

Embedding Formulations, Complexity and Representability for Unions of Convex Sets

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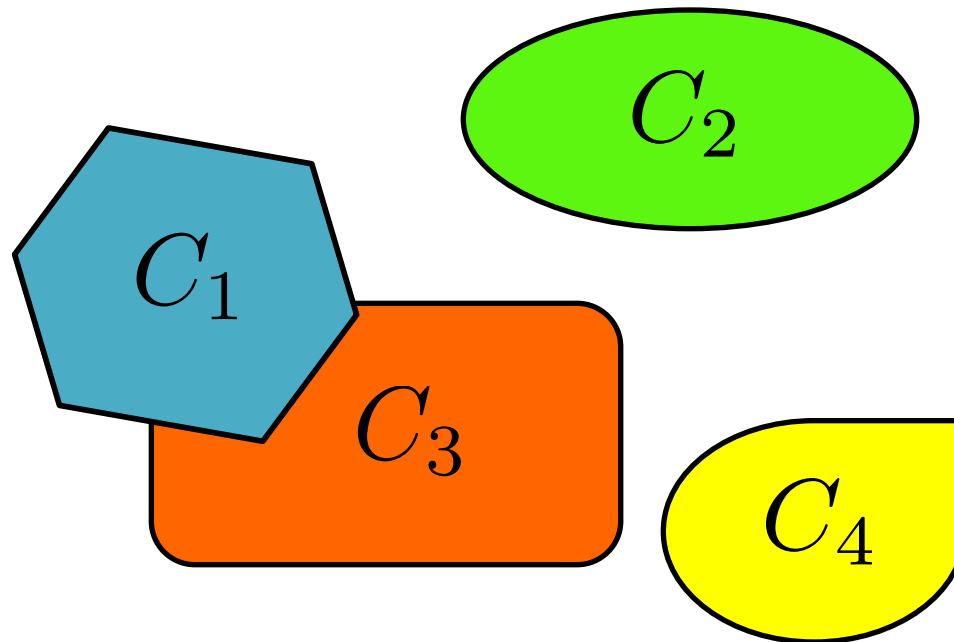
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Nonlinear Mixed 0-1 Integer Formulations

- Modeling Finite Alternatives = Unions of Convex Sets

$$x \in \bigcup_{i=1}^n C_i \subseteq \mathbb{R}^d$$



Extended and Non-Extended Formulations for $\bigcup_{i=1}^n C_i$

$$C_i = \{x \in \mathbb{R}^d : A^i x \leq b^i\}$$

Extended

Non-Extended

$$\begin{aligned} A^i x^i &\leq b^i y_i && \forall i \in [n] \\ \sum_{i=1}^n x^i &= x \\ \sum_{i=1}^n y_i &= 1 \\ y &\in \{0, 1\}^n \\ x, x^i &\in \mathbb{R}^d && \forall i \in [n] \end{aligned}$$

$$\begin{aligned} A^i x - b^i &\leq M_i (1 - y_i) && \forall i \in [n] \\ \sum_{i=1}^n y_i &= 1 \\ y &\in \{0, 1\}^n \\ x &\in \mathbb{R}^d && \forall i \in [n] \end{aligned}$$

Large, but strong (ideal*)

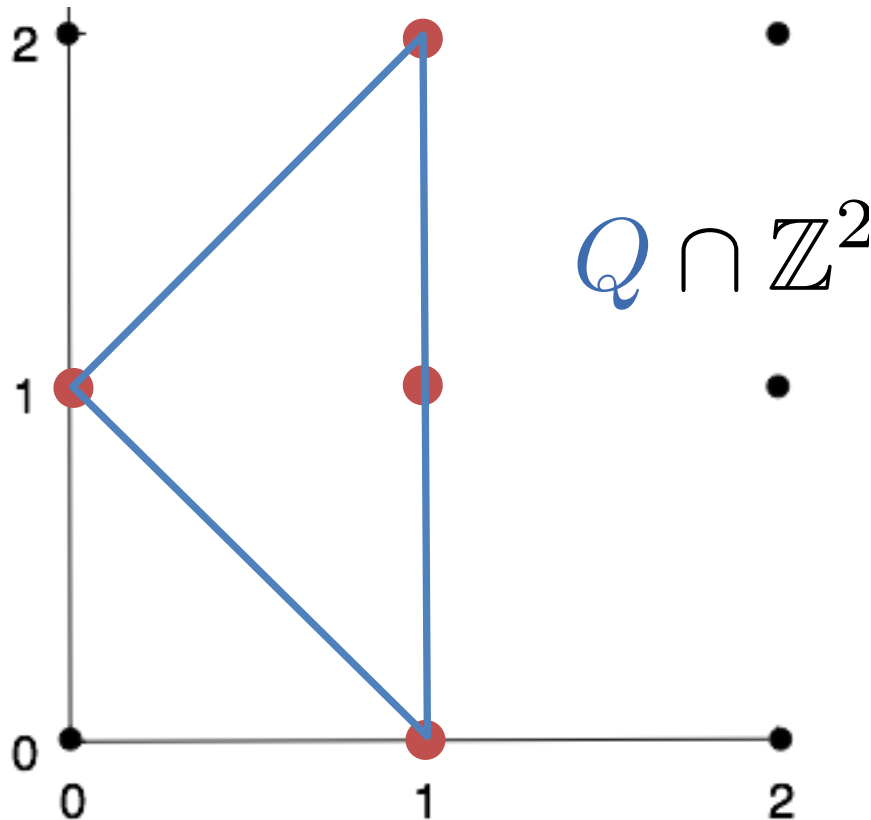
Small, but weak?

