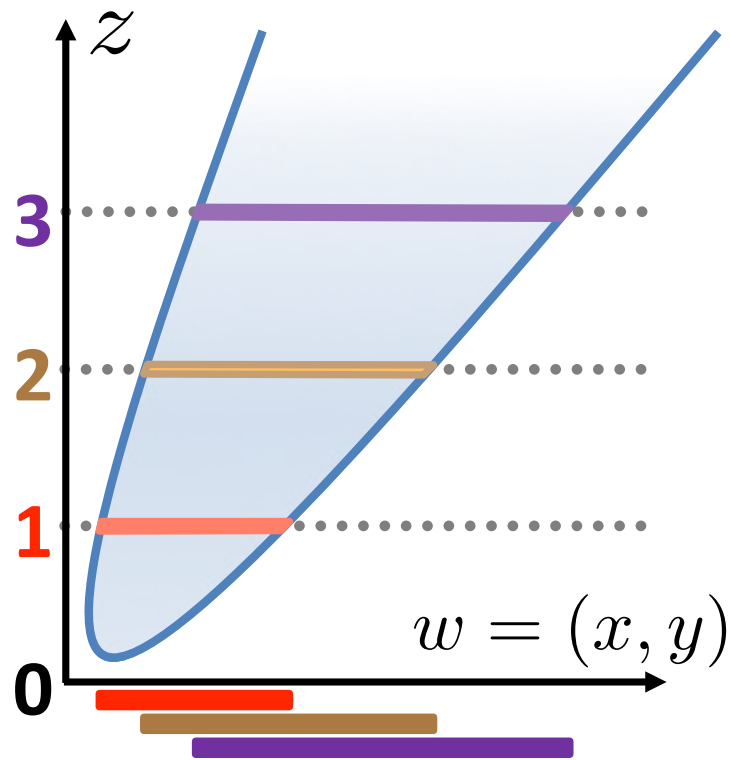
$$C_i = \lambda_i C + b^i + C^\infty$$



$$\gamma_C \left(x - \sum_{i=1}^n z_i b^i \right) \leq \sum_{i=1}^n \lambda_i z_i$$
$$\sum_{i=1}^n z_i = 1, \quad z \in \{0, 1\}^n$$

≈ to polyhedral results from Balas '85, Jeroslow '88 and Blair '90

Embedding Formulation Construction



$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x(B_z)$$

↓

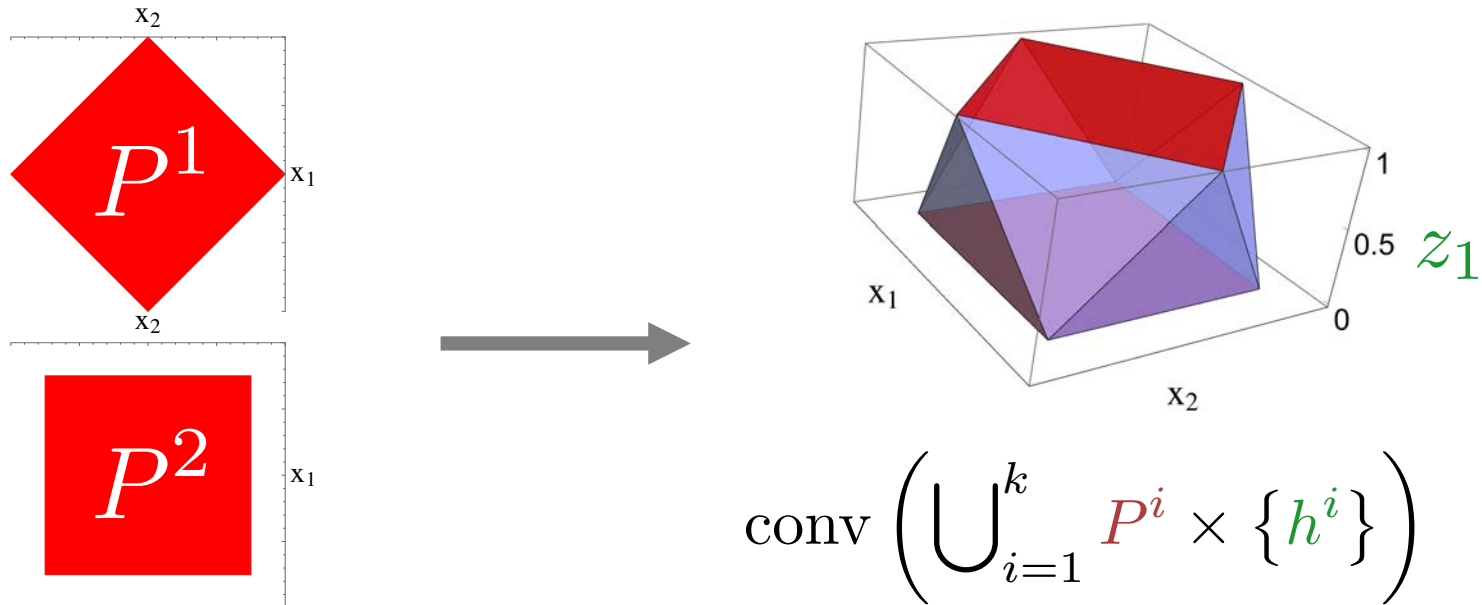
$$M = \overline{\text{conv}} \left(\bigcup_{z \in I \cap \mathbb{Z}^d} B_z \times \{z\} \right)$$

