

# Mixed Integer Programming (MIP) for Daily Fantasy Sports, Statistics and Marketing

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# MIP & Daily Fantasy Sports



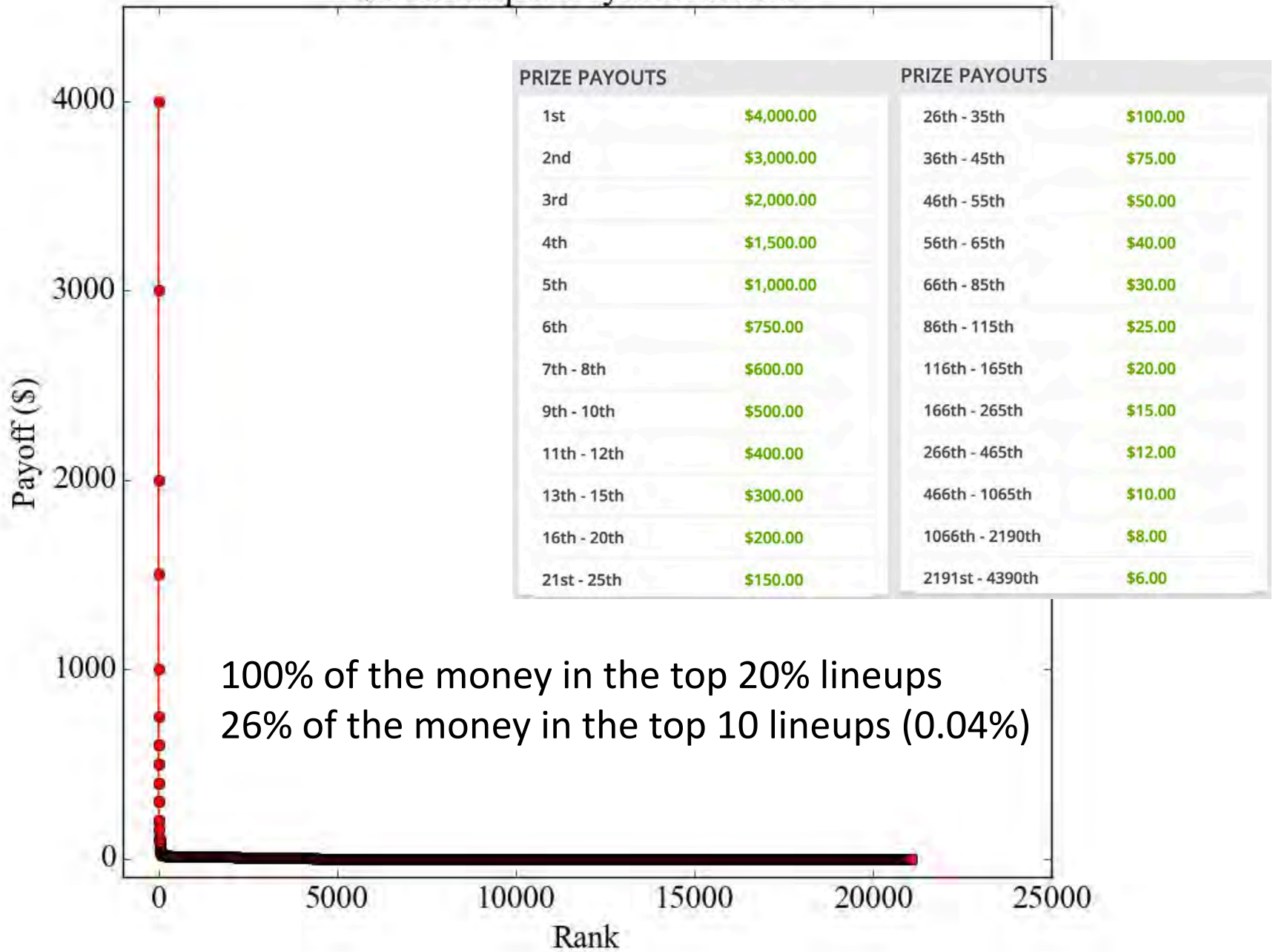
# Example Entry



Avg. Rem. / Player: \$0  
Rem. Salary: \$0

POS	PLAYER	OPP	FPPG	SALARY	
C	Jussi Jokinen	Fla@Anh	3.1	\$5,300	✘
C	Brandon Sutter	Pit@Van	3.0	\$4,400	✘
W	Nikolaj Ehlers	Wpg@Tor	3.9	\$4,800	✘
W	Daniel Sedin	Pit@Van	3.8	\$6,400	✘
W	Radim Vrbata	Pit@Van	3.4	\$5,800	✘
D	Brian Campbell	Fla@Anh	2.6	\$4,100	✘
D	Morgan Rielly	Wpg@Tor	3.5	\$4,200	✘
G	Corey Crawford	StL@Chi	6.3	\$7,800	✘
UTIL	Blake Wheeler	Wpg@Tor	4.8	\$7,200	✘

# \$55K Sniper Payoff Structure



# Building a Lineup



# Integer Programming Formulation

- We will make a bunch of lineups consisting of 9 players each
- Use an integer programming approach to find these lineups

Decision variables

$$x_{pl} = \begin{cases} 1, & \text{if player } p \text{ in lineup } l \\ 0, & \text{otherwise} \end{cases}$$

# Basic Feasibility

- 9 different players
- Salary less than \$50,000

## Basic constraints

$$\sum_{p=1}^N c_p x_{pl} \leq \$50,000, \quad (\text{budget constraint})$$

$$\sum_{p=1}^N x_{pl} = 9, \quad (\text{lineup size constraint})$$

$$x_{pl} \in \{0, 1\}, \quad 1 \leq p \leq N.$$



# Position Feasibility

- Between 2 and 3 centers
- Between 3 and 4 wingers
- Between 2 and 3 defensemen
- 1 goalie

## Position constraints

$$2 \leq \sum_{p \in C} x_{pl} \leq 3, \quad (\text{center constraint})$$

$$3 \leq \sum_{u \in W} x_{pl} \leq 4, \quad (\text{winger constraint})$$

$$2 \leq \sum_{u \in D} x_{pl} \leq 3, \quad (\text{defensemen constraint})$$

$$\sum_{u \in G} x_{pl} = 1 \quad (\text{goalie constraint})$$



# Team Feasibility

- At least 3 different NHL teams

## Team constraints

$$t_i \leq \sum_{p \in T_i} x_{pl}, \quad \forall i \in \{1, \dots, N_T\}$$

$$\sum_{i=1}^{N_T} t_i \geq 3,$$

$$t_i \in \{0, 1\}, \quad \forall i \in \{1, \dots, N_T\}.$$

# Maximize Points

- Forecasted points for player  $p$ :  $f_p$



Score type	Points
Goal	3
Assist	2
Shot on Goal	0.5
Blocked Shot	0.5
Short Handed Point Bonus (Goal/Assist)	1
Shootout Goal	0.2
Hat Trick Bonus	1.5
Win (goalie only)	3
Save (goalie only)	0.2
Goal allowed (goalie only)	-1
Shutout Bonus (goalie only)	2

Table 1 Points system for NHL contests in DraftKings.

## Points Objective Function

$$\sum_{p=1}^N f_p x_{pl}$$

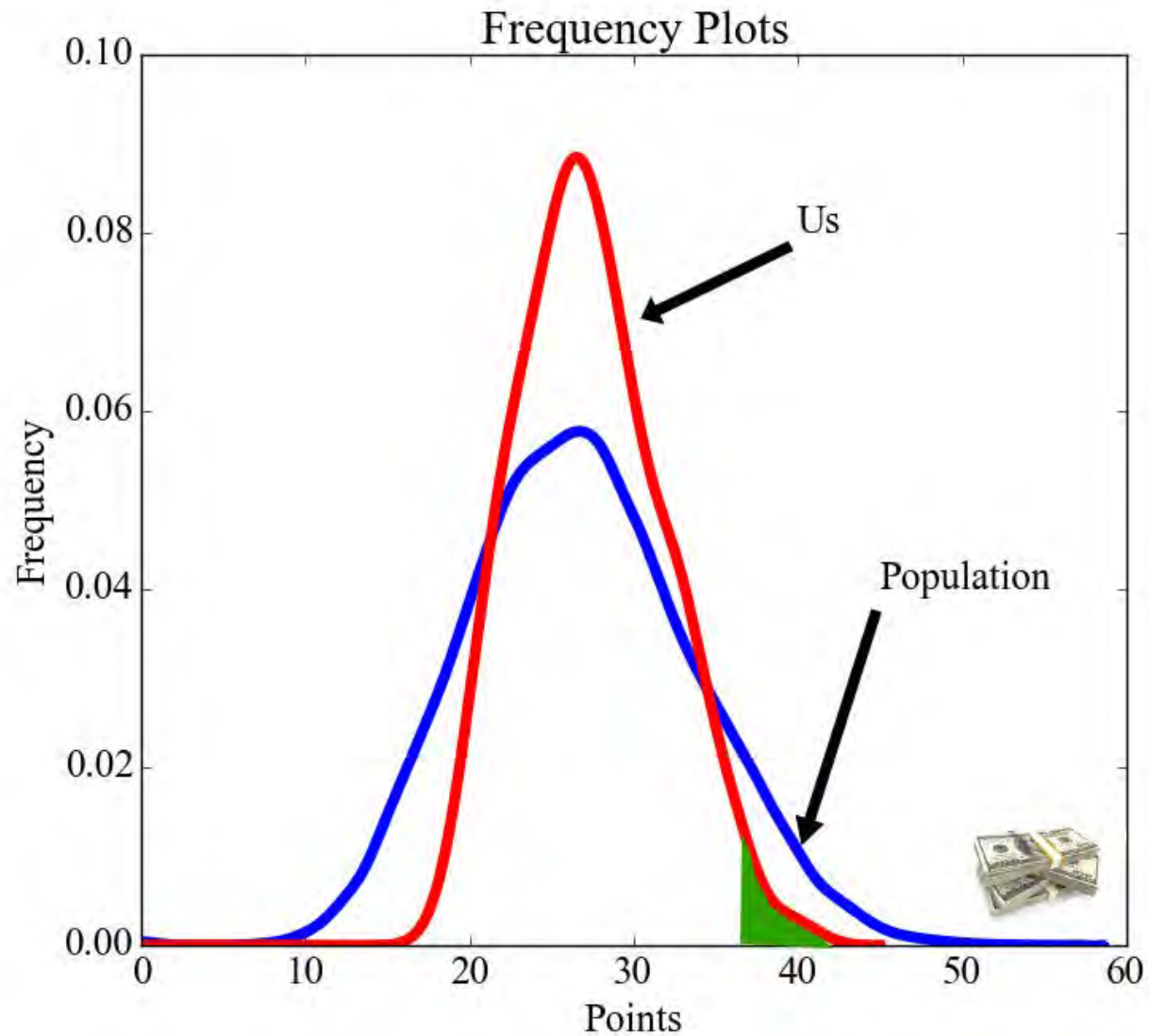
# Lineup

Projections: 5.4 2.5 3.4 3.0 3.2 4.2 3.5 3.4 5.7  
\$9500 \$2700 \$4600 \$3800 \$4600 \$6400 \$5200 \$5100 \$8000  
W UTIL D D C C W W G