*Integral y in extreme points of LP relaxation

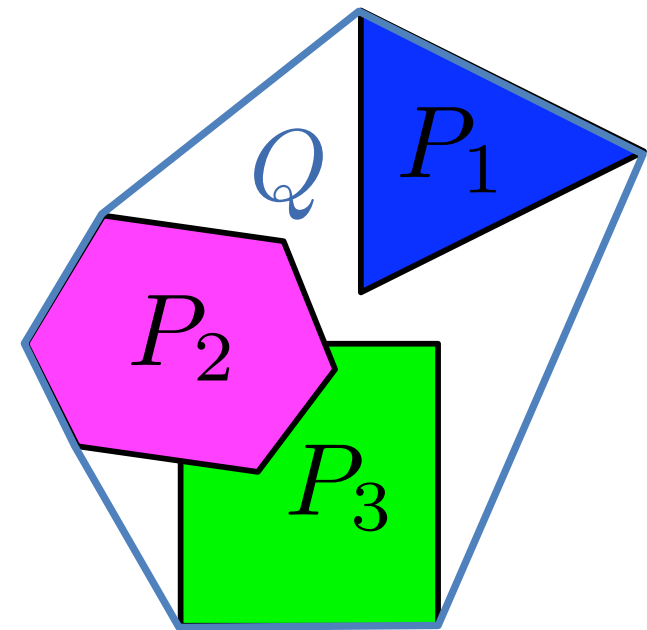
Constructing Non-extended Ideal Formulations

- Pure Integer:

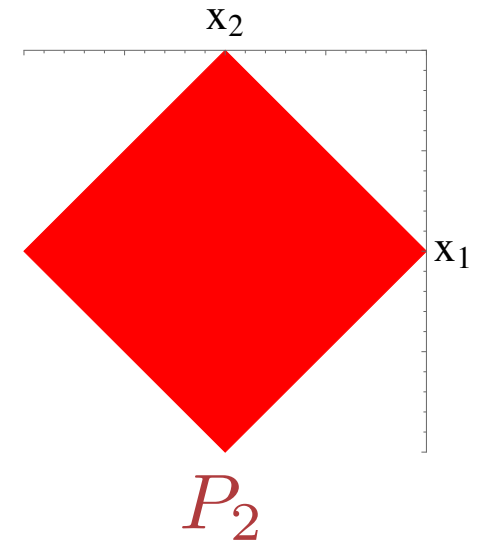
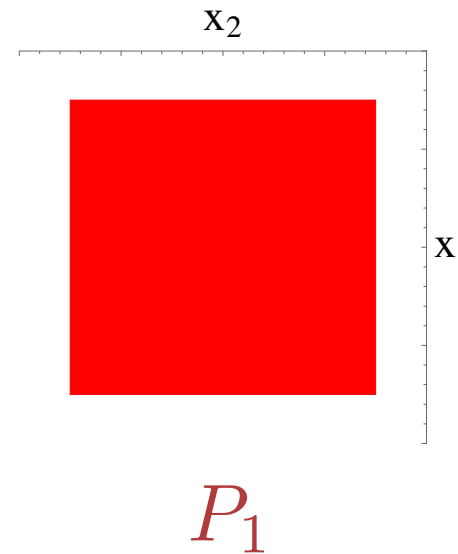
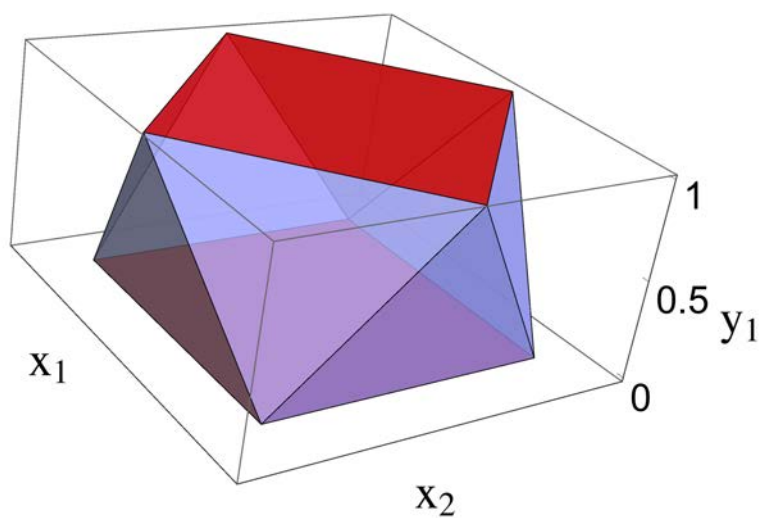
$$Q := \text{conv} \left(\{p^i\}_{i=1}^n \right)$$