I convex and $B_z : \mathbb{R}^d \Rightarrow \mathbb{R}^{n+p}$ closed and convex:

$$C_z^\infty = C_{z'}^\infty \quad \bullet \quad (w_m \in B_{z_m}, (w_m, z_m) \xrightarrow{m} (w, z)) \Rightarrow w \in B_z$$

$$I \cap \mathbb{Z}^d \subseteq \{0, 1\}^d \quad \bullet \quad \lambda B_z + (1 - \lambda) B_{z'} \subseteq B_{\lambda z + (1 - \lambda) z'}$$

Embedding Formulation = Automatically **Ideal**



- Originally for Polyhedra (V. '17)
 - Small size with careful choice of **encoding** $\{h^i\}_{i=1}^k \subseteq \{0, 1\}^d$
- Extensions to **general integers**, practical construction techniques, computations, applications and software tools:
 - Huchette and V. '17a,b,c; Huchette, Dey and V. '17

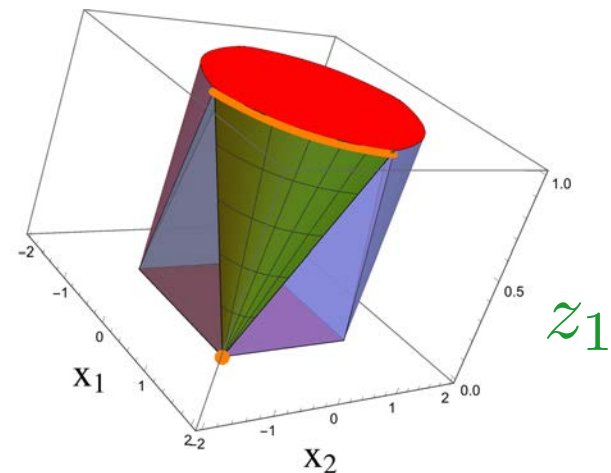
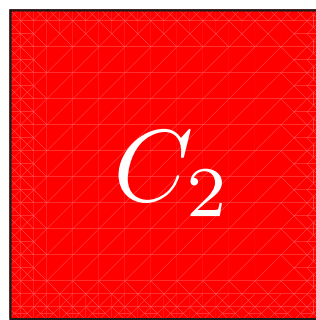
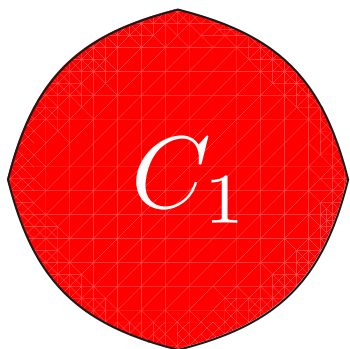