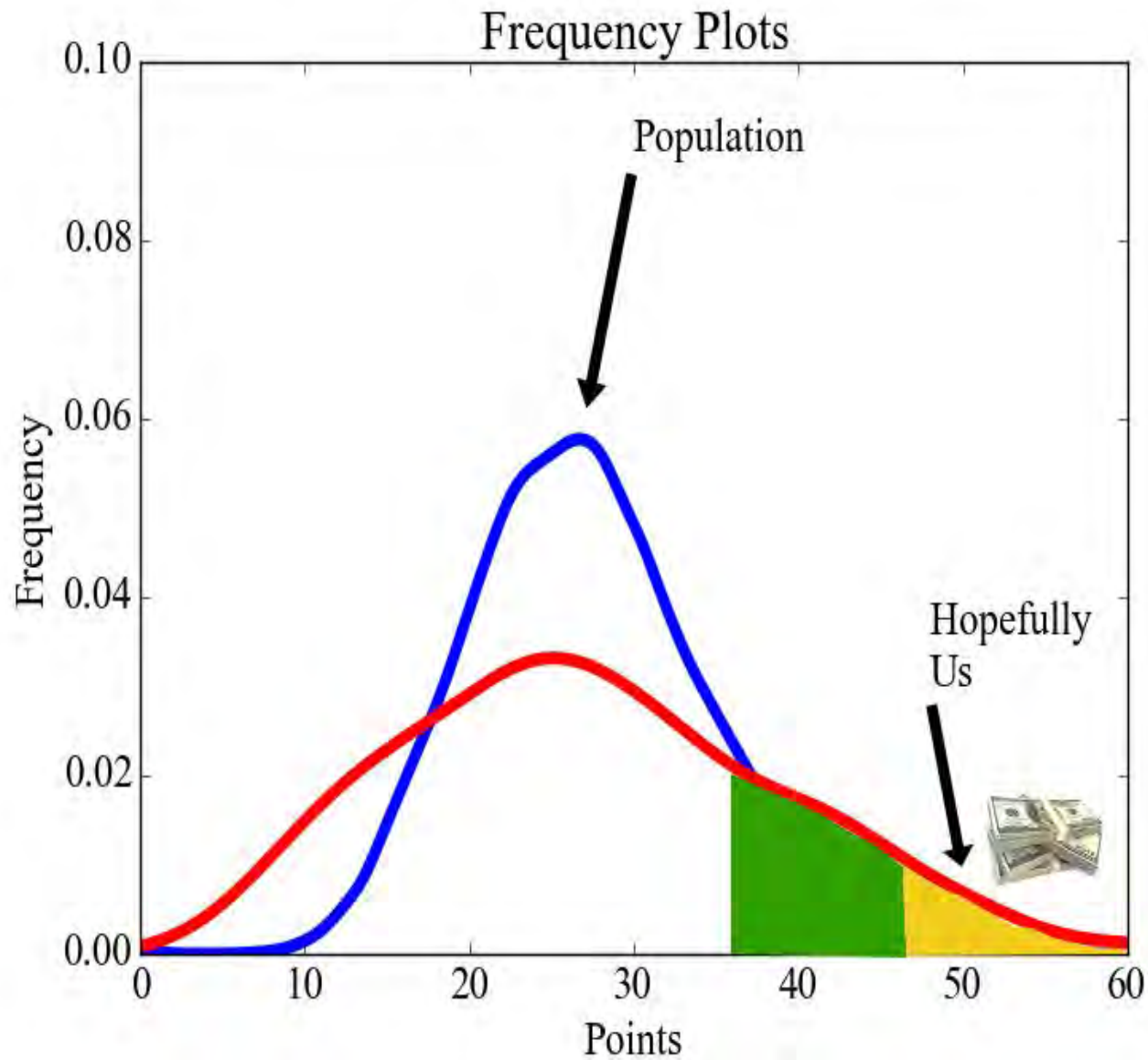


23 points on average

# Need $> 38$ points for a chance to win



# Increase variance to have a chance



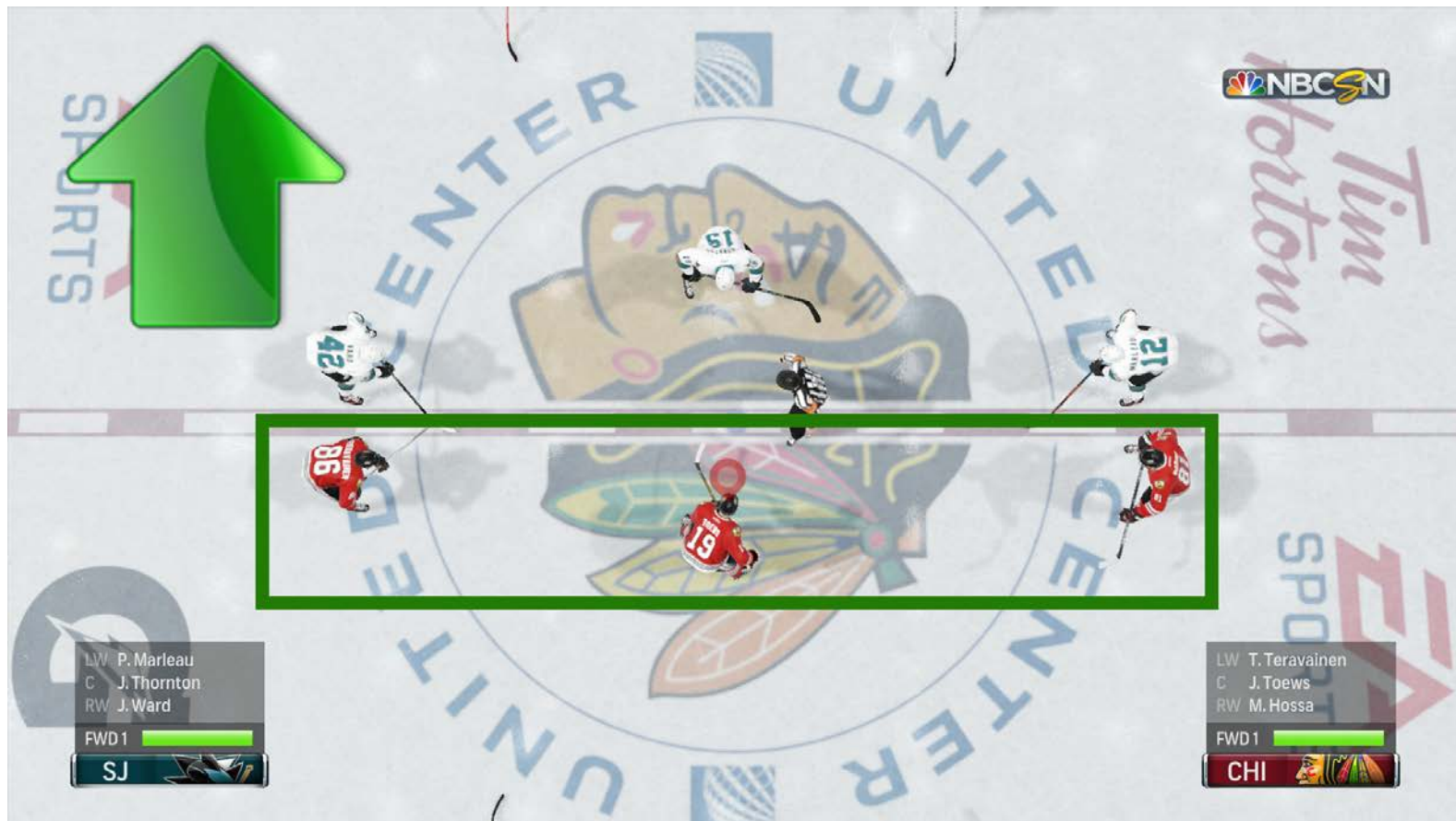


# Structural Correlations - Teams



# Structural Correlations - Lines

- Goal = 3 pt, assist = 2 pt





# Structural Correlations – Lines = Stacking

- At least 1 complete line (3 players per line)
- At least 2 partial lines (at least 2 players per line)

## 1 complete line constraint

$$3v_i \leq \sum_{p \in L_i} x_{pl}, \quad \forall i \in \{1, \dots, N_L\}$$

$$\sum_{i=1}^{N_L} v_i \geq 1$$

$$v_i \in \{0, 1\}, \quad \forall i \in \{1, \dots, N_L\}.$$

## 2 partial lines constraint

$$2w_i \leq \sum_{p \in L_i} x_{pl}, \quad \forall i \in \{1, \dots, N_L\}$$

$$\sum_{i=1}^{N_L} w_i \geq 2$$

$$w_i \in \{0, 1\}, \quad \forall i \in \{1, \dots, N_L\}.$$

# Structural Correlations – Goalie Against Opposing Players



# Structural Correlations – Goalie Against Skaters

- No skater against goalie

No skater against goalie constraint

$$6x_{pl} + \sum_{q \in \text{Opponents}_p} x_{ql} \leq 6, \quad \forall p \in G$$

# Good, but not great chance

Feasible

Team

Goalie

Line

Line

Not  
Against



# Play many diverse Lineups

- Make sure lineup  $l$  has no more than  $\gamma$  players in common with lineups 1 to  $l-1$

Diversity constraint

$$\sum_{p=1}^N x_{pk}^* x_{pl} \leq \gamma, k = 1, \dots, l - 1$$



# Were we able to do it?

NHL \$2K Sniper [\$2,000 Guaranteed]

1st	zlisto	54.50
3rd	zlisto	51.50
9th	zlisto	49.50
23rd	zlisto	46.00
28th	zlisto	45.50
28th	zlisto	45.50

November 15, 2015

NHL \$40K Sniper [\$40,000 Guaranteed]

2nd	zlisto	61.30
21st	zlisto	57.30
21st	zlisto	57.30
40th	zlisto	56.10
42nd	zlisto	55.70
81st	zlisto	54.10

November 16, 2015

NHL \$80K Tuesday Special [\$80,000 Guaranteed]

3rd	zlisto	54.60
6th	zlisto	52.80
7th	zlisto	52.30
10th	zlisto	50.60
11th	zlisto	50.30
15th	zlisto	50.10

November 17, 2015

NHL \$45K Sniper [\$45,000 Guaranteed]

1st	zlisto	52.60
8th	zlisto	49.60
57th	zlisto	45.60
57th	zlisto	45.60
83rd	zlisto	44.60
83rd	zlisto	44.60

November 23, 2015

200 lineups

# Policy Change



200 lineups -> 100 lineups



# Were we able to continue it?

The screenshot shows the GameCenter app interface for an NHL \$12K Sniper tournament. The player 'zlisto' is ranked 1st with a score of 62.50 and a PMR of 0. The tournament has a \$12,000 guaranteed prize pool. A progress bar at the top indicates that 100 lineups have been entered.

Rank	Player	Score	PMR	Prize
1st	zlisto	62.50	0	\$1,000.00
6th	zlisto	58.80	0	\$150.00
8th	zlisto	57.40	0	\$125.00
13th	zlisto	55.80	0	\$80.00
16th	zlisto	55.30	0	\$60.00
20th				



> \$15K

December 12, 2015

100 lineups



How can you do it?



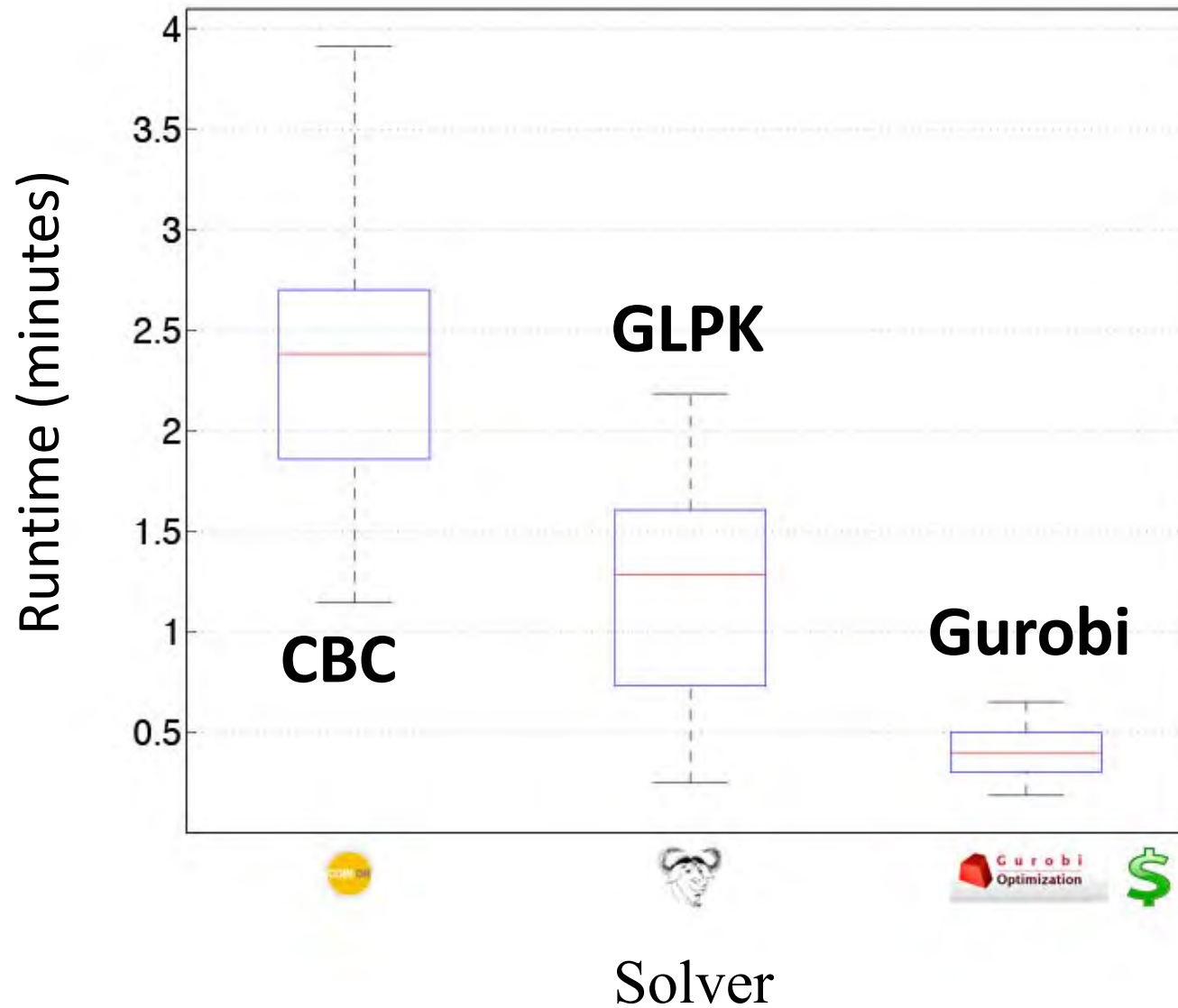
Download Code from Github:

