Discrete Geometry for Small and Strong Mixed Integer Programming Formulations

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Nonlinear Mixed <u>0-1</u> Integer Formulations

• Modeling Finite Alternatives = Unions of Convex Sets

$$x \in \bigcup_{i=1}^{n} C_i \subseteq \mathbb{R}^d$$

$$\overbrace{C_1}_{C_3}$$

$$\overbrace{C_4}$$

Extended and Non-Extended Formulations for $\bigcup_{i=1}^{n} C_i$

$$C_{i} = \left\{ x \in \mathbb{R}^{d} : A^{i}x \leq b^{i} \right\}$$
Extended
$$A^{i}x^{i} \leq b^{i}y_{i} \quad \forall i \in [n]$$

$$\sum_{i=1}^{n} x^{i} = x$$

$$\sum_{i=1}^{n} y_{i} = 1$$

$$y \in \{0, 1\}^{n}$$

$$x, x^{i} \in \mathbb{R}^{d} \quad \forall i \in [n]$$
Non-Extended
$$A^{i}x - b^{i} \leq M_{i}(1 - y_{i}) \quad \forall i \in [n]$$

$$A^{i}x - b^{i} \leq M_{i}(1 - y_{i}) \quad \forall i \in [n]$$

Small? and strong (ideal^{*})

Small, but weak?

*Integral y in extreme points of LP relaxation

Embedding Formulations

Outline

- Smallest ideal non-extended linear formulation
 - Generalization of Cayley Embedding of Polytopes
 - Hyperplane arrangements
- Small ideal non-extended nonlinear formulations
 - Boundary structure of Cayley Embedding: from polytopes to closed convex sets
- "Formulations" with a general integer variables (possibly a fixed number)

Ideal Non-Extended Formulations and Hyperplane Arrangements

Embedding Formulation = Ideal non-Extended



Embedding Formulation = Ideal non-Extended



Special Ordered Sets = Simplex Faces = $\mathcal{P} := \{P_i\}_{i=1}^n$

•
$$\Delta^{d+1} := \left\{ x \in \mathbb{R}^{d+1}_{+} : \sum_{i=1}^{d+1} x_i = 1 \right\} = \operatorname{conv} \left(\left\{ e^i \right\}_{i=1}^{d+1} \right)$$

 $P_i := \operatorname{conv} \left(\left\{ e^j \right\}_{j \in T_i} \right) = \left\{ x \in \Delta^{d+1} : \sum_{j \notin T_i} x_i \le 0 \right\}$
 $T_i \subseteq \{1, \dots, d+1\}$
• $\operatorname{mc} \left(\mathcal{P} \right) := \operatorname{min}_H \left\{ \operatorname{size} \left(Q \left(H \right) \right) \right\},$

size (Q(H)) := # facets

 $\operatorname{mc}_{G}(\mathcal{P}) := \operatorname{min}_{H} \left\{ \operatorname{size}_{G} \left(Q\left(H \right) \right) \right\},$ $\operatorname{size}_{G} \left(Q\left(H \right) \right) := \# \text{ non-bound facets}$

Embedding Formulations

Special Ordered Sets of Type 2 (SOS2) = $\mathcal{P} := \{P_i\}_{i=1}^n$

•
$$P_i := \operatorname{conv}(\{e^i, e^{i+1}\}) \subseteq \Delta^{n+1}, \quad i \in [n]$$



Special Ordered Sets of Type 2 (SOS2) = $\mathcal{P} := \{P_i\}_{i=1}^n$

•
$$P_i := \operatorname{conv}\left(\left\{e^i, e^{i+1}\right\}\right) \subseteq \Delta^{n+1}, \quad i \in [n]$$

Claim: $\operatorname{mc}_G(\mathcal{P}) = 2 \left[\log_2 n\right],$
 $n+1 \leq \operatorname{mc}(\mathcal{P}) \leq n+1+2 \left[\log_2 n\right]$

•
$$(x, y) \in Q(H) = \operatorname{conv}\left(\bigcup_{i=1}^{n} P_i \times \{h^i\}\right)$$

= $\operatorname{conv}\left(\bigcup_{i=1}^{n} \{e^i, e^{i+1}\} \times \{h^i\}\right)$