Focus for Non-Polyhedral Embedding Formulations

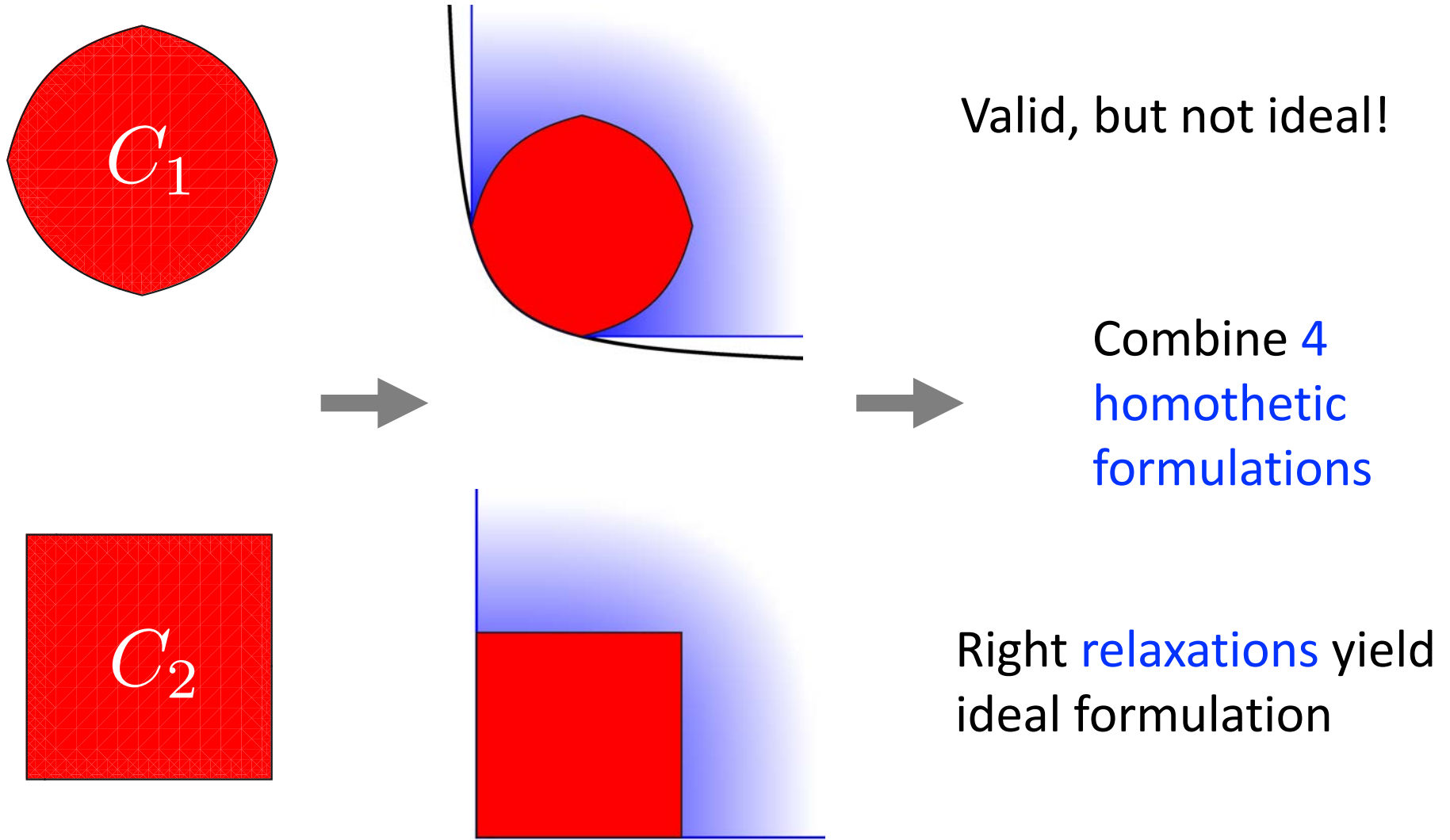
- **Unary encoding:** $\{h^i\}_{i=1}^k = \{e^i\}_{i=1}^k \subseteq \{0, 1\}^k$
 - Related to Cayley Embedding for Minkowski sums
 - Homothetic formulation