<https://github.com/dscotthunter/Fantasy-Hockey-IP-Code>

```
175 function solve!(skaters, goals, lines, num_overlap, num_skaters, num_goals, centers, wingers, defenders, num_teams, skaters_teams, goals_opponents, team_lines, num_lines, P1_info!  
176     = Model{InfeasibleOptimization})  
177  
178     # Skaters  
179     @variable{Int} x1[1:num_skaters] # Skaters  
180     @variable{Int} x2[1:num_goals] # Goals  
181     # Constraints  
182     @constraint{Int} sum(x1) == num_skaters  
183     @constraint{Int} sum(x2) == num_goals  
184  
185     # Lines  
186     @constraint{Int} sum(centers[i]:skaters_lines[i], 1:num_skaters) == 1  
187     @constraint{Int} sum(wingers[i]:skaters_lines[i], 1:num_skaters) == 1  
188     @constraint{Int} sum(defenders[i]:skaters_lines[i], 1:num_skaters) == 1  
189     # Salaries  
190     @constraint{Int} sum(salary[i]:skaters_lines[i], 1:num_skaters) >= sum(goal[i]:goals_lines[i], 1:num_goals) * M  
191  
192     # Team assignments  
193     @variable{Int} x3[1:num_teams] # Team assignments  
194     @constraint{Int} sum(x3) == num_teams  
195     @constraint{Int} x3[i] >= 0  
196     @constraint{Int} sum(x3) == num_teams  
197  
198     # Goalies  
199     @variable{Int} x4[1:num_lines] # Goalies  
200     @constraint{Int} sum(x4) == num_lines  
201     # Opponents  
202     @variable{Int} x5[1:num_teams] # Opponents  
203     @constraint{Int} sum(x5) == num_teams  
204  
205     # Objective  
206     @objective{Int} minimize sum(x1) + sum(x2) + sum(x3) + sum(x4) + sum(x5)  
207  
208     # Solve  
209     solve!(P1_info!)  
210  
211     # Get results  
212     num_skaters = getvalue(x1)  
213     num_goals = getvalue(x2)  
214     x3 = getvalue(x3)  
215     x4 = getvalue(x4)  
216     x5 = getvalue(x5)  
217  
218     # Return results  
219     return num_skaters, num_goals, x3, x4, x5  
220 end
```

<http://arxiv.org/pdf/1604.01455v1.pdf>

# Performance Time < 30 Minutes



# MIP and Statistics: Inference for the Chilean Earthquake

# The 2010 Chilean Earthquake

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# 6<sup>th</sup> Strongest in Recorded History (8.8)

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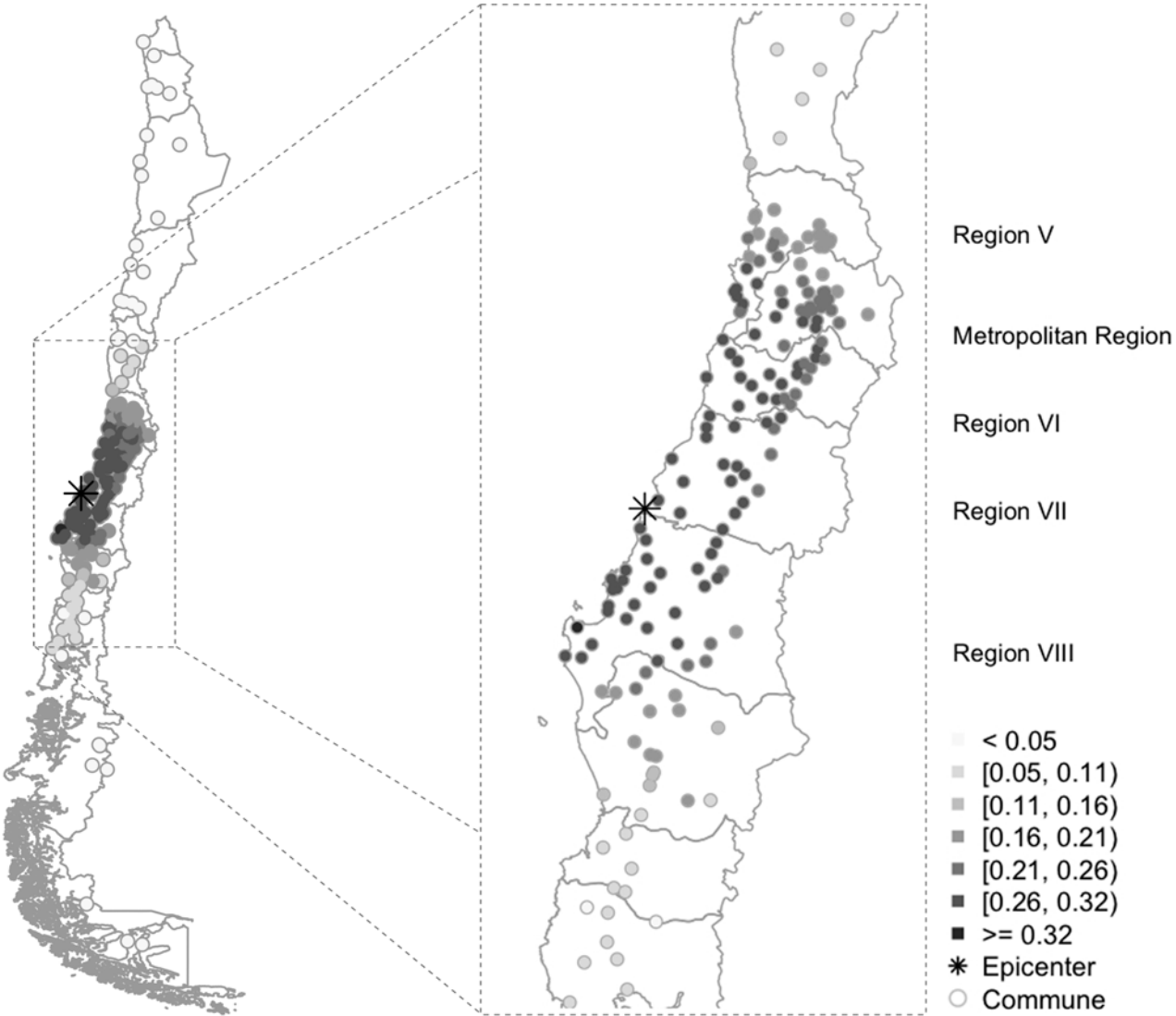
# Impact on Educational Achievement? PSU = SAT

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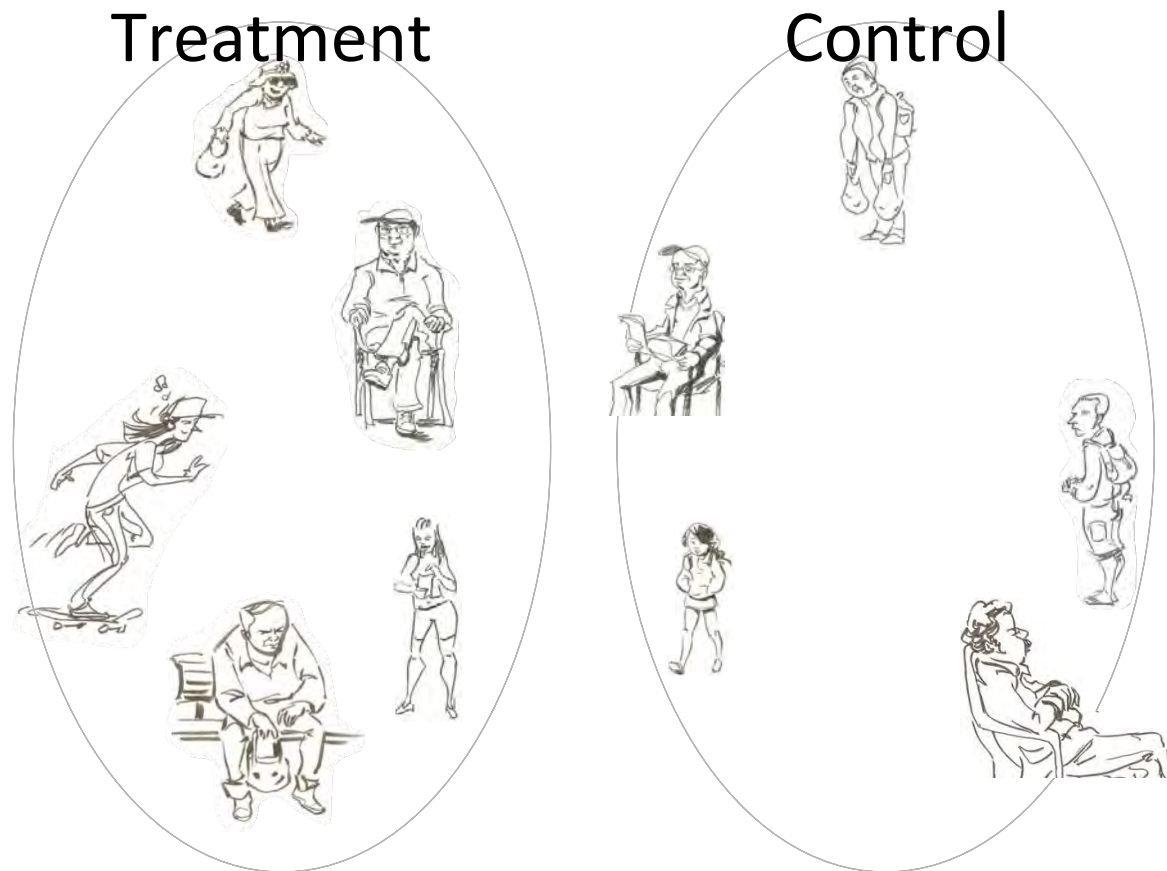
# Earthquake Intensity + Great Demographic Info



# Randomized experiment

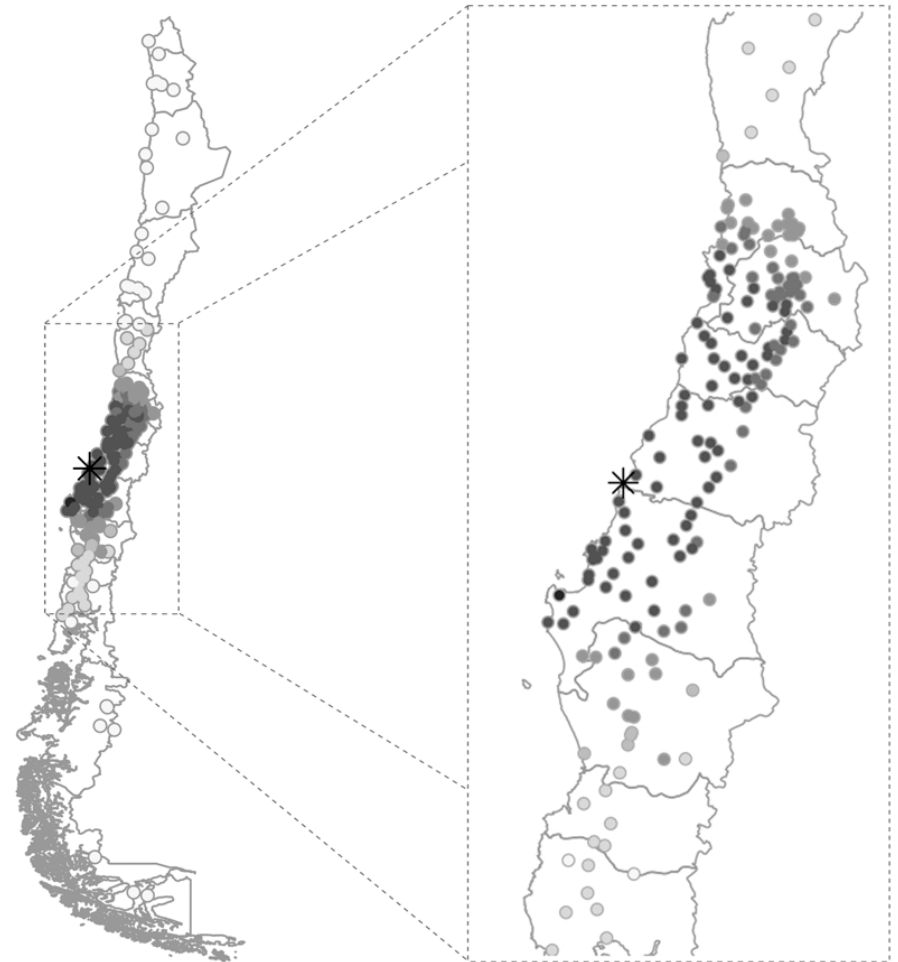
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- Treatment / control have similar characteristics (covariates).