- Mixed Integer:



Embedding Formulation = Ideal non-Extended



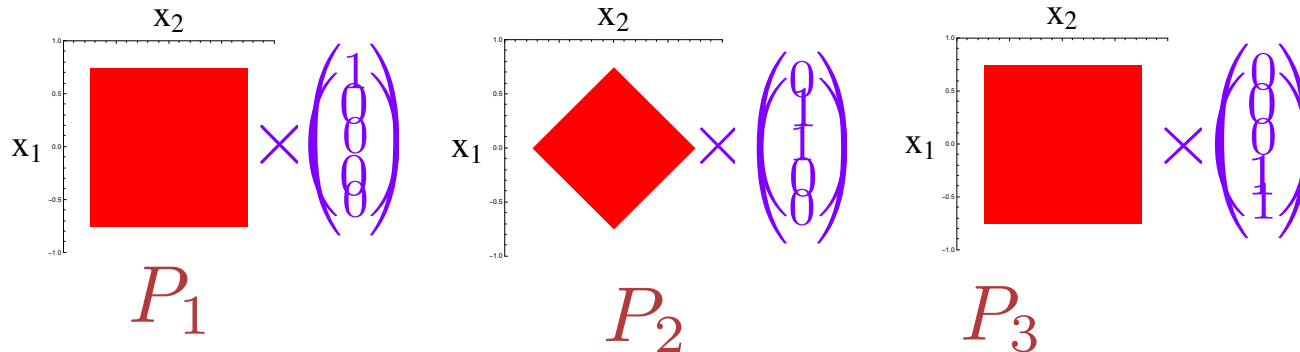
$$Q(H) := \text{conv} \left(\bigcup_{i=1}^n P_i \times \{h^i\} \right)$$

$$(x, y) \in Q \cap (\mathbb{R}^d \times \mathbb{Z}^k) \iff y = h^i \wedge x \in P_i$$

$$\text{ext}(Q) \subseteq \mathbb{R}^d \times \mathbb{Z}^k \quad H := \{h^i\}_{i=1}^n \subseteq \{0, 1\}^k, \quad h^i \neq h^j$$

Alternative Encodings

- 0-1 encodings guarantee validity



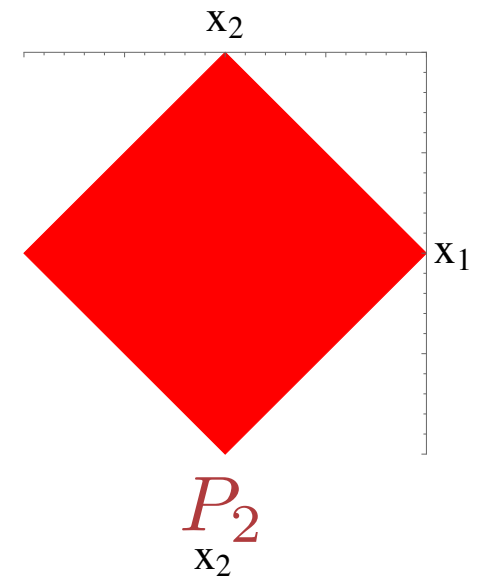
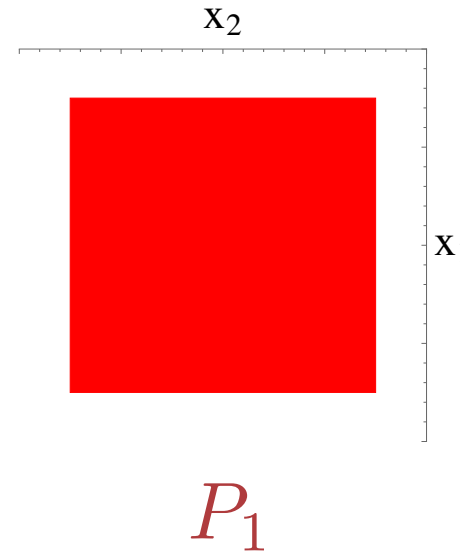
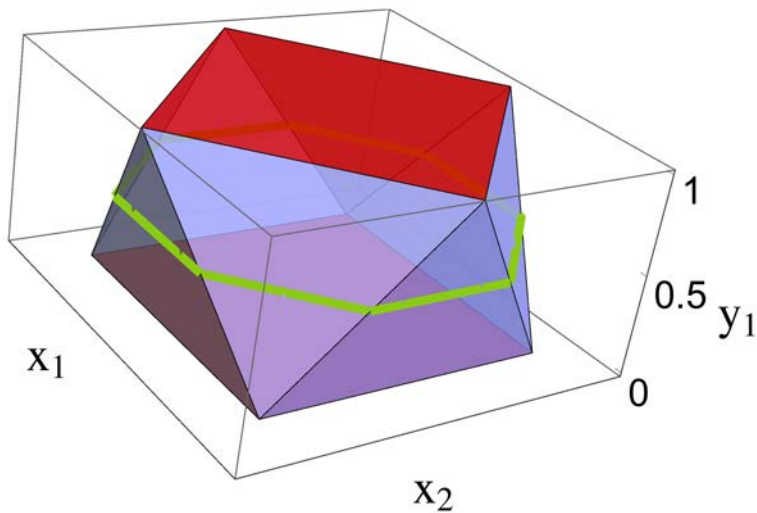
- Options for 0-1 encodings:
 - Traditional or **Unary** encoding