$$\gamma_C \left(x - \sum_{i=1}^n z_i b^i \right) \leq \sum_{i=1}^n \lambda_i z_i$$
$$\sum_{i=1}^n z_i = 1, z \in \{0, 1\}^n$$

- How to write convex hull:

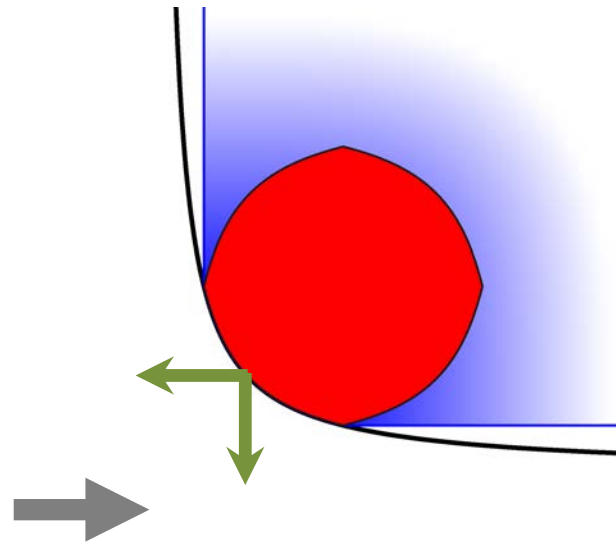
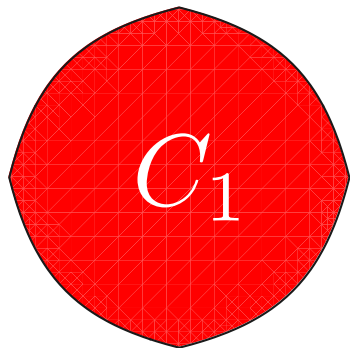


Sticking Homothetic Formulations Together



Sufficient Conditions For Ideal Formulation

$$\sigma_S(u) := \sup\{u \cdot x : x \in S\}$$



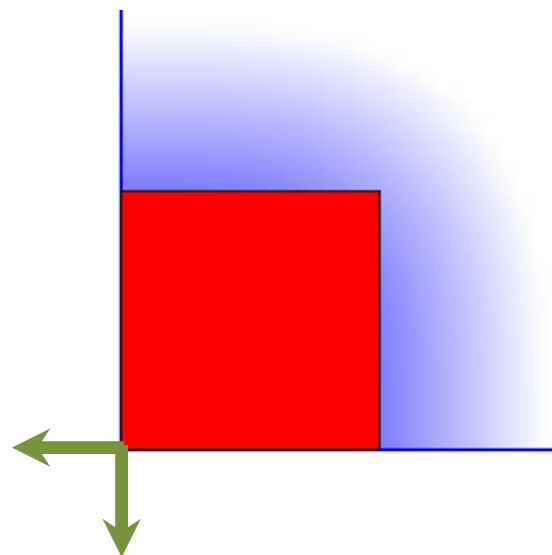
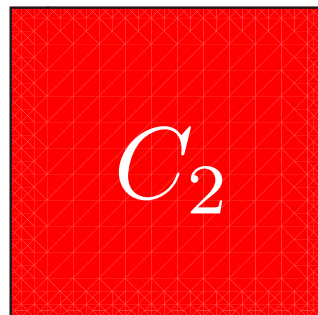
C_1^j

$$\forall u \in \mathbb{R}^n \quad \exists j$$

s.t.