# Covariate Balance Important for Inference

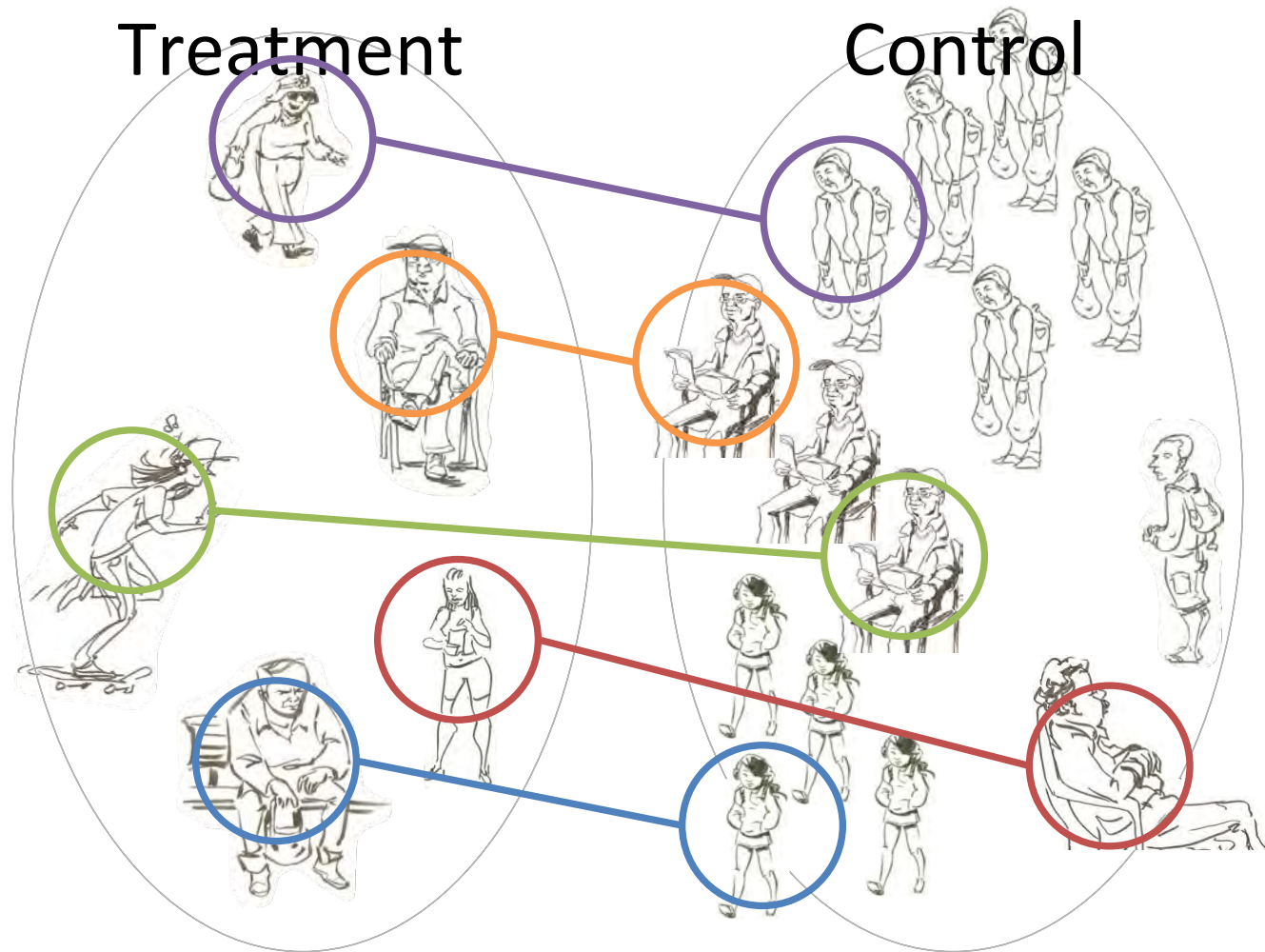
Covariate	Dose	
	1	2
Gender		
Male	462	462
Female	538	538
School SES		
Low	75	75
Mid-low	327	327
Medium	294	294
Mid-high	189	189
High	115	115
Mother's education		
Primary	335	335
Secondary	426	426
Technical	114	114
College	114	114
Missing	11	11
⋮		



# Observational Study: e.g. After Earthquake

- Treatment / control can have different characteristics.

Solution  
=  
Matching?



# Matching

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Treated Units:  $\mathcal{T} = \{t_1, \dots, t_T\}$

Control Units:  $\mathcal{C} = \{c_1, \dots, c_C\}$

Observed Covariates:  $\mathcal{P} = \{p_1, \dots, p_P\}$

Covariate Values:  $\mathbf{x}^t = (x_p^t)_{p \in \mathcal{P}}, \quad t \in \mathcal{T}$

$\mathbf{x}^c = (x_p^c)_{p \in \mathcal{P}}, \quad c \in \mathcal{C}$

# Nearest Neighbor Matching

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$$\text{minimize}_{\mathbf{m}} \quad \sum_{t \in \mathcal{T}} \sum_{c \in \mathcal{C}} \delta_{t,c} m_{t,c}$$

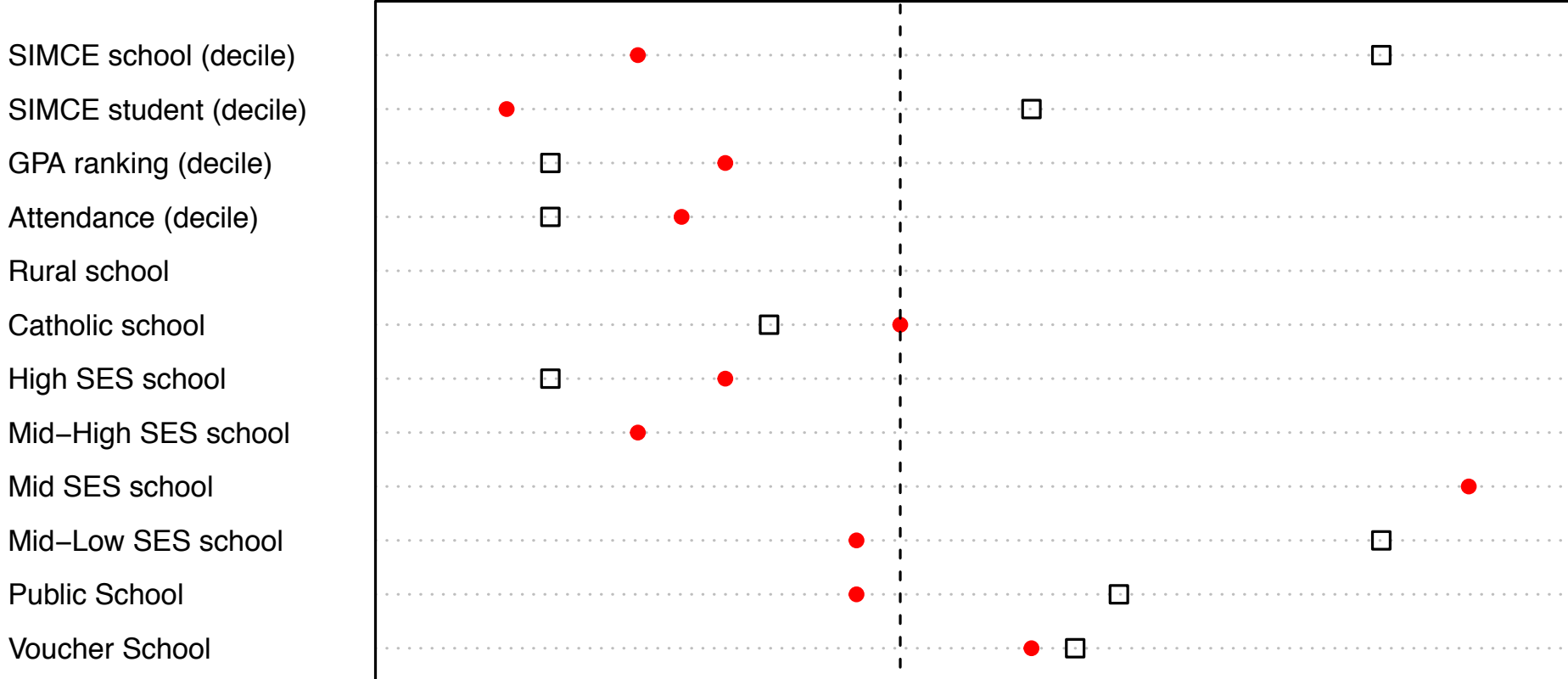
$$\text{subject to} \quad \sum_{c \in \mathcal{C}} m_{t,c} = 1, \quad t \in \mathcal{T}$$

$$\sum_{t \in \mathcal{T}} m_{t,c} \leq 1, \quad c \in \mathcal{C}$$

$$0 \leq m_{t,c} \leq 1 \quad \text{---} m_{t,c} \in \{0, 1\}, \quad t \in \mathcal{T}, c \in \mathcal{C}$$

- e.g.  $\delta_{t,c} = \|\mathbf{x}^t - \mathbf{x}^c\|_2$
- Easy to solve

# Balance Before After Matching





# Maximum Cardinality Matching

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$$\begin{aligned} \max \quad & \sum_{t \in \mathcal{T}} \sum_{c \in \mathcal{C}} m_{t,c} \\ \text{s.t.} \quad & \end{aligned}$$

$$\sum_{t \in \mathcal{T}} m_{t,c} \leq 1,$$

$$\sum_{c \in \mathcal{C}} m_{t,c} \leq 1,$$

$$\sum_{t \in \mathcal{T}_{p,k}} \sum_{c \notin \mathcal{C}_{p,k}} m_{t,c} = \sum_{t \notin \mathcal{T}_{p,k}} \sum_{c \in \mathcal{C}_{p,k}} m_{t,c}$$

$$m_{t,c} \in \{0, 1\}$$

$$\mathcal{K}(p) = \{\mathbf{x}_p^c\}_{c \in \mathcal{P}} \cup \{\mathbf{x}_p^t\}_{t \in \mathcal{T}}$$

$$\mathcal{C}_{p,k} = \{c \in \mathcal{C} : \mathbf{x}_p^c = k\}$$

$$\mathcal{T}_{p,k} = \{t \in \mathcal{T} : \mathbf{x}_p^t = k\}$$

$$\forall c \in \mathcal{C}$$

$$\forall t \in \mathcal{T}$$

$$\forall p \in \mathcal{P}, k \in \mathcal{K}(p)$$

$$\forall t \in \mathcal{T}, c \in \mathcal{C}.$$

- Very hard to solve ( and very hard to understand! )

# Advanced Maximum Cardinality Matching

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$$\max \sum_{t \in \mathcal{T}} x_t$$