$$H = \left\{ y \in \{0, 1\}^n : \sum_{i=1}^n y_i = 1 \right\} \quad e_j^i = \begin{cases} 1 & i = j \\ 0 & i \neq j \end{cases}$$

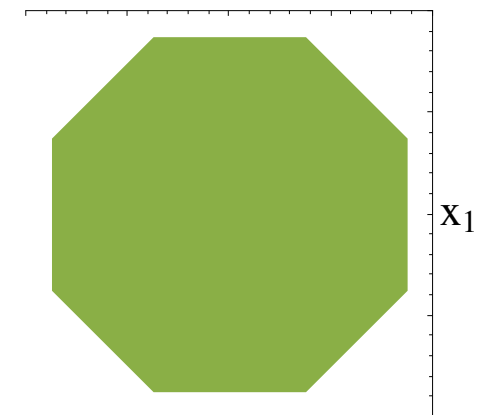
$$= \{e^i\}_{i=1}^n$$

- **Binary** encodings: $H \equiv \{0, 1\}^{\log_2 n}$
- Others (e.g. **incremental** encoding \equiv unary)

Unary Encoding, Minkowski Sum and Cayley Trick



$$Q \cap (\mathbb{R}^2 \times \{0.5\}) \equiv P_1 + P_2 =$$



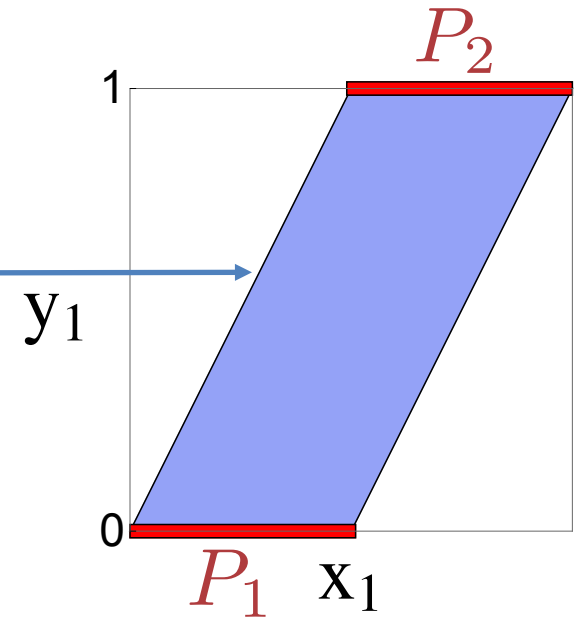
For traditional or unary encoding:

$$Q(H) \cap \left(\mathbb{R}^d \times \left\{ \frac{1}{n} \sum_{i=1}^n \mathbf{e}^i \right\} \right) \equiv \sum_{i=1}^n P_i$$

Complexity of Family of Polyhedra $\mathcal{P} := \{P_i\}_{i=1}^n$

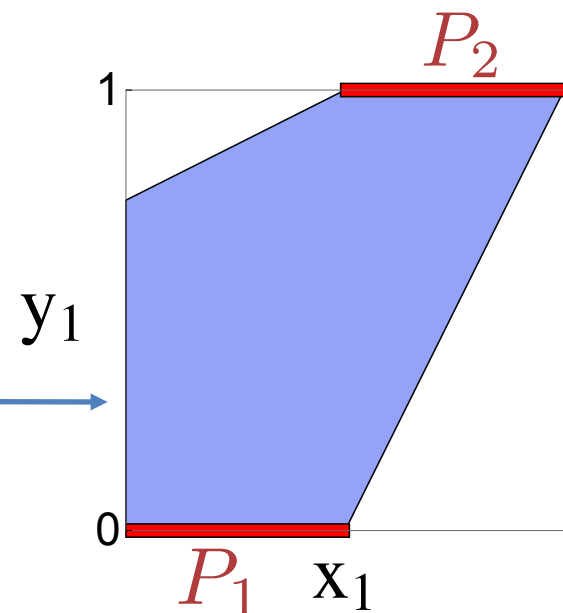
- Embedding complexity = smallest **ideal** formulation

$$mc(\mathcal{P}) := \min_H \{\text{size}(Q(H))\}$$



- Relaxation complexity = smallest formulation

$$rc(\mathcal{P}) := \min_{Q,H} \{\text{size}(Q)\}$$



$\text{size}(Q) := \#$ of facets of Q

Complexity Results

- Lower and Upper bounds for special structures:
 - e.g. for Special Order Sets of Type 2 (SOS2) on n variables

