## Extended Formulations for Quadratic Mixed Integer Programming

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Massachusetts Institute of Technology
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## Introduction

## Nonlinear MIP B\&B Algorithms

$$
\max \quad \sum_{i=1}^{n} c_{i} x_{i}
$$

s.t.

$$
\begin{aligned}
g_{i}(x) & \leq 0, i \in I \\
x & \in \mathbb{Z}^{n_{1}} \times \mathbb{R}^{n_{2}}
\end{aligned}
$$


$\uparrow$

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## Nonlinear MIP B\&B Algorithms

- NLP (QCP) Based B\&B

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- NLP (QCP) Based B\&B
- (Dynamic) LP Based B\&B
- Few cuts = high speed.
- Possible slow convergence.

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\left.\begin{array}{l}
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x
\end{array}\right)=\mathbb{Z}^{n_{1}} \times \mathbb{R}^{n_{2}} \quad .
$$



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$$

$$
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$$

$$
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$$

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- NLP (QCP) Based B\&B
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- Lifted LP B\&B
- Fixed extended relaxation.
- Mimic NLP B\&B.

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$\max \sum_{i=1}^{n} c_{i} x_{i}$

$$
\mathscr{C} \subset \mathbb{T}^{n_{1}} \times \mathbb{R}^{n_{2}}
$$



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$$
\text { s.t. } \quad A x+D z \leq b,
$$



$$
x \in \mathbb{Z}^{n_{1}} \times \mathbb{R}^{n_{2}}
$$



## Problem 1: Classical

$\max _{x, y} \quad \bar{a} y$
s.t.

$$
\begin{aligned}
\left\|Q^{1 / 2} y\right\|_{2} & \leq \sigma \\
\sum_{j=1}^{n} y_{j} & =1
\end{aligned}
$$

$$
y_{j} \leq x_{j} \quad \forall j \in\{1, \ldots, n\}
$$

$$
\sum_{j=1}^{n} x_{j} \leq 10
$$

$$
x \in\{0,1\}^{n}
$$

$$
y \in \mathbb{R}_{+}^{n}
$$

- $y$ fraction of the portfolio invested in each of $n$ assets.
- $\bar{a}$ expected returns of assets.
- $Q^{1 / 2}$ positive semidefinite square root of the covariance matrix $Q$ of returns.
- Hold at most 10 assets.


## Avg. of Solve Times [s] for $n \in\{20,30\}$ (CPLEX v11)



## Introduction

## Extended Formulation for Lifted LP

- Approximation of Second

Order Cone by Ben-Tal and Nemirovski (Glineur).

- $O(d \log (1 / \varepsilon))$ variables and constraints for quality $\varepsilon$.
o Problem:
o Fixed a-priori quality: no dynamic improvement.

oe.g. $\varepsilon=0.01$ for portfolio had to be calibrated.


## Dynamic Lifted Approximations

## Towards a Dynamic Lifted LP

- Separable approach by Tawarmalani and Sahinidis '05 and Hijazi et al. '14
$B^{n}:=\left\{x \in \mathbb{R}^{n}: \sum_{i=1}^{n}\left(x_{i}-\frac{1}{2}\right)^{2} \leq \frac{n-1}{4}\right\}$
Showing $B^{n} \cap \mathbb{Z}^{n}=\emptyset$ requires $2^{n}$ cuts. ${ }_{0.8}$



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## Towards a Dynamic Lifted LP

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$$
B^{n}:=\left\{x \in \mathbb{R}^{n}: \sum_{i=1}^{n}\left(x_{i}-\frac{1}{2}\right)^{2} \leq \frac{n-1}{4}\right\}
$$

Extended formulation of $B^{n}$ :

$$
\begin{aligned}
\left(x_{i}-\frac{1}{2}\right)^{2} & \leq z_{i} \quad \forall i \in[n] \\
\sum_{i=1}^{n} z_{i} & \leq \frac{n-1}{4}
\end{aligned}
$$



## Towards a Dynamic Lifted LP

- Separable approach by Tawarmalani and Sahinidis '05 and Hijazi et al. '14

$$
B^{n}:=\left\{x \in \mathbb{R}^{n}: \sum_{i=1}^{n}\left(x_{i}-\frac{1}{2}\right)^{2} \leq \frac{n-1}{4}\right\}_{0.4}^{z_{i}}
$$

Extended formulation of $B^{n}$ :

$$
\left(x_{i}-\frac{1}{2}\right)^{2} \leq z_{i} \quad \forall i \in[n]
$$

$$
\sum_{i=1}^{n} z_{i} \leq \frac{n-1}{4}
$$

$B^{n} \cap \mathbb{Z}^{n}=\emptyset$ with only $2 n$ cuts on extended formulation.