*s.t.*

$$\sum_{t \in \mathcal{T}} x_t = \sum_{c \in \mathcal{C}} y_c,$$

$$\sum_{t \in \mathcal{T}_{p,k}} x_t = \sum_{c \in \mathcal{C}_{p,k}} y_c, \quad \forall p \in \mathcal{P}, \quad k \in \mathcal{K}(p)$$

$$x_t \in \{0, 1\} \quad \forall t \in \mathcal{T}$$

$$y_c \in \{0, 1\} \quad \forall c \in \mathcal{C}.$$

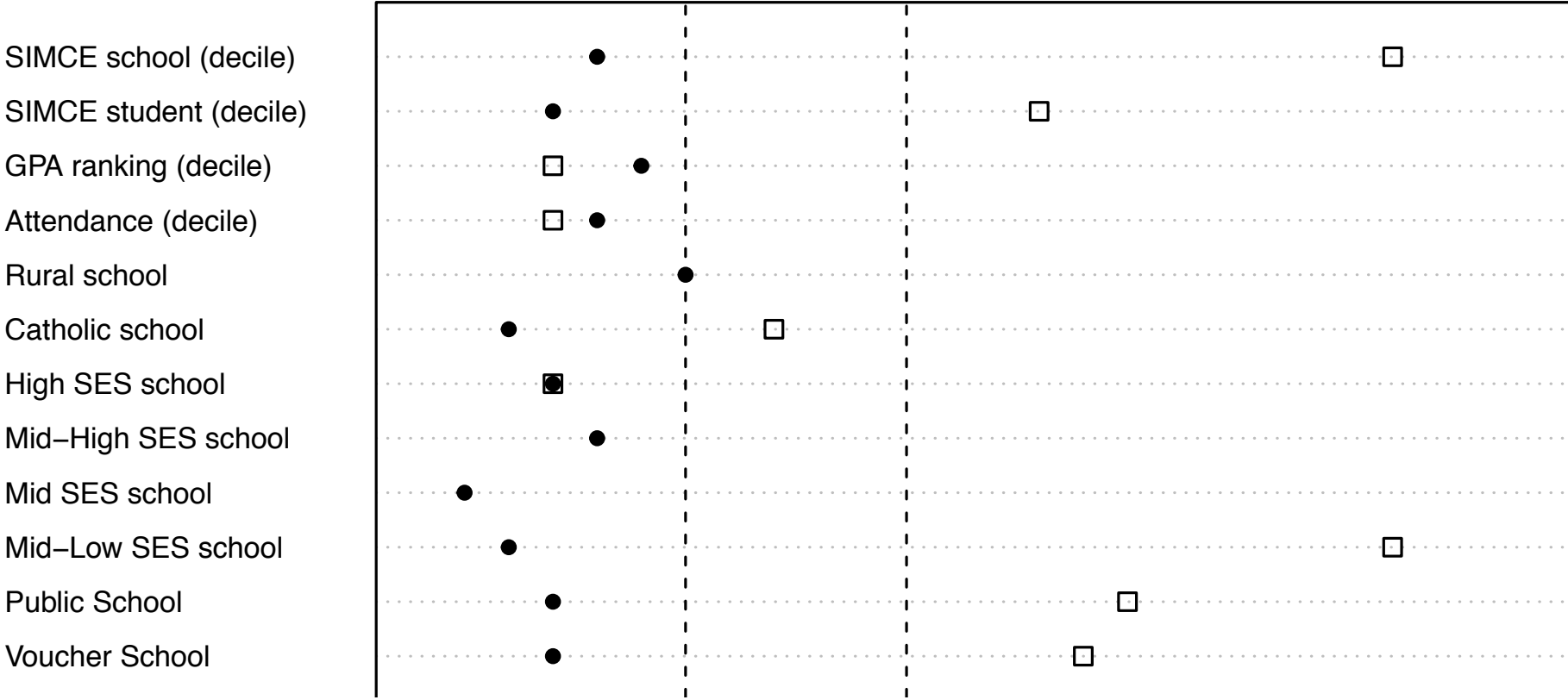
$$\mathcal{K}(p) = \{\mathbf{x}_p^c\}_{c \in \mathcal{P}} \cup \{\mathbf{x}_p^t\}_{t \in \mathcal{T}}$$

$$\mathcal{C}_{p,k} = \{c \in \mathcal{C} : \mathbf{x}_p^c = k\}$$

$$\mathcal{T}_{p,k} = \{t \in \mathcal{T} : \mathbf{x}_p^t = k\}$$

- Matching without matching variables
- Easy to solve: Small, but inherits matching properties

# Balance Before After Cardinality Matching



# Can Also do Multiple Doses

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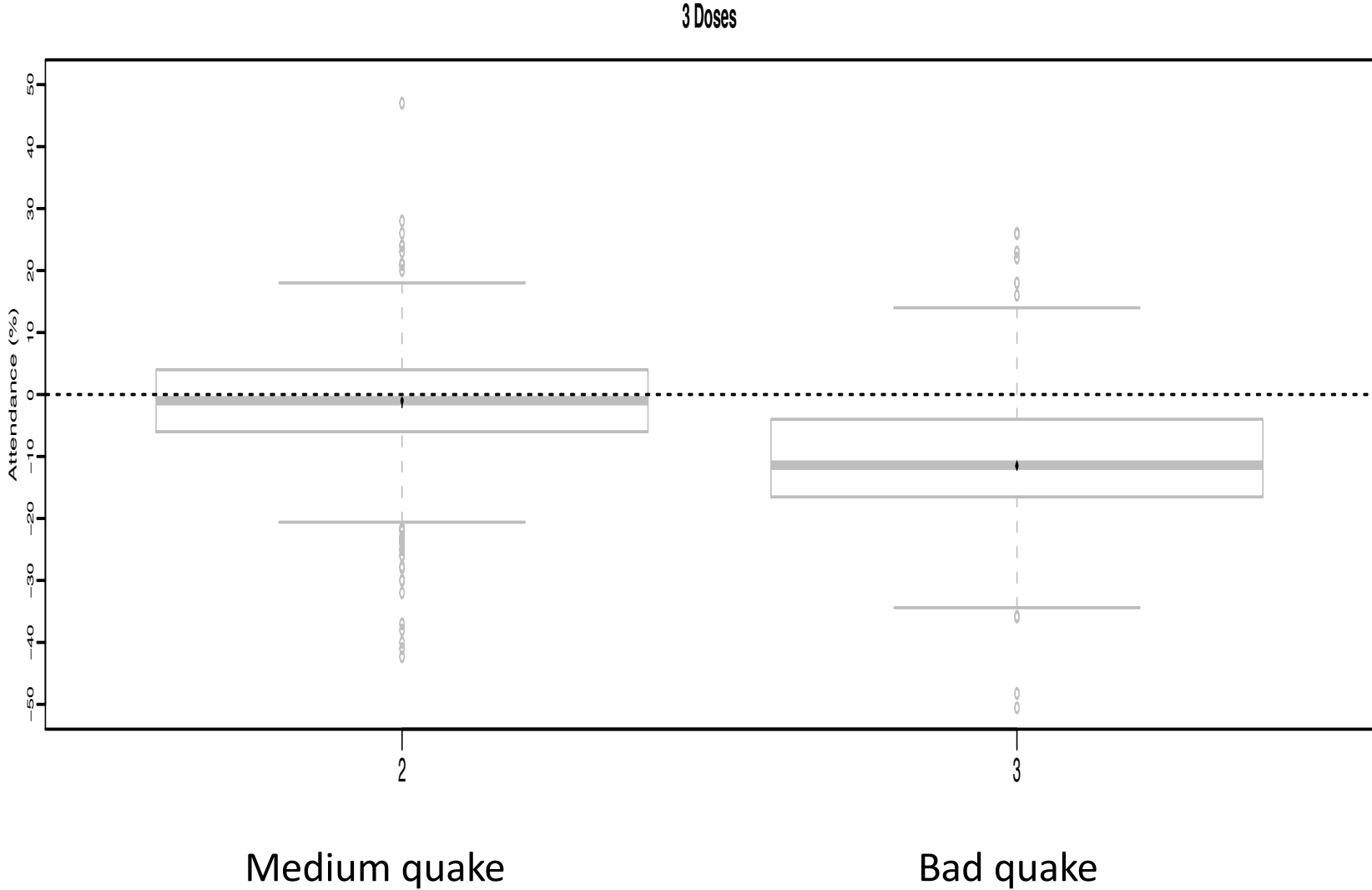
- Dose
  1. No quake
  2. Medium quake
  3. Bad quake

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Covariate	Dose		
	1	2	3
Gender			
Male	462	462	462
Female	538	538	538
School SES			
Low	75	75	75
Mid-low	327	327	327
Medium	294	294	294
Mid-high	189	189	189
High	115	115	115
Mother's education			
Primary	335	335	335
Secondary	426	426	426
Technical	114	114	114
College	114	114	114
Missing	11	11	11
⋮			

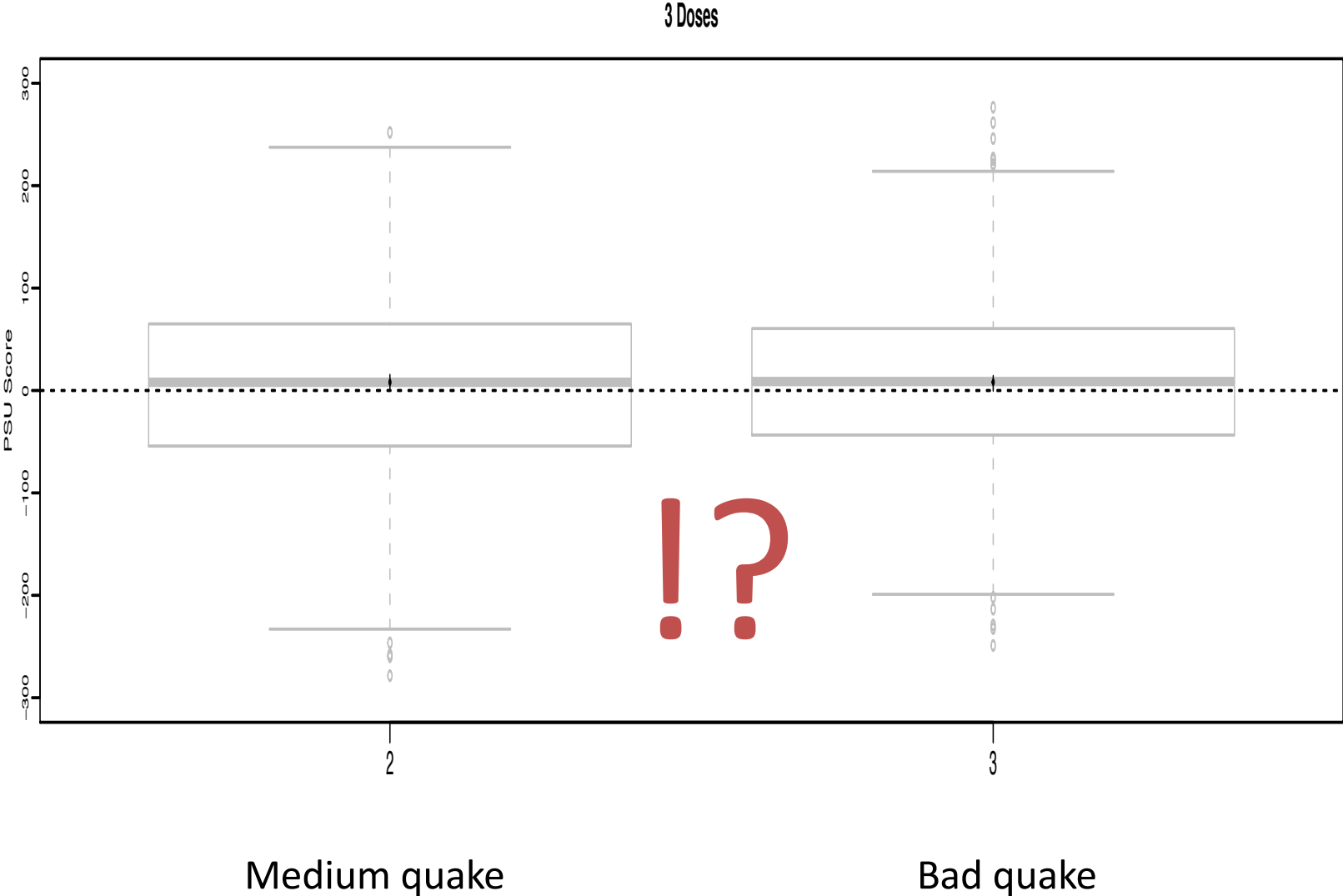
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# Relative (To no Quake) Attendance Impact





# Relative (To no Quake) PSU Score Impact



MIP and Marketing:  
Chewbacca or BB-8?