- Embedding complexity (ideal)

$$2^{\lceil \log_2 n \rceil} \longleftarrow \text{General Inequalities}$$

$$n + 1 \leq \dots \leq n + 1 + 2^{\lceil \log_2 n \rceil} \longleftarrow \text{Total}$$

- Relaxation complexity (non-ideal)

$$2 \leq \dots \leq 4 \longleftarrow \text{General Inequalities}$$

$$2 \leq \dots \leq 5 + 2n \longleftarrow \text{Total}$$

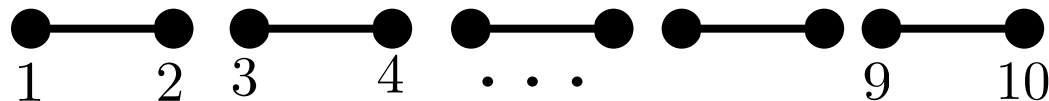
- Relation to other complexity measures

$$\text{hc}(\mathcal{P}) := \text{size} \left(\text{conv} \left(\bigcup_{i=1}^n P_i \right) \right)$$

$$\text{xc}(\mathcal{P}) := \min_R \left\{ \text{size}(R) : \text{proj}_x(R) = \text{conv} \left(\bigcup_{i=1}^n P_i \right) \right\}$$

- Still open questions (see V. 2015)

Example of Constant Sized Non-Ideal Formulation



$$x \in \bigcup_{i=1}^n P_i$$

$$P_i \subseteq \Delta^{2n} := \left\{ x \in \mathbb{R}_+^{2n} : \sum_{i=1}^{2n} x_i = 1 \right\}$$

$$P_i = \left\{ x \in \Delta^{2n} : x_j = 0 \quad \forall j \notin \{2i-1, 2i\} \right\}$$

Faces for Ideal Formulation with Unary Encoding

- Two types of facets (or faces):