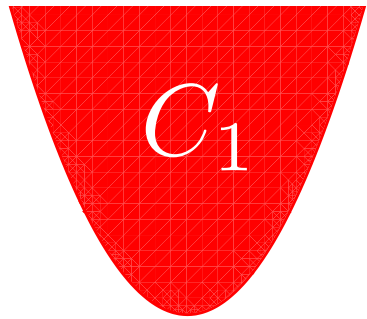
$$\sigma_{C_i}(u) = \sigma_{C_i^j}(u)$$

$$\forall i \in \{1, 2\}$$



C_2^j

May Need to “Find” Homothetic Constraints



$$x_1^2 \leq x_2 \leq 1$$

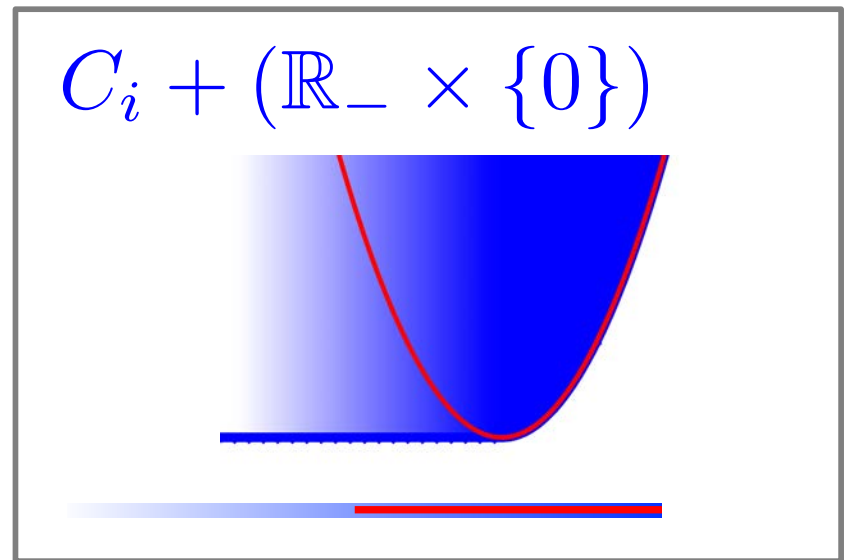
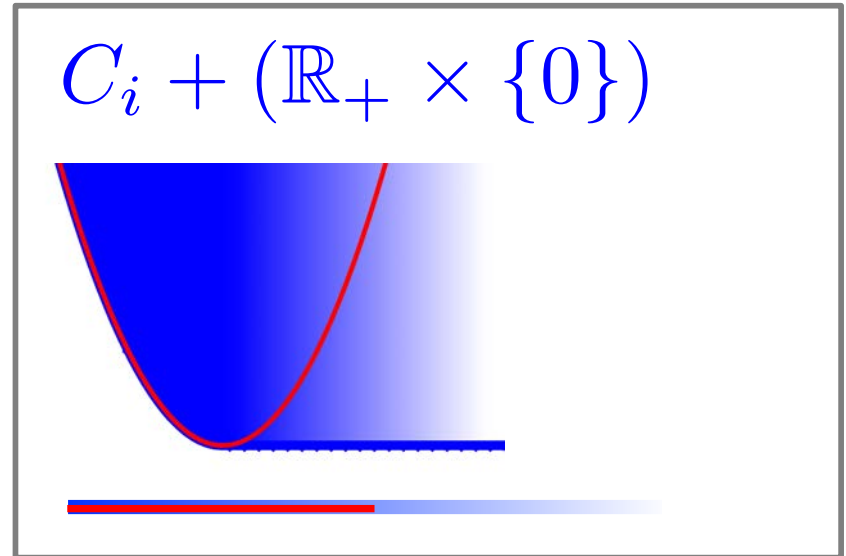


$$[-1, 1] \times 0$$

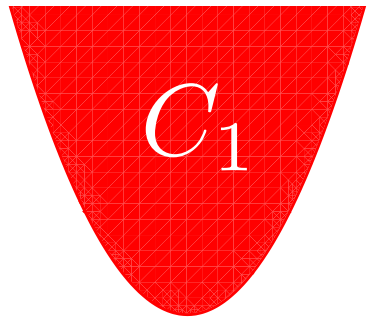
$$C_1 + (\mathbb{R}_+ \times \{0\}) :$$

$$(\max\{x_1, 0\})^2 \leq x_2 \leq 1$$

Similar to Bestuzheva et al.
'16 who divide sets in two.