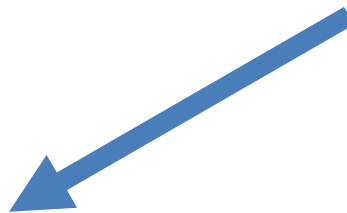
# Adaptive Preference Questionnaires



Feature	SX530	RX100
Zoom	50x	3.6x
Prize	\$249.99	\$399.99
Weight	15.68 ounces	7.5 ounces
Prefer	<input checked="" type="checkbox"/>	<input type="checkbox"/>



Feature	TG-4	G9
Waterproof	Yes	No
Prize	\$249.99	\$399.99
Weight	7.36 lb	7.5 lb
Prefer	<input type="checkbox"/>	<input checked="" type="checkbox"/>



Feature	TG-4	Galaxy 2
Waterproof	Yes	No
Prize	\$249.99	\$399.99
Viewfinder	Electronic	Optical
Prefer	<input checked="" type="checkbox"/>	<input type="checkbox"/>



We recommend:



# Choice-based Conjoint Analysis (CBCA)

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Feature	Chewbacca	BB-8
Wookiee	Yes	No
Droid	No	Yes
Blaster	Yes	No
I would buy toy	<input checked="" type="checkbox"/>	<input type="checkbox"/>

$$\begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} = x^2$$

Product Profile       $x^1$        $x^2$

# Preference Model and Geometric Interpretation

- Utilities for 2 products, d features, logit model

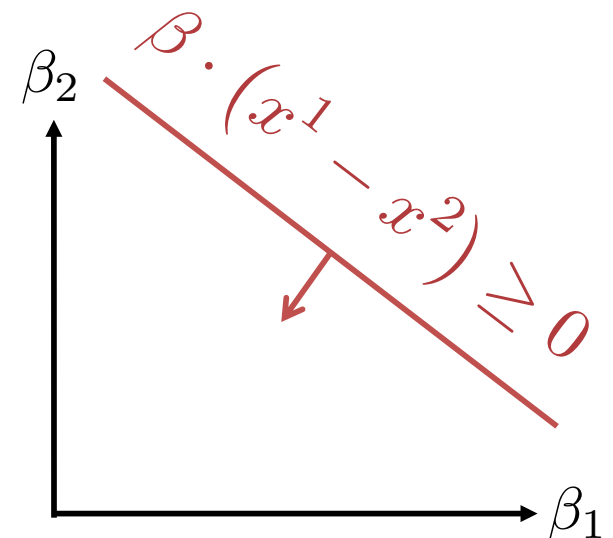
$$U_1 = \beta \cdot x^1 + \epsilon_1 = \sum_{i=1}^d \beta_i x_i^1 + \epsilon_1$$

$$U_2 = \beta \cdot x^2 + \epsilon_2 = \sum_{i=1}^d \beta_i x_i^2 + \epsilon_2$$

part-worths  $\uparrow$   
product profile  $\uparrow$  noise (gumbel)  $\uparrow$

- Utility maximizing customer
  - Geometric interpretation of preference for product 1 **without error**

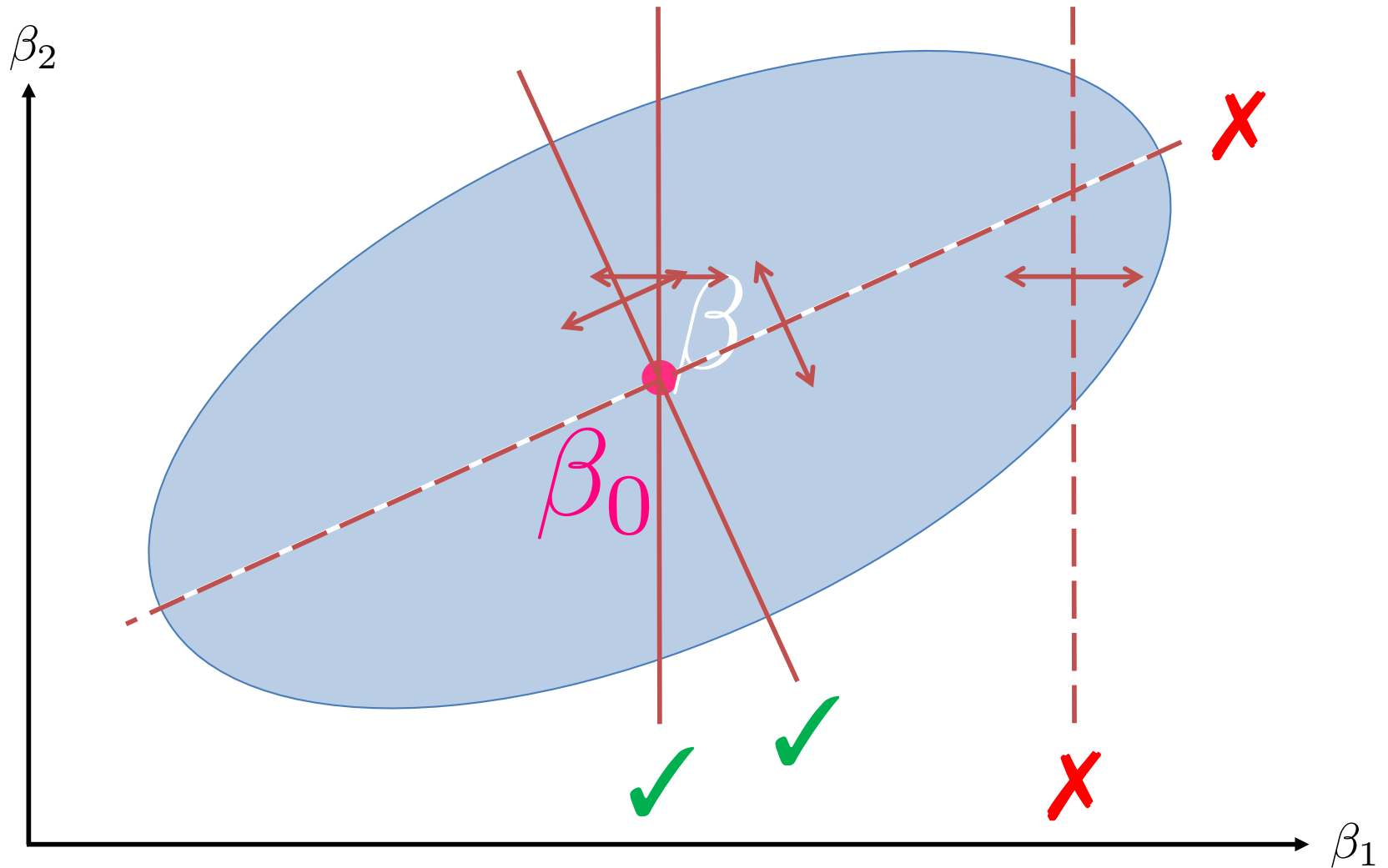
$$x^1 \succeq x^2 \Leftrightarrow U_1 \geq U_2$$





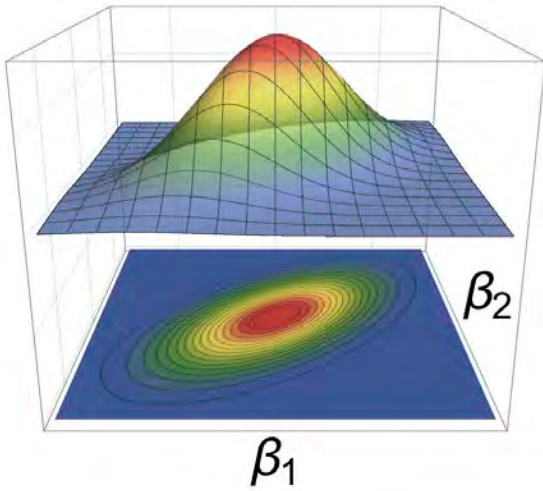
# Next Question = Minimize (Expected) Volume

Good Estimator? for  $\beta$ ? ~~Central tendency~~  $\beta_0$

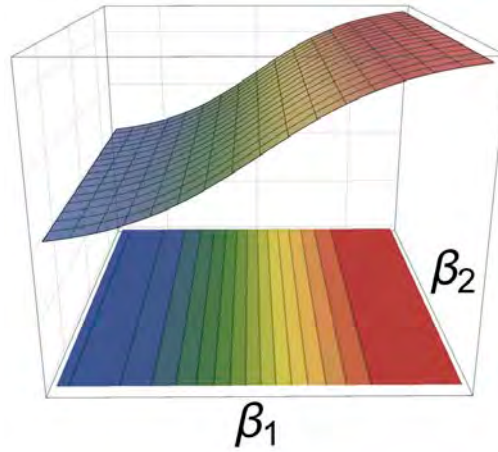


# With Error = Volume of Ellipsoid $f(x^1, x^2)$

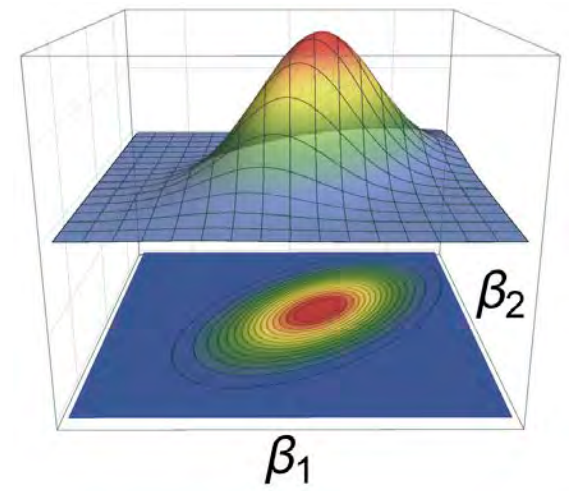
Prior distribution



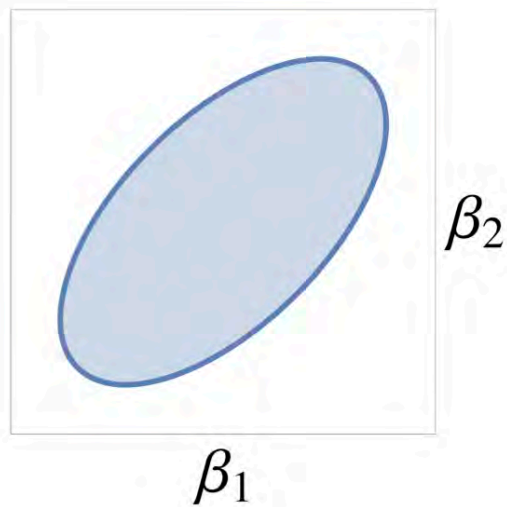
Answer likelihood



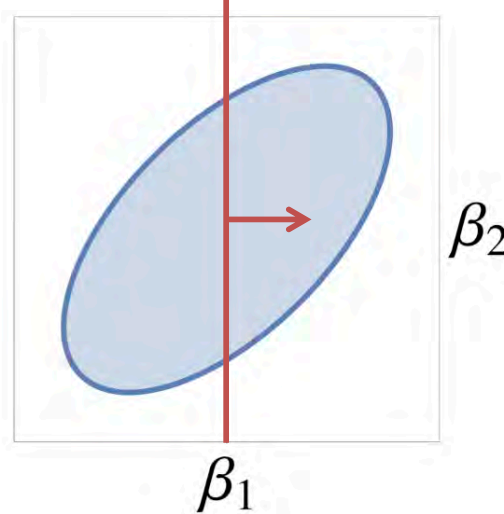
Posterior distribution



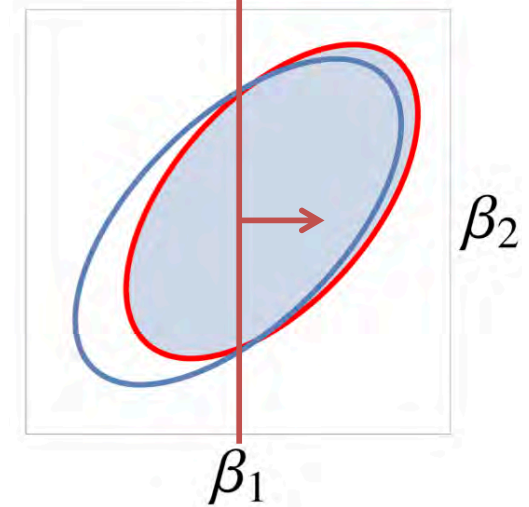
Prior ellipsoid



Question/Answer



Posterior ellipsoid



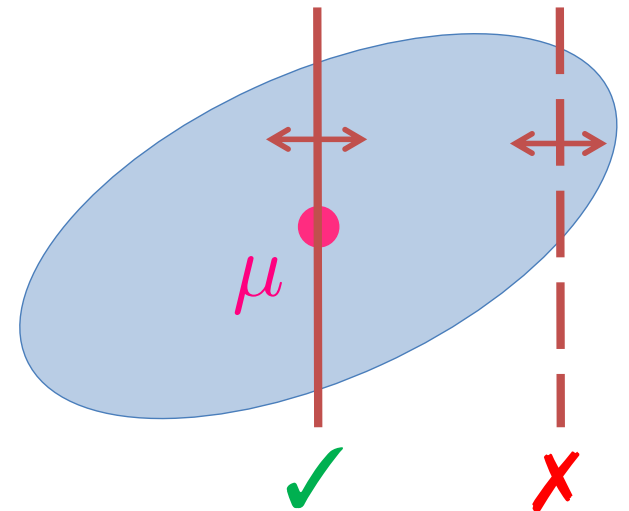
# Rules of Thumb Still Good For Ellipsoid Volume

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$$(\beta - \mu)' \cdot \Sigma^{-1} \cdot (\beta - \mu) \leq r$$