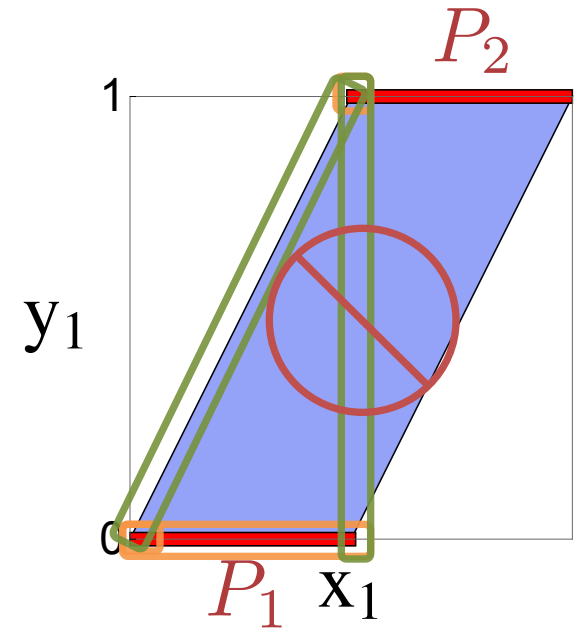
- $P_1 \times \{0\} \equiv y_i \geq 0$

- $\text{conv}((F_1 \times 0) \cup (F_2 \times 1))$

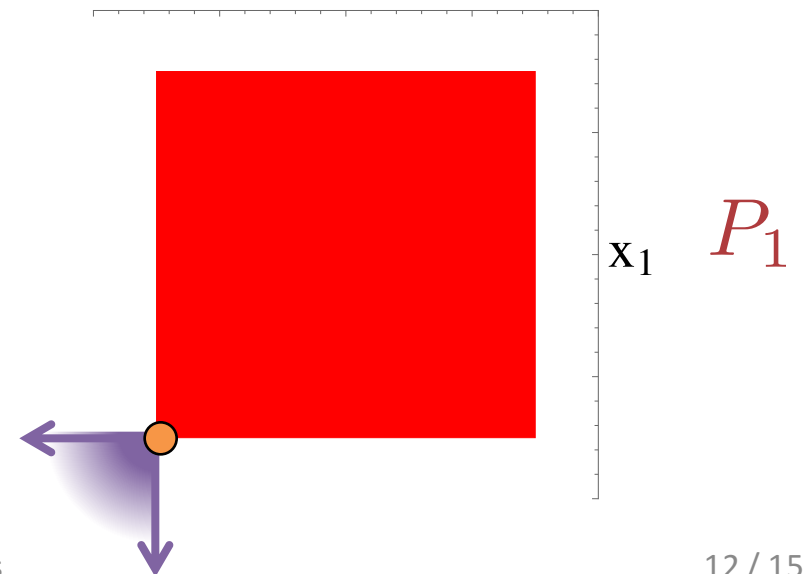
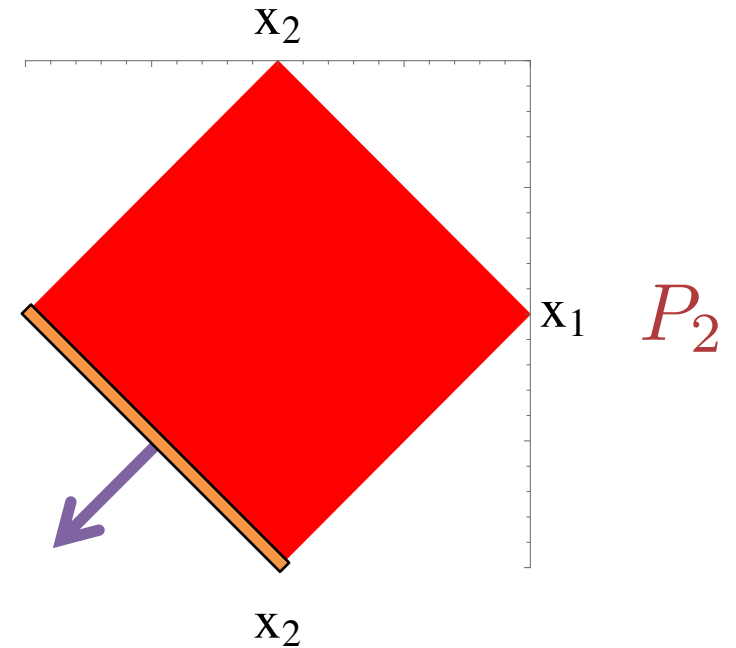
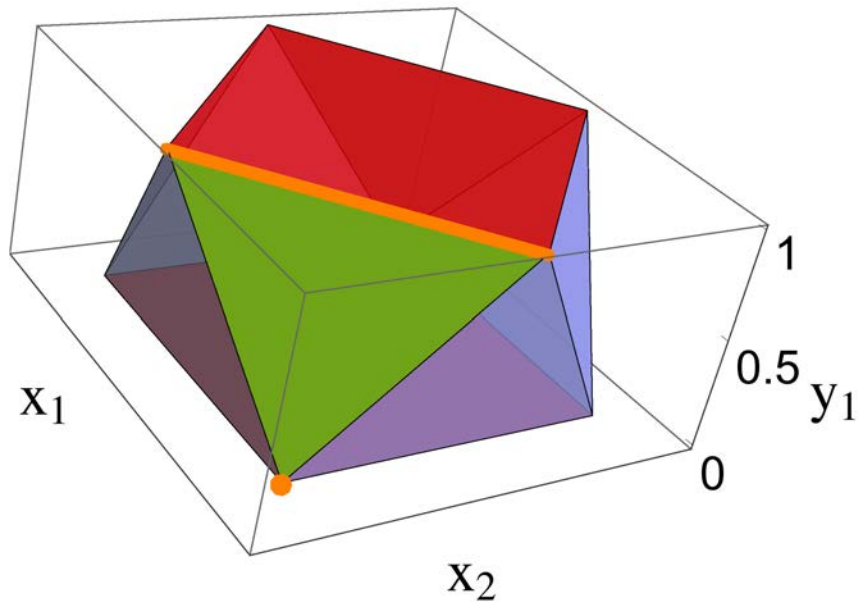
F_i proper face of P_i

- Not all combinations of faces

- Which ones are valid?