## Towards a Dynamic Lifted LP

- Separable approach by Tawarmalani and Sahinidis '05 and Hijazi et al. '14

Separable approach works for any set of the form:

$$
C=\left\{x \in \mathbb{R}: \sum_{i=1}^{n} f_{i}\left(x_{i}\right) \leq 1\right\}
$$

or

$$
C=\left\{(x, t) \in \mathbb{R} \times \mathbb{R}: \sum_{i=1}^{n} f_{i}\left(x_{i}\right) \leq t\right\}
$$

for convex $f_{i}: \mathbb{R} \rightarrow \mathbb{R}$

## Problem 1: Classical

$$
\begin{aligned}
& \max _{x, y} \\
& \text { s.t. }
\end{aligned}
$$

$$
\begin{aligned}
\left\|Q^{1 / 2} y\right\|_{2} & \leq \sigma \\
\sum_{j=1}^{n} y_{j} & =1
\end{aligned}
$$

$$
y_{j} \leq x_{j}
$$

$$
\sum_{j=1}^{n} x_{j} \leq K
$$

$$
x \in\{0,1\}^{n}
$$

$$
y \in \mathbb{R}_{+}^{n}
$$

- $y$ fraction of the portfolio invested in each of $n$ assets.
- $\bar{a}$ expected returns of assets.
- $Q^{1 / 2}$ positive semidefinite square root of the covariance matrix $Q$ of returns.
- $K$ maximum number of assets to hold.


## Problem 2 : Shortfall

$\max _{x, y}$
S.t.

$$
\begin{aligned}
\left\|Q^{1 / 2} y\right\|_{2} & \leq \sigma \\
\sum_{j=1}^{n} y_{j} & =1 \\
y_{j} & \leq x_{j} \\
\sum_{j=1}^{n} x_{j} & \leq K
\end{aligned}
$$

$$
x \in\{0,1\}^{n}
$$

$$
y \in \mathbb{R}_{+}^{n}
$$

- $y$ fraction of the portfolio invested in each of $n$ assets.
- $\bar{a}$ expected returns of assets.
- $Q^{1 / 2}$ positive semidefinite square root of the covariance matrix $Q$ of returns.
- $K$ maximum number of assets to hold.


## Problem 2 : Shortfall

$\max _{x, y}$
s.t.

$$
\begin{aligned}
\left\|Q^{1 / 2} y\right\|_{2} & \leq \frac{\bar{a} y-W_{i}^{l o w}}{\Phi^{-1}\left(\eta_{i}\right)} \quad i \in\{1,2\} \\
\sum_{j=1}^{n} y_{j} & =1 \\
y_{j} & \leq x_{j} \quad \forall j \in\{1, \ldots, n\} \\
\sum_{j=1}^{n} x_{j} & \leq K \\
x & \in\{0,1\}^{n} \\
y & \in \mathbb{R}_{+}^{n}
\end{aligned}
$$

- $y$ fraction of the portfolio invested in each of $n$ assets.
- $\bar{a}$ expected returns of assets.
- $Q^{1 / 2}$ positive semidefinite square root of the covariance matrix $Q$ of returns.
- $K$ maximum number of assets to hold.
- Approximation of $\operatorname{Prob}\left(\bar{a} y \geq W_{i}^{l o w}\right) \geq \eta_{i}$


## Extended Formulation for SOCP

$$
L^{n}=\{(x, t) \in \mathbb{R} \times \mathbb{R}:\|x\| \leq t\}
$$

Extended formulation of $L^{n}=$
homogenization of $B^{n}$ formulation:

$$
\begin{aligned}
x_{i}^{2} & \leq z_{i} \cdot t \quad \forall i \in[n] \\
\sum_{i=1}^{n} z_{i} & \leq t
\end{aligned}
$$

## Extended Formulation for SOCP

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L^{n}=\{(x, t) \in \mathbb{R} \times \mathbb{R}:\|x\| \leq t\}
$$

Extended formulation of $L^{n}=$
homogenization of $B^{n}$ formulation:

$$
\sum_{i=1}^{n} z_{i} \leq t \quad \begin{aligned}
& x_{i}^{2} \leq z_{i} \cdot t
\end{aligned} \quad \forall i \in[n]
$$

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O LP B\&B on extended form.

$$
\max \quad \sum_{i=1}^{n} c_{i} x_{i}
$$

s.t.

$$
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$$
\begin{aligned}
& \max \sum_{i=1}^{n} c_{i} x_{i} \\
& \text { s.t. } g_{i}(x, z) \leq 0, i \in I, \\
& \qquad x \in \mathbb{Z}^{n_{1}} \times \mathbb{R}^{n_{2}}
\end{aligned}
$$



## Computational Results

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- Averages over 20 instances:
- Classical and Shortfall. 40, 50 and 60 stocks.
- Solvers:
o CPLEX/Gurobi QCP-BB on original formulation .
- Lifted LP: Implemented in JuMP using CPLEX's branch, incumbent and heuristic callback.
o CPLEX/Gurobi LP-BB on extended "separable" reformulation.


## LiftedLP v/s QCP: Classical


$\square$ LiftedLP
$\square$ CPLEX QCP
$\square$ Gurobi QCP

## LiftedLP v/s QCP: Shortfall


$\square$ LiftedLP
$\square$ CPLEX QCP
$\square$ Gurobi QCP

## LiftedLP v/s Dynamic : Classical


$\square$ LiftedLP

- CPLEX SEP-LP
$\square$ Gurobi SEP-LP


## LiftedLP v/s Dynamic : Shortfall


$\square$ LiftedLP
C CPLEX SEP-LP
$\square$ Gurobi SEP-LP

## Summary

- Lifted LP: 200 lines of JuMP code in a weekend.
- Developed by ORC students lain Dunning, Joey Huchette and Miles Lubin

O https://github.com/JuliaOpt/JuMP.jl

- Poster at MIP 2014. OSU, July 21st
- Talk at INFORMS. San Francisco, November
- Dynamic Lifted LP:
oComparable performance with simple reformulation.