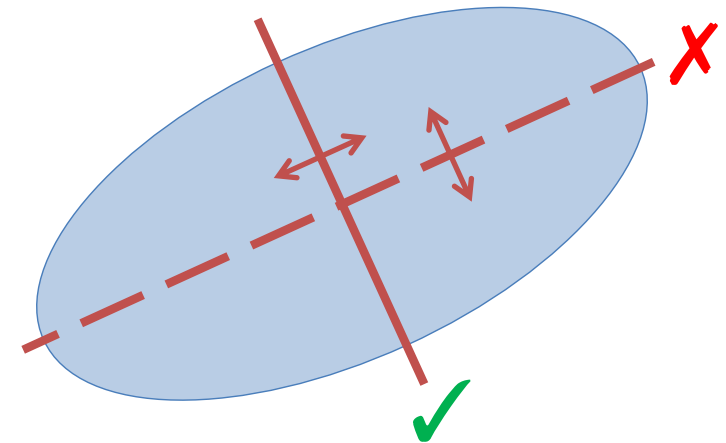
- Choice balance:
  - Minimize **distance** to center

$$\mu \cdot (x^1 - x^2)$$



- Postchoice symmetry:
  - Maximize **variance** of question

$$(x^1 - x^2)' \cdot \Sigma \cdot (x^1 - x^2)$$

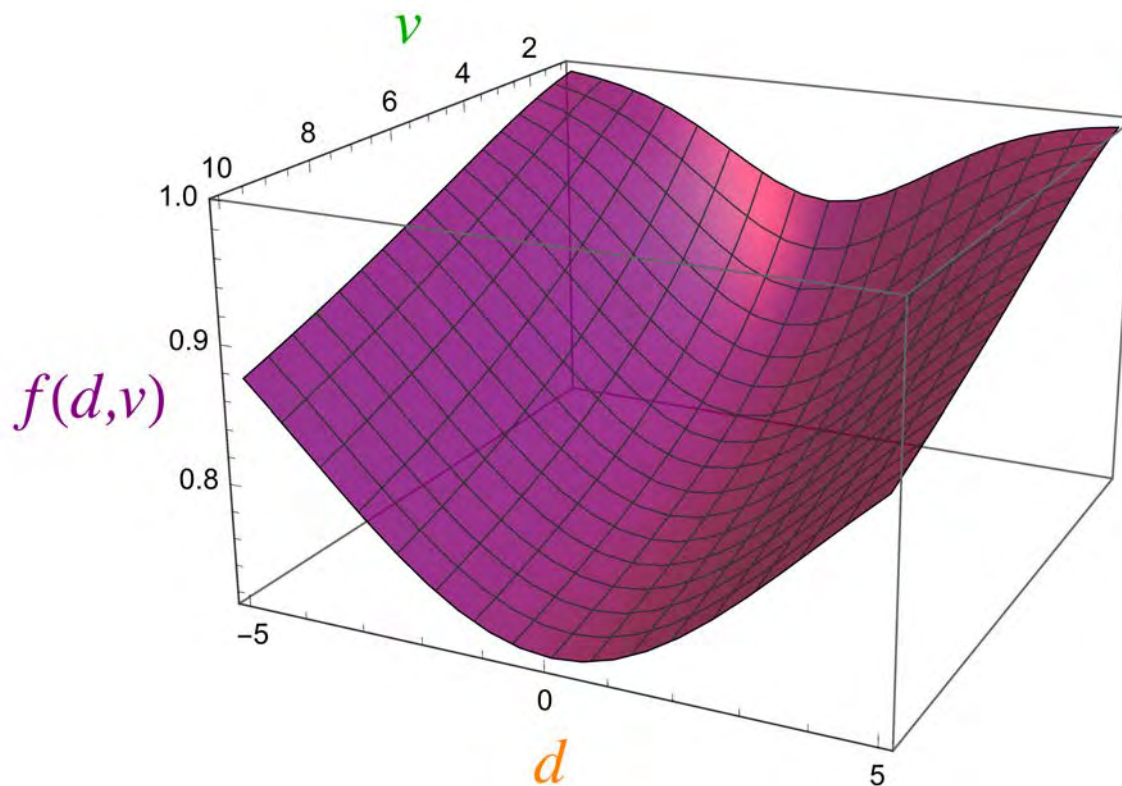


# “Simple” Formula for Expected Volume

- Expected Volume = Non-convex function  $f(d, v)$  of

distance:  $d := \mu \cdot (x^1 - x^2)$

variance:  $v := (x^1 - x^2)' \cdot \Sigma \cdot (x^1 - x^2)$



Can evaluate  $f(d, v)$   
with 1-dim integral 😊

# Optimization Model

---

min

$$f(d, v)$$

~~X~~

s.t.

$$\mu \cdot (x^1 - x^2) = d \quad \checkmark$$

$$(x^1 - x^2)' \cdot \Sigma \cdot (x^1 - x^2) = v \quad \checkmark \times$$

$$A^1 x^1 + A^2 x^2 \leq b \quad \checkmark$$

Formulation trick:

linearize  $x_i^k \cdot x_j^l$

$$x^1 \neq x^2 \quad \checkmark \times$$

$$x^1, x^2 \in \{0, 1\}^n$$

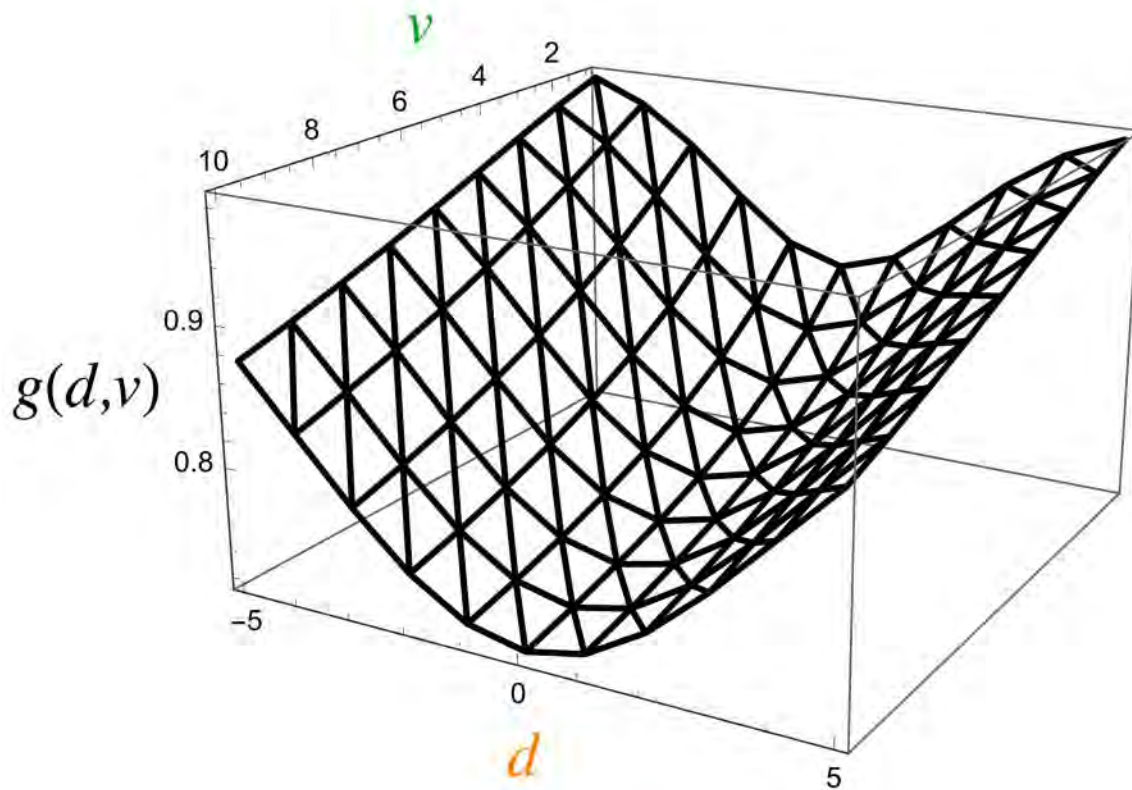


# Technique 2: Piecewise Linear Functions

- D-efficiency = Non-convex function  $f(d, v)$

distance:  $d := \mu \cdot (x^1 - x^2)$

variance:  $v := (x^1 - x^2)' \cdot \Sigma \cdot (x^1 - x^2)$



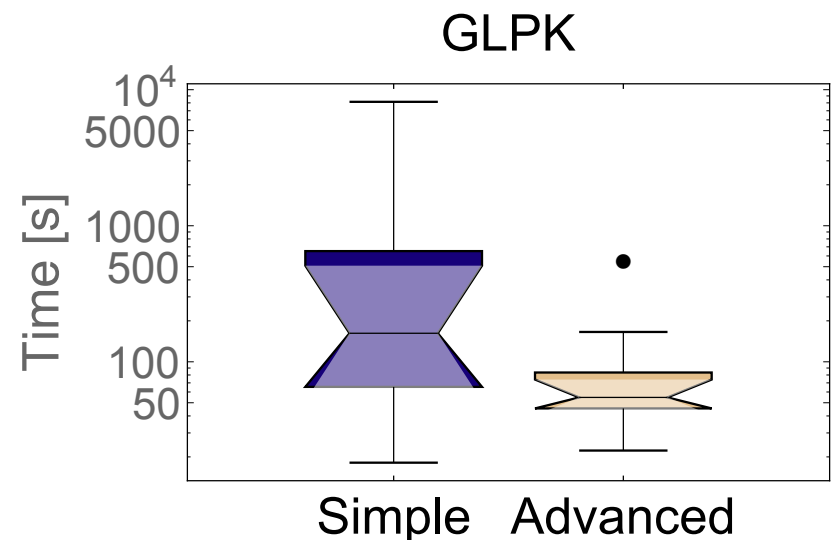
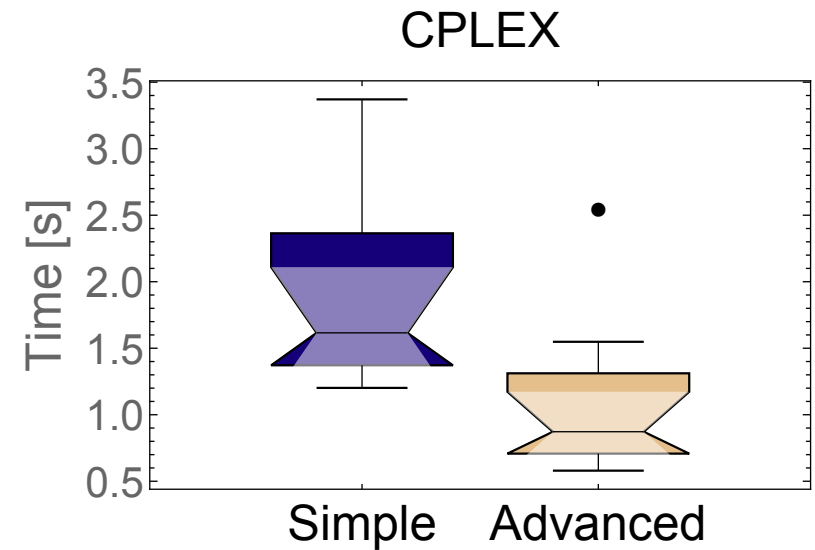
Can evaluate  $f(d, v)$   
with 1-dim integral 😊

Piecewise Linear  
Interpolation

MIP formulation

# Computational Performance

- Advanced formulations provide an computational advantage
- Advantage is significantly more important for free solvers
- State of the art commercial solvers can be significantly better than free solvers
- Still, free is free!



# Summary and Main Messages

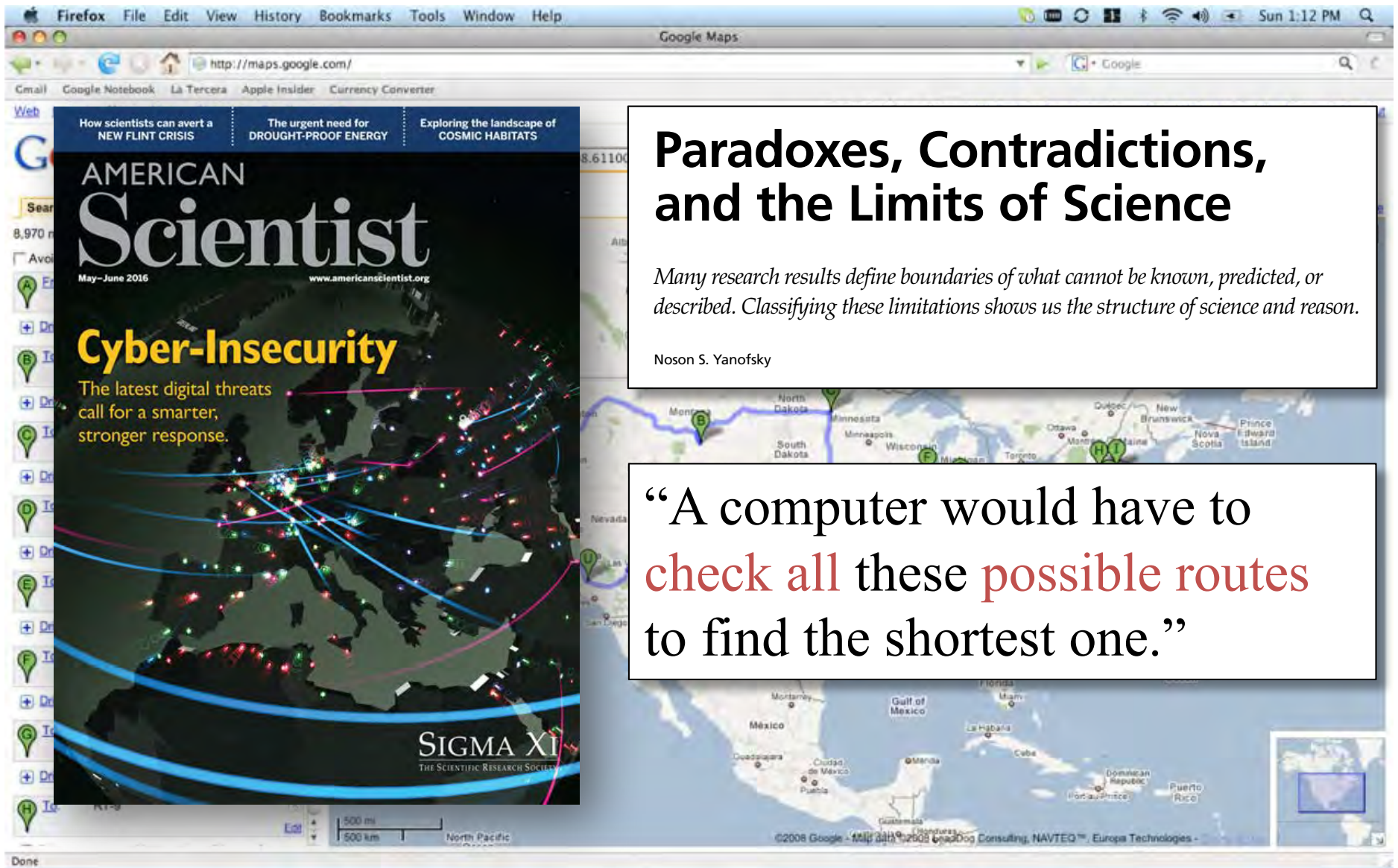
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- Always choose Chewbacca!
- How to YOU use MIP?
  - Study for the final!
  - Use JuMP and Julia Opt.
  - Write “good” formulations.
  - Use your domain expertise.