Valid Combinations = Common Normals

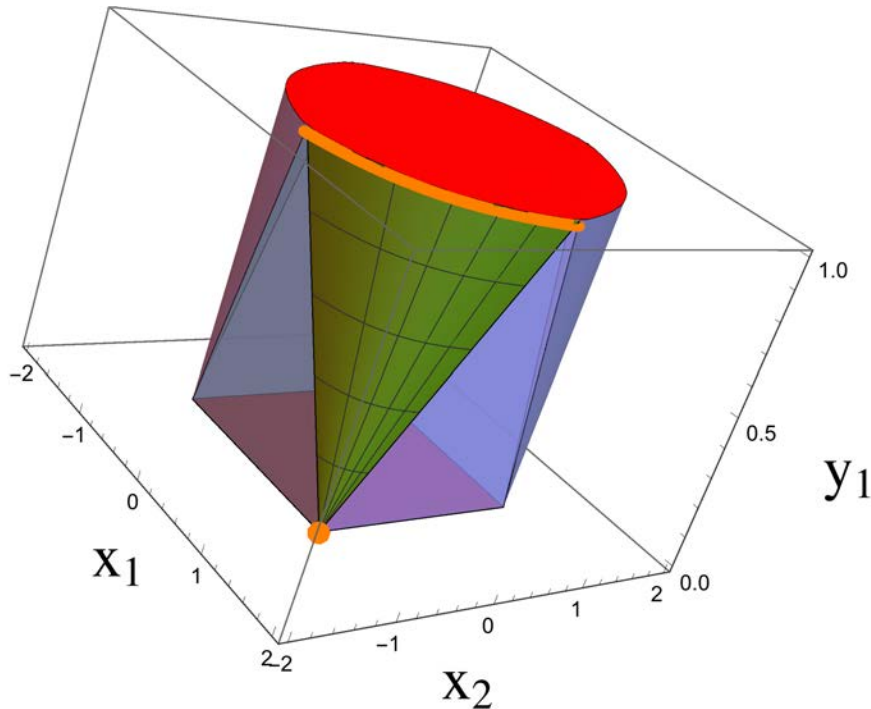


$$N(F_1) \cap N(F_2) \neq \emptyset$$

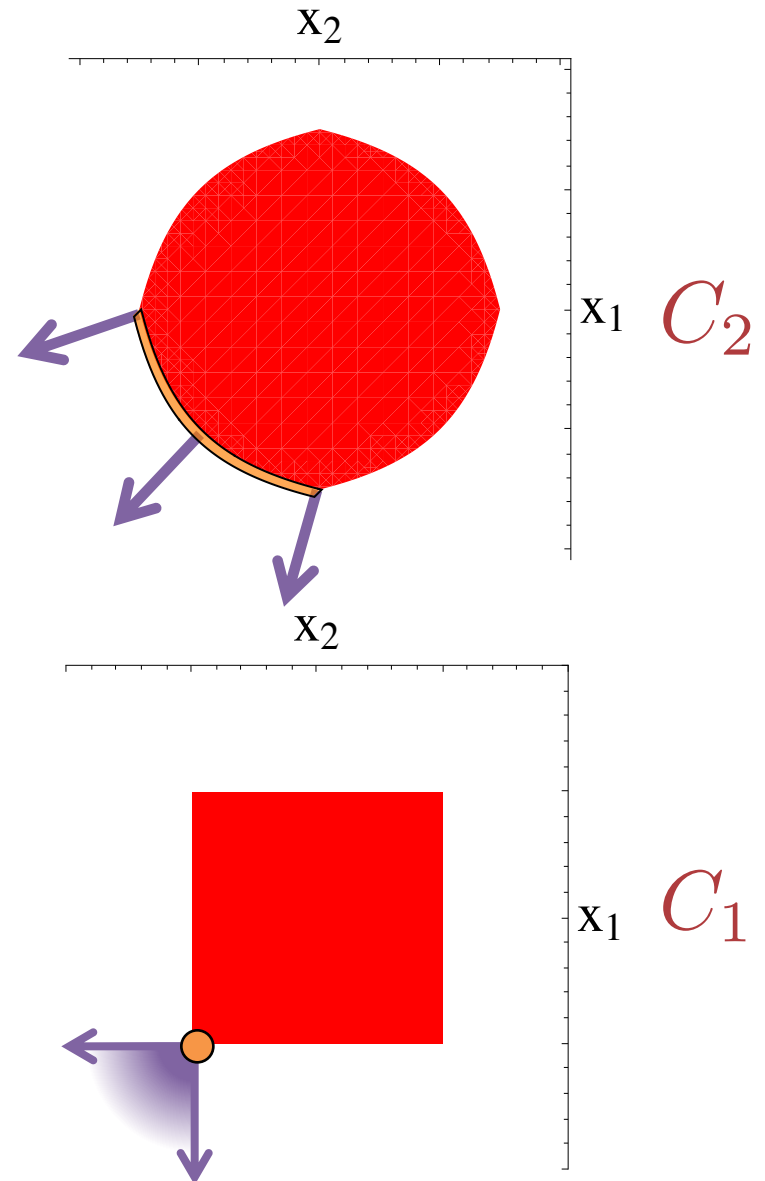


$\text{conv}((F_1 \times 0) \cup (F_2 \times 1))$
is face of $Q(H)$

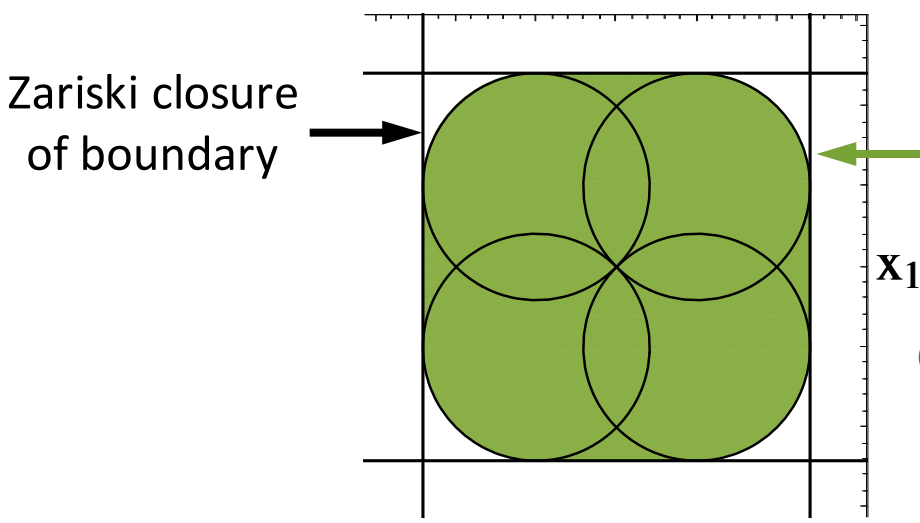
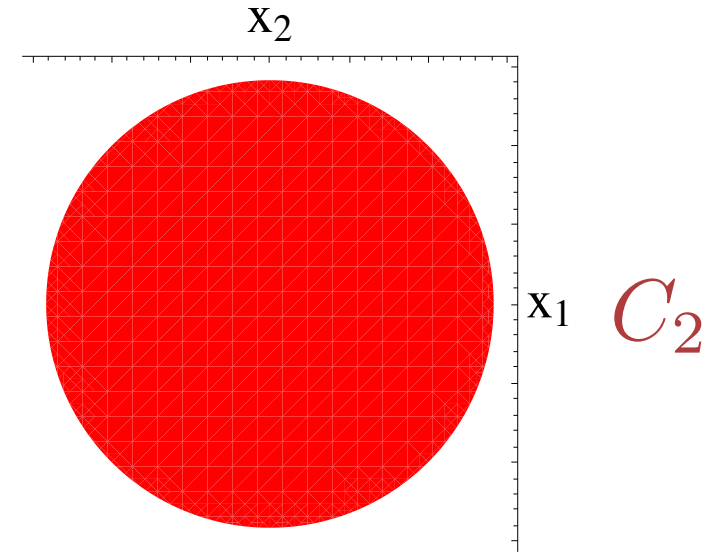
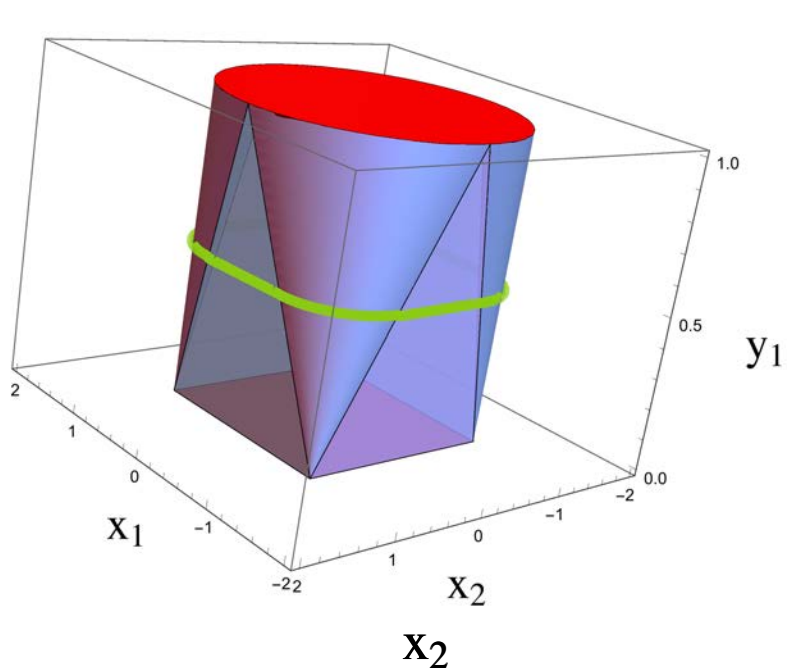
Unary Embedding for Unions of Convex Sets



- Description of boundary of $Q(H)$ is easy if “normals condition” yields convex hull of **1** nonlinear constraint and point(s)



Bad Example: Representability Issues



Description with finite number of (quadratic) polynomial inequalities?

$$Q := \text{conv} ((C_1 \times \{0\}) \cup (C_2 \times \{1\}))$$

can fail to be basic semi-algebraic

Summary

- Embedding Formulations = Systematic procedure for strong (ideal) non-extended formulations
 - Encoding can significantly affect size
- Complexity of Union of Polyhedra beyond convex hull
 - Embedding Complexity (non-extended ideal formulation)
 - Relaxation Complexity (any non-extended formulation)
 - Still open questions on relations between complexity
(Embedding Formulations and Complexity for Unions of Polyhedra, arXiv:1506.01417)
- Embedding Formulations for Convex Sets
 - MINLP formulations
 - Can have representability issues
 - Open question: minimum number of auxiliary variables for fixing this