• $a \cdot x \leq b \cdot y$ $a_{i+1} \leq \min\left\{b \cdot h^i, b \cdot h^{i+1}\right\}$

 $b \in L(H) := \operatorname{aff}(H) - h^1$

Embedding Formulations

 $b \cdot \left(h^{i+1} - h^i \right) = 0$

Embedding Formulation for SOS2: Part 1

• From encodings (H) to hyperplanes:



Embedding Formulation for SOS2: Part 1

• From encodings (H) to hyperplanes:



non-bound facets = 2 × # of hyperplanes





Embedding Complexity for SOS2

- Lower Bound: $L(H) := \operatorname{aff} (H) h^1$ $\operatorname{mc}_G (\mathcal{P}) \ge 2 \times \min \# \text{ of hyperplanes}$ $\min \# \text{ of hyperplanes} \ge \dim (L(H))$ $\dim (L(H)) \ge \lceil \log_2 n \rceil$
 - Upper Bound: $H = \{0, 1\}^{\lceil \log_2 n \rceil}$ - Gray code: $\{h^i - h^{i+1}\}_{i=1}^{n-1} \equiv \{e^i\}_{i=1}^{\lceil \log_2 n \rceil}$ size_G $(Q(H)) = 2 \lceil \log_2 n \rceil$

$$n+1 \le \operatorname{mc}\left(\mathcal{P}\right) \le n+1+2\left\lceil \log_2 n\right\rceil$$

Minkowski Sums and Nonlinear MIP Formulations

Unary Encoding, Minkowski Sum and Cayley Trick



Faces of Cayley Embedding

- Two types of facets (or faces): $-P_1 \times \{0\} \equiv y_i \geq 0$
 - $-\operatorname{conv}\left(\left(F_1\times 0\right)\cup\left(F_2\times 1\right)\right)$
 - F_i proper face of P_i
 - Not all combinations of faces
 - Which ones are valid?



Valid Combinations = Common Normals

Embedding Formulations

 P_2

X₁

 \mathbf{X}_{1}

Characterization Extends to Closed Convex Sets

Small Formulations for Isotone Sets

• Studied by Hijazi et al. '12 and Bonami et al. '15 (n=1, 2):

$$-C_i = \left\{ x \in \mathbb{R}^d : l^i \le x \le u^i, \quad f_i(x) \le 0 \right\}$$

• $f_i(x)$ component-wise monotonous (i=1,2 opposite).

• Ideal Formulation $y_1 l^1 + y_2 l^2 \le x \le y_1 u^1 + y_2 u^2$ $f_J^i(x, y) \le 0$ $\forall J \subseteq [d], i \in [2]$ $y_1 + y_2 = 1$ $y_i \in \{0, 1\}$ $i \in [2]$

Embedding Formulations

Boundary Structure = Redundancy Detection

$$y_{1}l^{1} + y_{2}l^{2} \leq x \leq y_{1}u^{1} + y_{2}u^{2}$$

$$f^{i}(x, y) \leq 0 \qquad \forall i \in [2]$$

$$y_{1} + y_{2} = 1$$

$$y_{i} \geq \{0, 1\} \qquad i \in [2]$$

Boundary Structure = Redundancy Detection

"Formulations" With a Fixed Number of General Integer Variables

Alternative Encodings

 \mathbf{X}_1

X₂

 P_1

 \mathbf{X}_1

 \mathbf{X}_2

General integer (rational?) encodings:

- Points in convex position
- Recover convex sets by (possibly non-axis aligned) sections

 \mathbf{X}_2

- Example: Integers in moment curve
- For SOS2: Hyperplane characterization still works

X1 •

- 2-dim moment curve = 2(n-1) general inequalities
- More bounds soon (with Joey Huchette)

3

2

 $P_i \times A$

i=1

Summary

- Embedding Formulations = Systematic procedure for ideal non-extended formulations
 - Encoding can significantly affect size
 - Results beyond SOS2, but many open questions
- Extension to General Convex Sets
 - Can yield practical formulations
 - Not always practical (basic semi-algebraic representability)
- Using General Integer Variables
 - Smaller formulations not likely (in general)
 - General convex MIP representability (w. M. Lubin and I. Zadik):
 - The set of prime numbers is not convex MIP representable