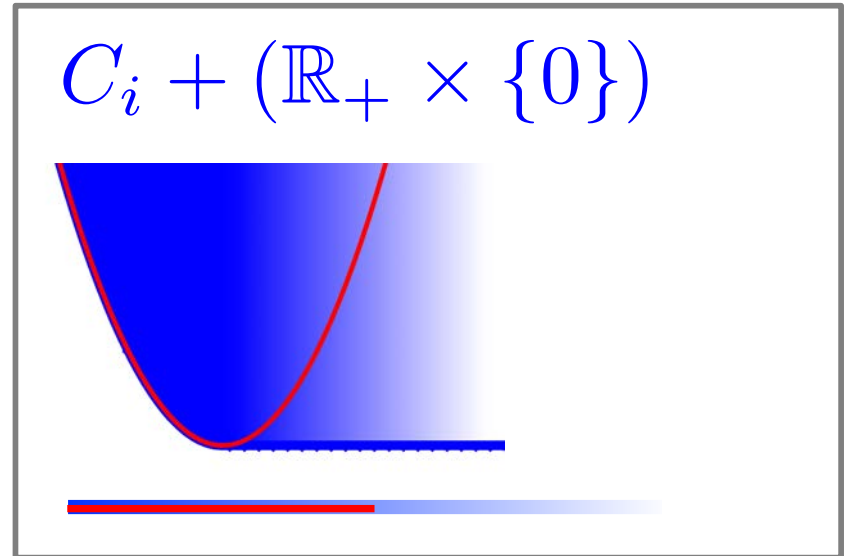
Algebraic Representation Issues



$$x_1^2 \leq x_2 \leq 1$$

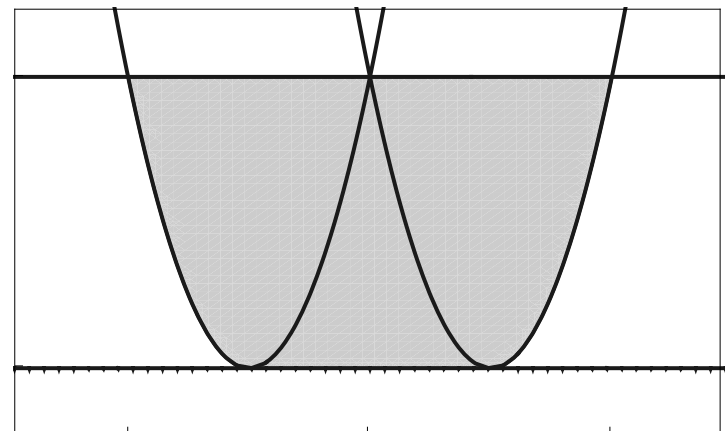


$$[-1, 1] \times 0$$



$$C_1 + (\mathbb{R}_+ \times \{0\}) : (\max\{x_1, 0\})^2 \leq x_2 \leq 1$$

- Non-basic semi-algebraic set contained in formulation.
- Finite polynomial inequalities requires max or auxiliary vars.



Summary

- General mixed integer convex representability (MICPR):
 - Infinite union of convex sets with special structure
 - More results/questions on regularity (arXiv:1706.05135)
- Bounded MICPR = Finite unions of Convex Sets
 - Variable copies = strong (ideal), but slow computation
 - Copies can be removed, but possibly at a price
 - More on the paper (arXiv:1704.03954):
 - MIP-solver compatible formulations = gauge calculus.
 - More examples: generalizations and size reductions
 - Conditions for piecewise formulations to be ideal