How Hard is MIP?

# How hard is MIP: Traveling Salesman Problem ?



The image is a screenshot of a Firefox browser window. The address bar shows 'http://maps.google.com/'. The browser's menu bar includes 'Firefox', 'File', 'Edit', 'View', 'History', 'Bookmarks', 'Tools', 'Window', and 'Help'. The page content is split into two main sections. On the left, the cover of the May-June 2016 issue of 'AMERICAN Scientist' is displayed. The cover features a dark background with a map of North America overlaid with a complex network of colorful lines (red, blue, green, yellow) representing data or connections. The main headline on the cover is 'Cyber-Insecurity' in large yellow letters, with a sub-headline 'The latest digital threats call for a smarter, stronger response.' in white. Above the main title, there are three smaller headlines: 'How scientists can avert a NEW FLINT CRISIS', 'The urgent need for DROUGHT-PROOF ENERGY', and 'Exploring the landscape of COSMIC HABITATS'. At the bottom of the cover, it says 'SIGMA XI THE SCIENTIFIC RESEARCH SOCIETY'. On the right side of the browser window, a Google Maps interface is visible, showing a map of the United States with several green location pins. A white text box is overlaid on the map, containing the title 'Paradoxes, Contradictions, and the Limits of Science' in bold black font. Below the title is a quote in italics: 'Many research results define boundaries of what cannot be known, predicted, or described. Classifying these limitations shows us the structure of science and reason.' followed by the author's name 'Noson S. Yanofsky'. A second white text box is overlaid on the map below the first one, containing a quote in black font: 'A computer would have to check all these possible routes to find the shortest one.' The word 'check' is in red, and 'possible routes' is in red. The browser's status bar at the bottom shows 'Done'.

AMERICAN  
**Scientist**  
May-June 2016  
www.americanscientist.org

**Cyber-Insecurity**  
The latest digital threats call for a smarter, stronger response.

SIGMA XI  
THE SCIENTIFIC RESEARCH SOCIETY

## Paradoxes, Contradictions, and the Limits of Science

*Many research results define boundaries of what cannot be known, predicted, or described. Classifying these limitations shows us the structure of science and reason.*

Noson S. Yanofsky

“A computer would have to check all these possible routes to find the shortest one.”



# MIP = Avoid Enumeration

---

- Number of tours for 49 cities =  $48!/2 \approx 10^{60}$
- Fastest supercomputer  $\approx 10^{17}$  flops
- Assuming one floating point operation per tour:  
>  $10^{35}$  years  $\approx 10^{25}$  times the age of the universe!
- How long does it take on an iphone?
  - Less than a second!
  - 4 iterations of **cutting plane** method!
  - Dantzig, Fulkerson and Johnson 1954 did it by hand!
  - For more info see tutorial in ConcordeTSP app
  - **Cutting planes** are the key for effectively solving (even NP-hard) MIP problems in practice.

# 50+ Years of MIP = Significant Solver Speedups

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- Algorithmic Improvements (Machine Independent):
  - **CPLEX** v1.2 (1991) – v11 (2007): 29,000x speedup
  - Gurobi v1 (2009) – v6.5 (2015): 48.7x speedup
  - Commercial, but free for academic use
- (Reasonably) effective free / open source solvers:
  - GLPK, **COIN-OR (CBC)** and SCIP (only for non-commercial)
- Easy to use, fast and versatile modeling languages
  - Julia based JuMP modelling language
  - <http://julialang.org>
  - <http://www.juliaopt.org>

# Technique 1: Binary Quadratic $x^1, x^2 \in \{0, 1\}^n$

---

$$(x^1 - x^2)' \cdot \Sigma \cdot (x^1 - x^2) = v$$

$$X_{i,j}^l = x_i^l \cdot x_j^l \quad (l \in \{1, 2\}, \quad i, j \in \{1, \dots, n\}) :$$

$$X_{i,j}^l \leq x_i^l, \quad X_{i,j}^l \leq x_j^l, \quad X_{i,j}^l \geq x_i^l + x_j^l - 1, \quad X_{i,j}^l \geq 0$$

$$W_{i,j} = x_i^1 \cdot x_j^2 :$$

$$W_{i,j} \leq x_i^1, \quad W_{i,j} \leq x_j^2, \quad W_{i,j} \geq x_i^1 + x_j^2 - 1, \quad W_{i,j} \geq 0$$

$$\sum_{i,j=1}^n (X_{i,j}^1 + X_{i,j}^2 - W_{i,j} - W_{j,i}) \Sigma_{i,j} = v$$

# Technique 1: Binary Quadratic $x^1, x^2 \in \{0, 1\}^n$

---

$$x^1 \neq x^2 \iff \|x^1 - x^2\|_2^2 \geq 1$$

$$X_{i,j}^l = x_i^l \cdot x_j^l \quad (l \in \{1, 2\}, \quad i, j \in \{1, \dots, n\}) :$$

$$X_{i,j}^l \leq x_i^l, \quad X_{i,j}^l \leq x_j^l, \quad X_{i,j}^l \geq x_i^l + x_j^l - 1, \quad X_{i,j}^l \geq 0$$

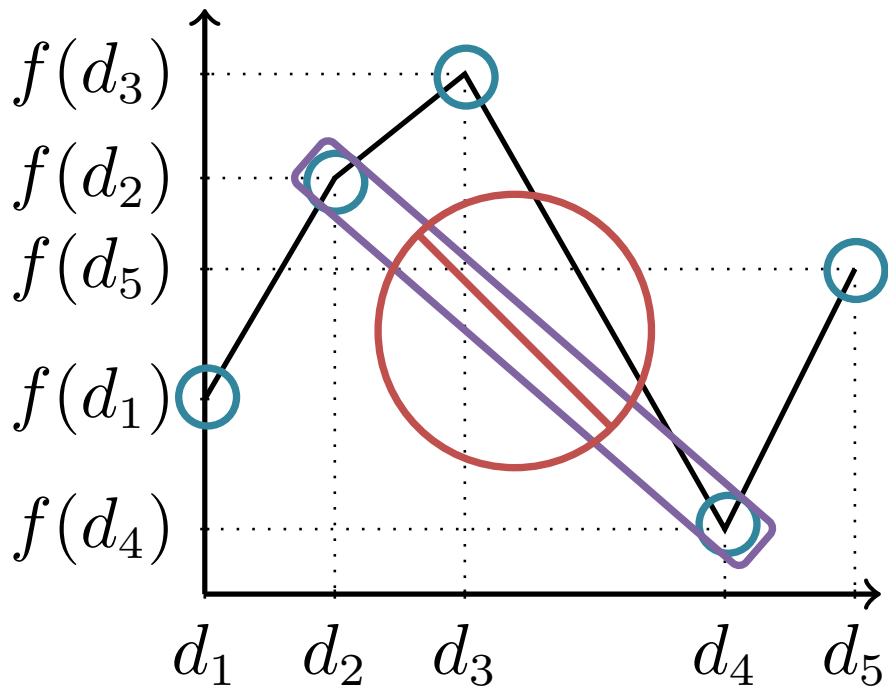
$$W_{i,j} = x_i^1 \cdot x_j^2 :$$

$$W_{i,j} \leq x_i^1, \quad W_{i,j} \leq x_j^2, \quad W_{i,j} \geq x_i^1 + x_j^2 - 1, \quad W_{i,j} \geq 0$$

$$\sum_{i,j=1}^n (X_{i,j}^1 + X_{i,j}^2 - W_{i,j} - W_{j,i}) \geq 1$$

# Simple Formulation for Univariate Functions

$$z = f(x)$$



Size =  $O(\# \text{ of segments})$

Non-Ideal: Fractional Extreme Points

$$\begin{pmatrix} x \\ z \end{pmatrix} = \sum_{j=1}^5 \begin{pmatrix} d_j \\ f(d_j) \end{pmatrix} \lambda_j$$

$$1 = \sum_{j=1}^5 \lambda_j, \quad \lambda_j \geq 0$$

$$y \in \{0, 1\}^4, \quad \sum_{i=1}^4 y_i = 1$$

$$0 \leq \lambda_1 \leq y_1$$

$$0 \leq \lambda_2 \leq y_1 + y_2$$

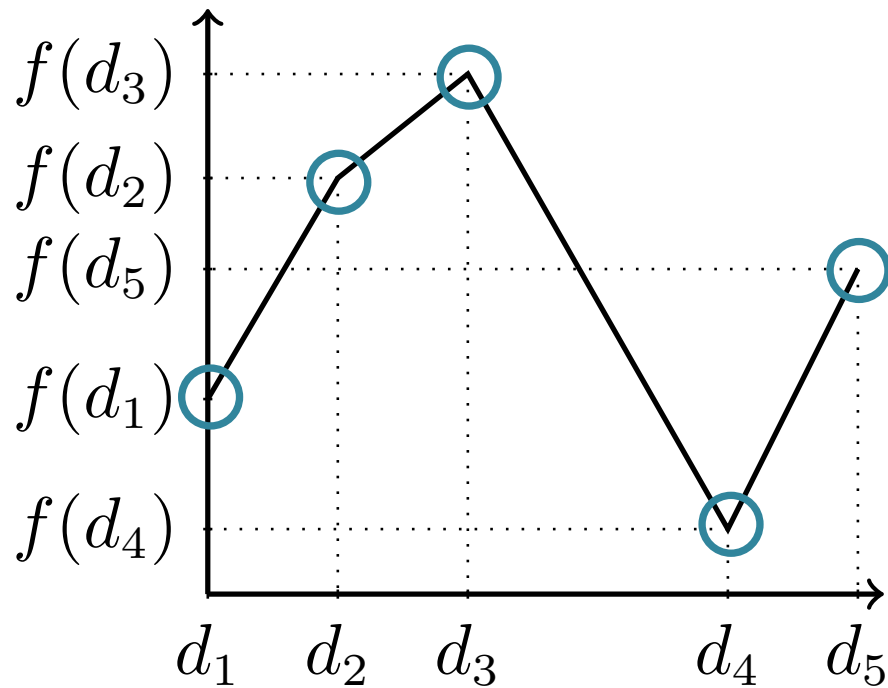
$$0 \leq \lambda_3 \leq y_2 + y_3$$

$$0 \leq \lambda_4 \leq y_3 + y_4$$

$$0 \leq \lambda_5 \leq y_4$$

# Advanced Formulation for Univariate Functions

$$z = f(x)$$



Size =  $O(\log_2 \# \text{ of segments})$

Ideal: Integral Extreme Points

$$\begin{pmatrix} x \\ z \end{pmatrix} = \sum_{j=1}^5 \begin{pmatrix} d_j \\ f(d_j) \end{pmatrix} \lambda_j$$

$$1 = \sum_{j=1}^5 \lambda_j, \quad \lambda_j \geq 0$$

$$y \in \{0, 1\}^2$$

$$0 \leq \lambda_1 + \lambda_5 \leq 1 - y_1$$

$$0 \leq \lambda_3 \leq y_1$$

$$0 \leq \lambda_4 + \lambda_5 \leq 1 - y_2$$

$$0 \leq \lambda_1 + \lambda_2 \leq y_2$$